

CS164: Simplicial Complexes, Homology, Persistence



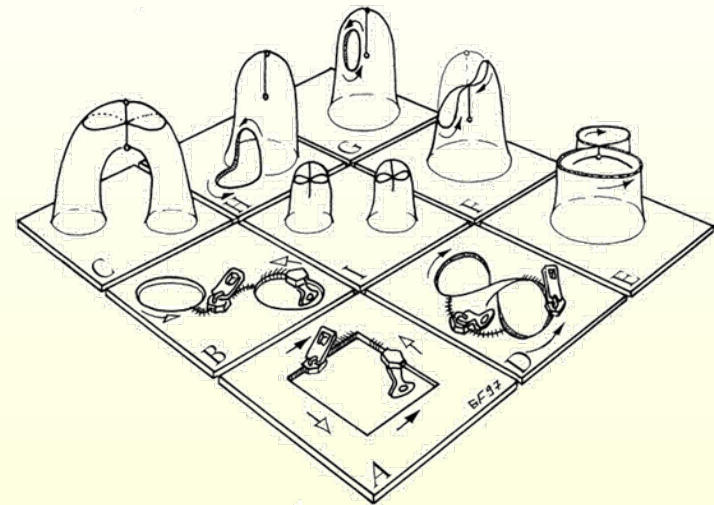
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[Most slides: Afra Zomorodian]

Why Topology?

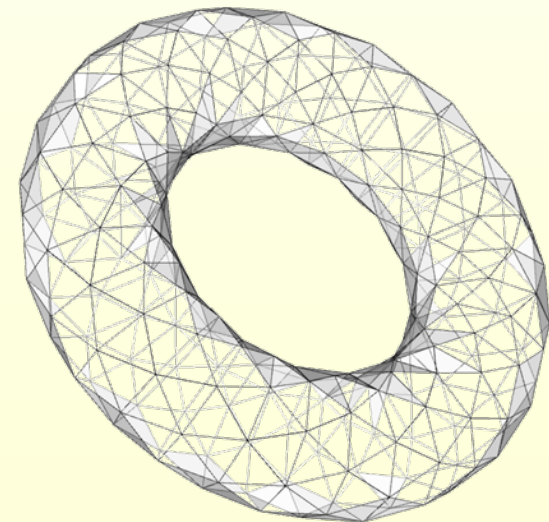
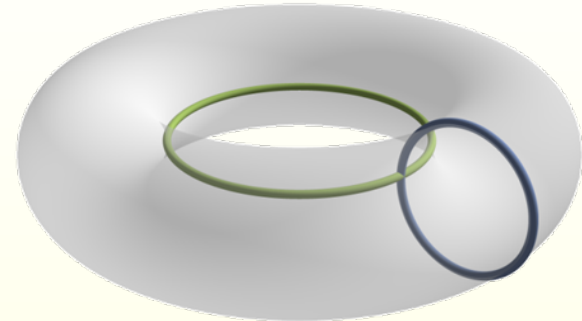
- Topology is the study of connectivity
- It investigates properties of shapes and spaces that are less sensitive to exact metric information
- As a consequence, it produces shape and space invariants that are robust to many kinds of deformation and noise



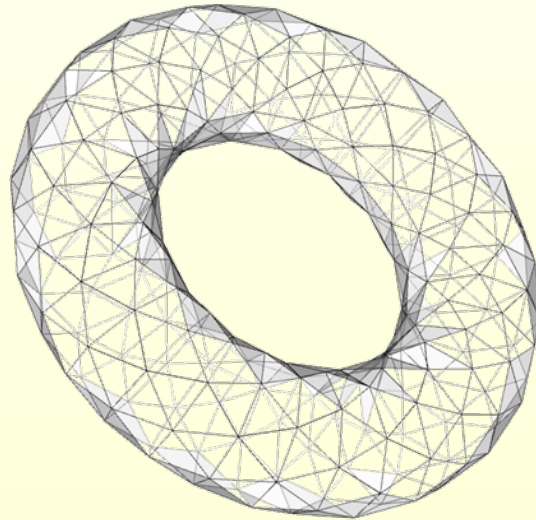
cross-handle = 2 cross-caps

Computational Representations of Topology

- We need to find discrete representations of infinite, continuous topological spaces
- We need to develop efficient algorithms for the manipulation of such representations, as well as for extracting topological information from them

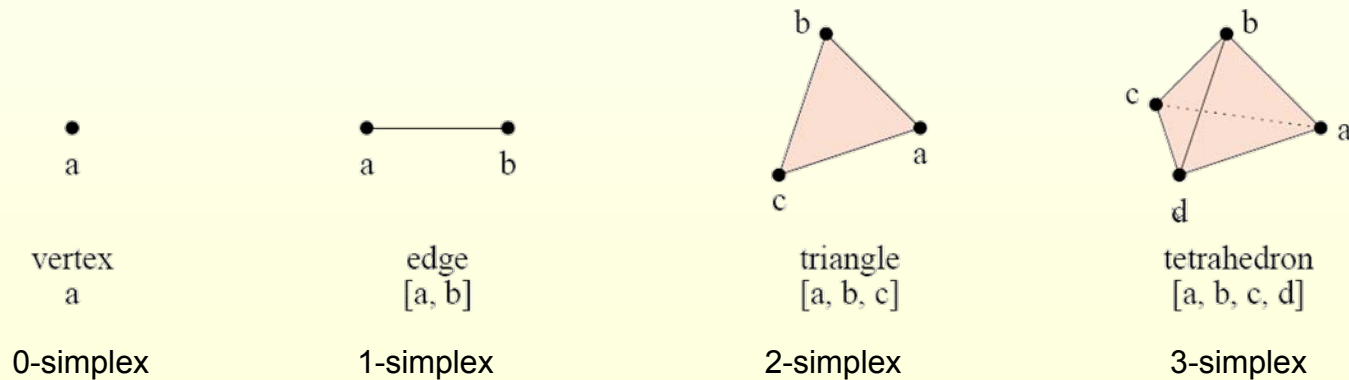


Simplicial Complexes



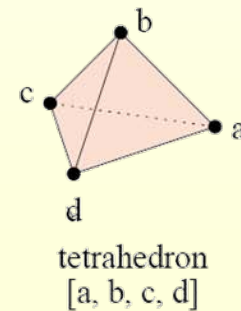
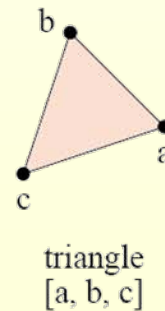
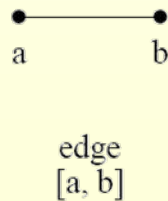
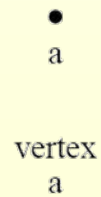
Simplices

- A k -simplex is the convex hull of $k + 1$ affinely independent points $S = \{v_0, v_1, \dots, v_k\}$. The points in S are the **vertices** of the simplex.
- A k -simplex is a k -dimensional subspace of \mathbb{R}^d , $\dim \sigma = k$.



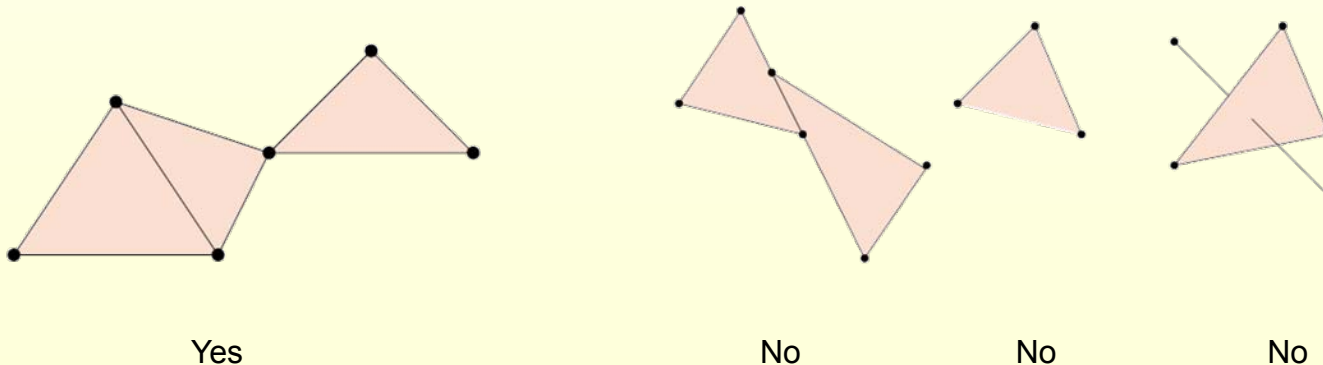
Faces

- σ : a k -simplex defined by $S = \{v_0, v_1, \dots, v_k\}$.
- τ defined by $T \subseteq S$ is a **face** of σ
- σ is its **coface**.
- $\sigma \geq \tau$ and $\tau \leq \sigma$.
- $\sigma \leq \sigma$ and $\sigma \geq \sigma$.

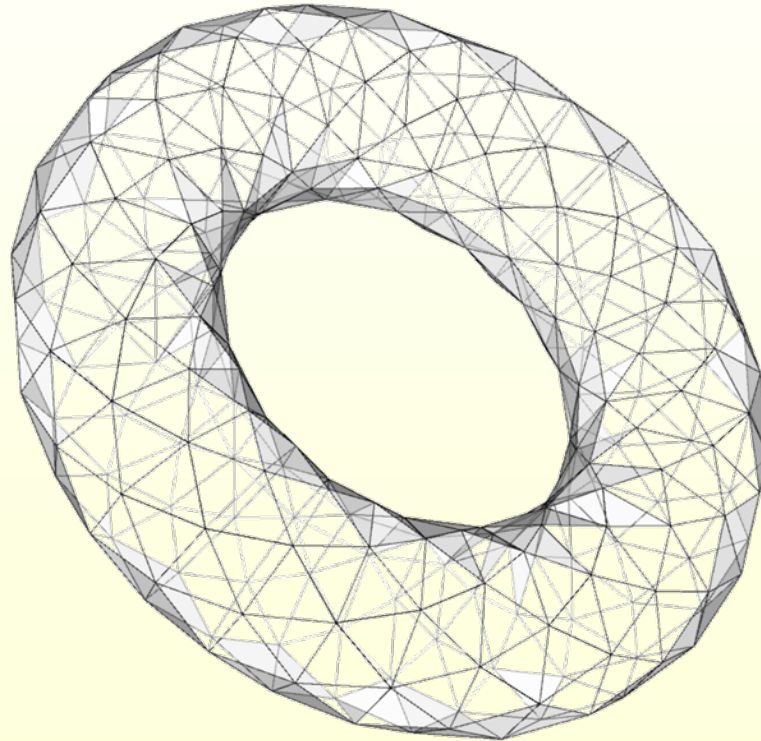


Simplicial Complex

- A **simplicial complex** K is a finite set of simplices such that
 1. $\sigma \in K, \tau \leq \sigma \Rightarrow \tau \in K$,
 2. $\sigma, \sigma' \in K \Rightarrow \sigma \cap \sigma' \leq \sigma, \sigma'$ or $\sigma \cap \sigma' = \emptyset$.
- The **dimension** of K is $\dim K = \max\{\dim \sigma \mid \sigma \in K\}$.
- The **vertices** of K are the zero-simplices in K .
- A simplex is **principal** if it has no proper coface in K .



A Triangle Mesh is a Simplicial Complex



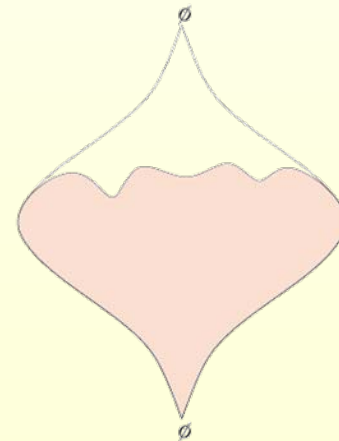
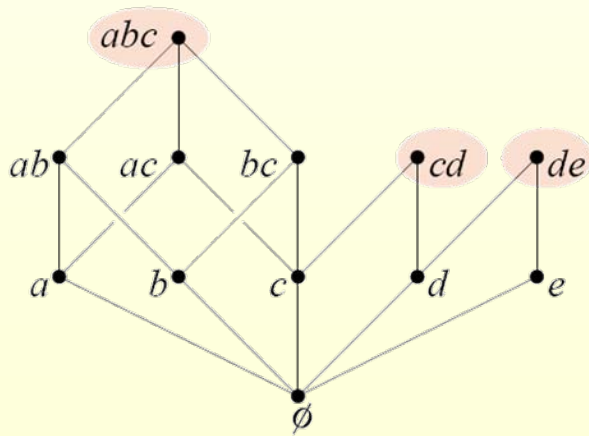
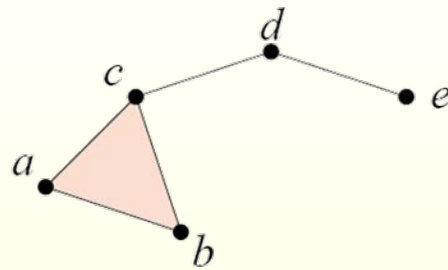
Abstract Simplicial Complex

- An **abstract simplicial complex** is a set K , together with a collection \mathcal{S} of subsets of K called **(abstract) simplices** such that:
 1. For all $v \in K$, $\{v\} \in \mathcal{S}$. We call the sets $\{v\}$ the **vertices** of K .
 2. If $\tau \subseteq \sigma \in \mathcal{S}$, then $\tau \in \mathcal{S}$.
- We call \mathcal{S} the complex.

Relationship

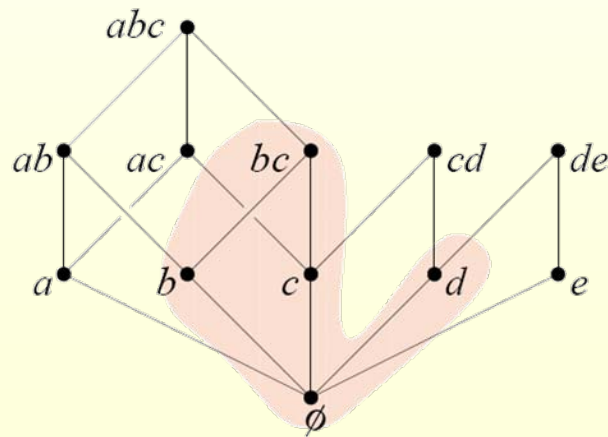
- Let K be a simplicial complex with vertices V and let \mathcal{S} be the collection of all subsets $\{v_0, v_1, \dots, v_k\}$ of V such that the vertices v_0, v_1, \dots, v_k span a simplex of K . Then, \mathcal{S} is the **vertex scheme** of K .
- K and \mathcal{S} form an abstract simplicial complex.
- Two abstract simplicial complexes are **isomorphic** if we can one from the other by renaming vertices.
- (Theorem) Every abstract complex \mathcal{S} is isomorphic to the vertex scheme of some simplicial complex K .
- We call K a **geometric realization** of \mathcal{S} .

An Example



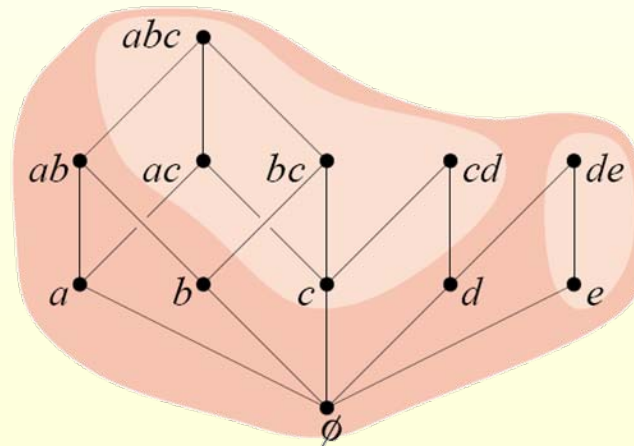
Subcomplex

- A **subcomplex** is a simplicial complex $L \subseteq K$. The smallest subcomplex containing a subset $L \subseteq K$ is its closure, $\text{Cl } L = \{\tau \in K \mid \tau \leq \sigma \in L\}$.
- Everything “below” is included.



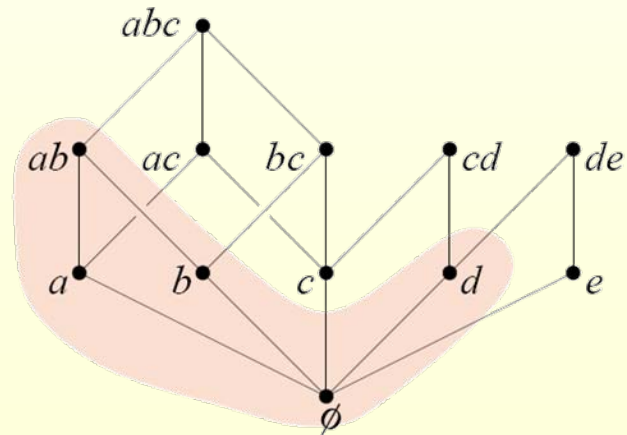
Star

- The **star of L** contains all of the cofaces of L ,
 $\text{St } L = \{\sigma \in K \mid \sigma \geq \tau \in L\}$.
- Everything “above” is included.
- Stars are analogs of neighborhoods (open).



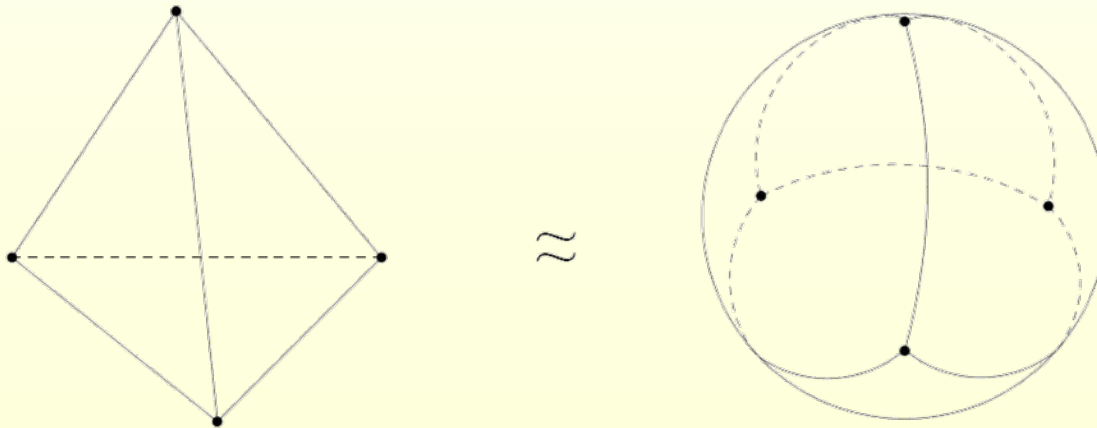
Link

- The **link of L** is the boundary of its star,
 $\text{Lk } L = \text{Cl St } L - \text{St}(\text{Cl } L - \{\emptyset\})$.



Triangulations

- The **underlying space** $|K|$ of a simplicial complex K is $|K| = \cup_{\sigma \in K} \sigma$.
- $|K|$ is a topological space.
- A **triangulation** of a topological space \mathbb{X} is a simplicial complex K such that $|K| \approx \mathbb{X}$.



We typically study shapes and spaces via triangulations of them

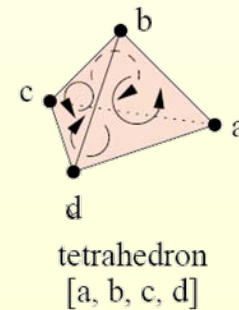
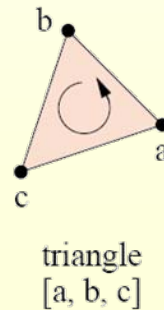
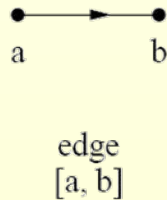
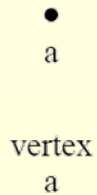
Orientability I

- An **orientation** of a k -simplex $\sigma \in K$, $\sigma = \{v_0, v_1, \dots, v_k\}$, $v_i \in K$ is an equivalence class of orderings of the vertices of σ , where

$$(v_0, v_1, \dots, v_k) \sim (v_{\tau(0)}, v_{\tau(1)}, \dots, v_{\tau(k)})$$

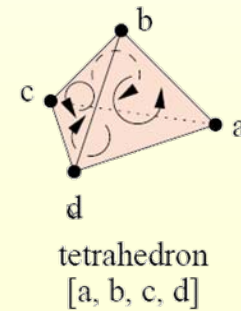
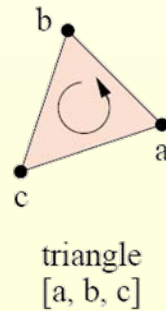
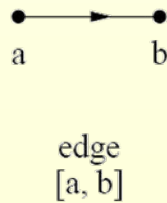
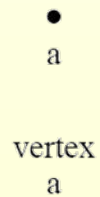
are equivalent orderings if the parity of the permutation τ is even.

- We denote an **oriented simplex**, a simplex with an equivalence class of orderings, by $[\sigma]$.



Orientability II

- Two k -simplices sharing a $(k - 1)$ -face σ are **consistently oriented** if they induce different orientations on σ .
- A triangulable d -manifold is **orientable** if all d -simplices can be oriented consistently.
- Otherwise, the d -manifold is **non-orientable**



Invariants

- A **(topological) invariant** is a map f that assigns the same object to spaces of the same topological type.
- $X \approx Y \implies f(X) = f(Y)$
- $f(X) \neq f(Y) \implies X \not\approx Y$ (contrapositive)
- $f(X) = f(Y) \implies$ nothing
- “coarser” differentiation

Invariants provide partial information about a space

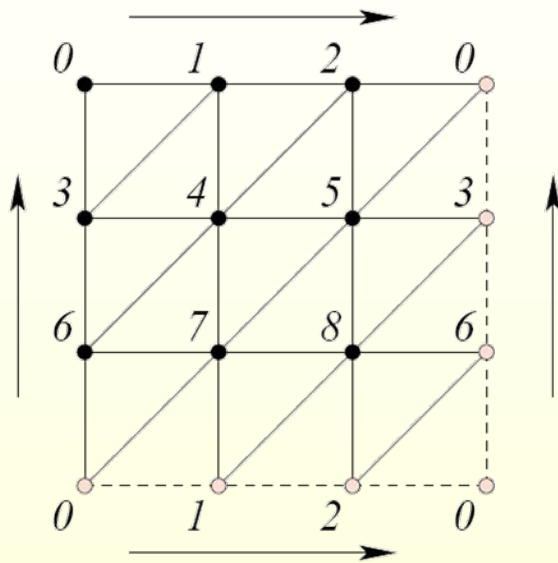
Euler Characteristic

- K a simplicial complex with s_k k -simplices.
- The Euler characteristic $\chi(K)$ is

$$\chi(K) = \sum_{i=0}^{\dim K} (-1)^i s_i = \sum_{\sigma \in K - \{\emptyset\}} (-1)^{\dim \sigma}.$$

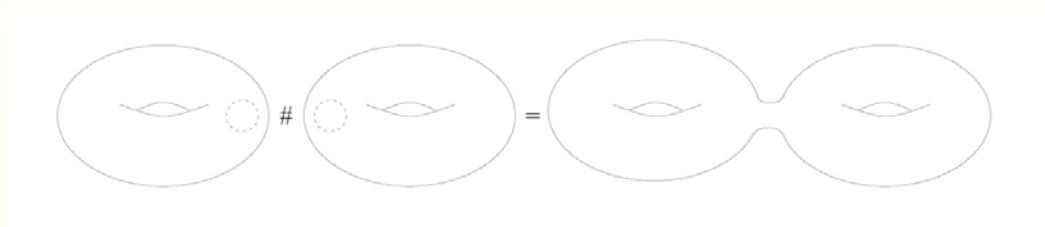
- $v - e + f = 1$ (Graph Theory)
- Invariant for $|K|$
- **Any** triangulation gives the same answer!
- Intrinsic property

Basic 2-Manifolds

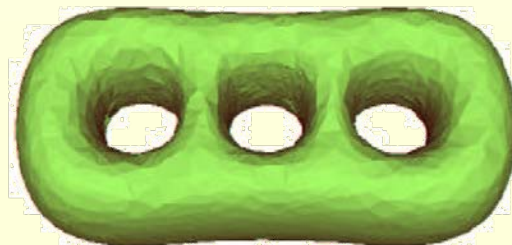


2-Manifold	χ
Sphere S^2	2
Torus T^2	0
Klein bottle K^2	0
Projective plane $\mathbb{R}P^2$	1

Euler and Connected Sums



- (Theorem) For compact surfaces M_1, M_2 ,
 $\chi(M_1 \# M_2) = \chi(M_1) + \chi(M_2) - 2$.
- $\chi(g\mathbb{T}^2) = 2 - 2g$
- $\chi(g\mathbb{RP}^2) = 2 - g$
- The connected sum of g tori is called a surface with **genus** g .



Compact 2-Manifolds

- $d = 2$: orientable



Torus



Double Torus



Triple Torus

add handles

• • •

- $d = 2$: non-orientable



Projective Plane \mathbb{P}^2

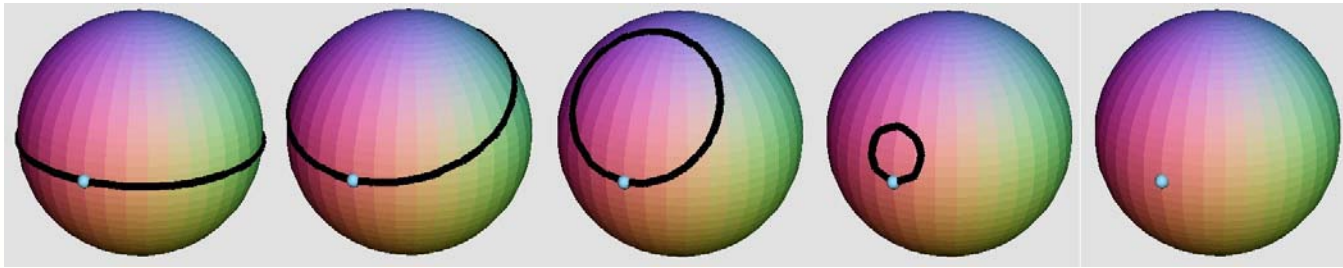


Klein Bottle

add cross-caps

• • •

Euler characteristic and orientability are two invariants providing a full classification of compact 2-manifolds

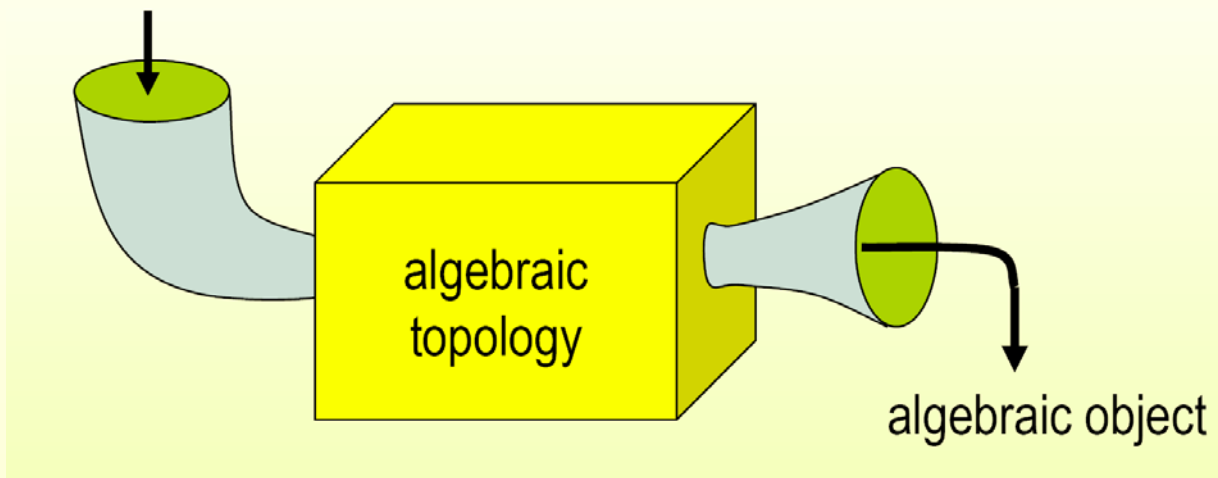


$d = 3$: Very hard

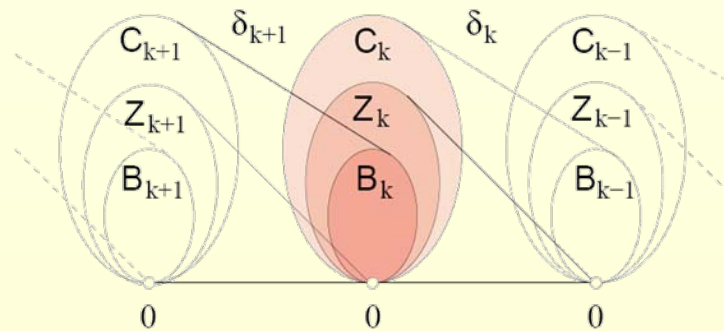
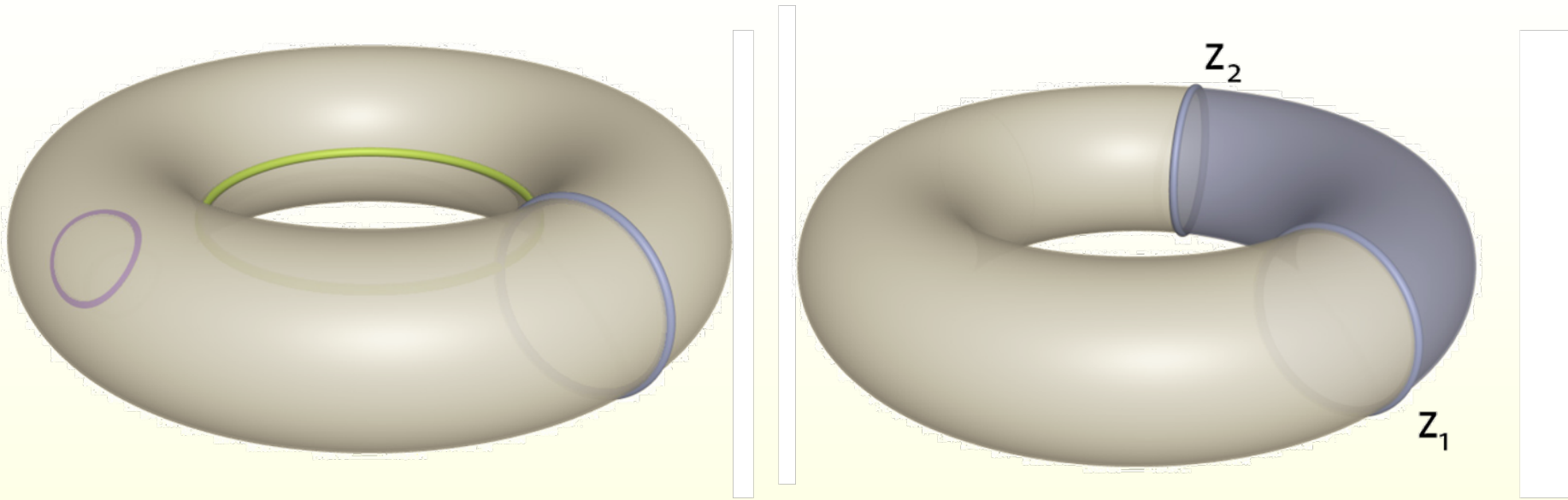
$d \geq 4$: *Undecidable* [Markov 1958]

[Grigori Perelman, 2003]

Homology

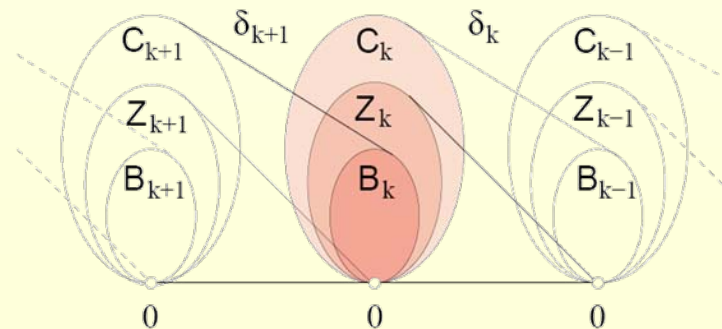


Homology



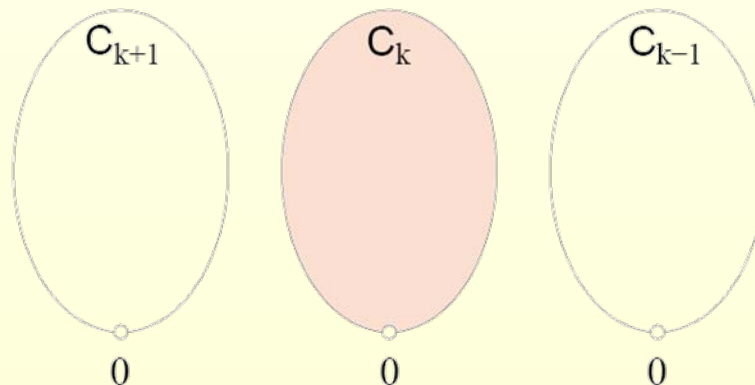
Why Homology?

- Algebraization of first layer of geometry in structures
- How cells of dimension n attach to cells of dimension $n - 1$
- Less transparent, more machinery
- Combinatorial
- Finite description
- Computable



Chain Group

- Simplicial complex K
- **k -chain**: $c = \sum_i n_i [\sigma_i]$, $n_i \in \mathbb{Z}$, $\sigma_i \in K$ (like a path)
- $[\sigma] = -[\tau]$ if $\sigma = \tau$ and σ and τ have different orientations.
- The **k th chain group \mathbf{C}_k** of K is the free abelian group on its set of oriented k -simplices
- $\text{rank } \mathbf{C}_k = ?$



Boundary Operator

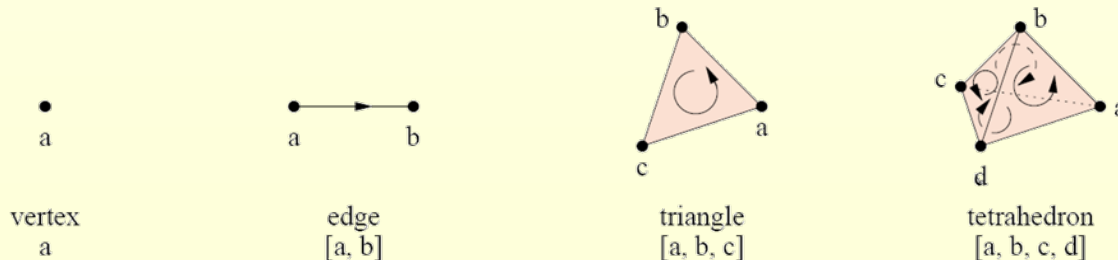
- The boundary operator $\partial_k : \mathbf{C}_k \rightarrow \mathbf{C}_{k-1}$ is a homomorphism defined linearly on a chain c by its action on any simplex

$$\sigma = [v_0, v_1, \dots, v_k] \in c,$$

$$\partial_k \sigma = \sum_i (-1)^i [v_0, v_1, \dots, \hat{v}_i, \dots, v_k],$$

where \hat{v}_i indicates that v_i is deleted from the sequence.

- $\partial_1 [a, b] = b - a.$
- $\partial_2 [a, b, c] = [b, c] - [a, c] + [a, b] = [b, c] + [c, a] + [a, b].$
- $\partial_3 [a, b, c, d] = [b, c, d] - [a, c, d] + [a, b, d] - [a, b, c].$
- $\partial_1 \partial_2 [a, b, c] = [c] - [b] - [c] + [a] + [b] - [a] = 0.$



Boundary Theorem

- (Theorem) $\partial_{k-1}\partial_k = 0$, for all k .

- Proof:

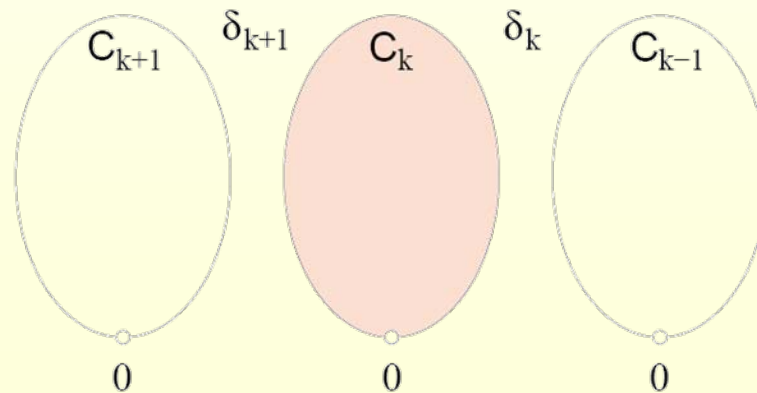
$$\begin{aligned}\partial_{k-1}\partial_k[v_0, v_1, \dots, v_k] &= \\ &= \partial_{k-1} \sum_i (-1)^i [v_0, v_1, \dots, \hat{v}_i, \dots, v_k] \\ &= \sum_{j < i} (-1)^i (-1)^j [v_0, \dots, \hat{v}_j, \dots, \hat{v}_i, \dots, v_k] \\ &\quad + \sum_{j > i} (-1)^i (-1)^{j-1} [v_0, \dots, \hat{v}_i, \dots, \hat{v}_j, \dots, v_k] \\ &= 0,\end{aligned}$$

as switching i and j in the second sum negates the first sum.

Chain Complex

- The boundary operator connects the chain groups into a **chain complex** C_* :

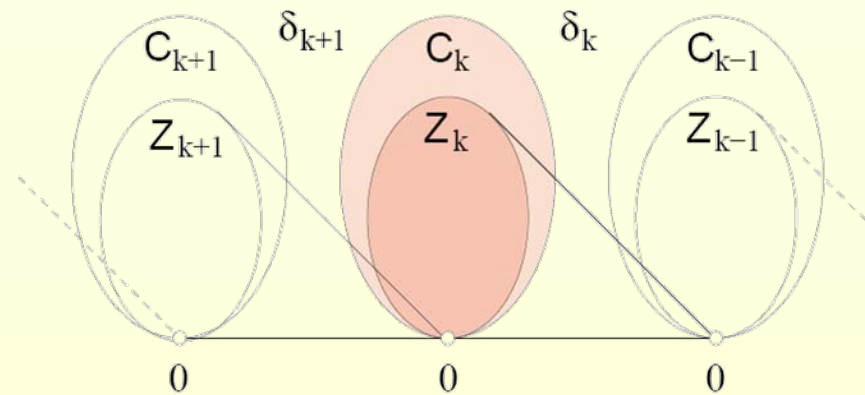
$$\dots \rightarrow C_{k+1} \xrightarrow{\partial_{k+1}} C_k \xrightarrow{\partial_k} C_{k-1} \rightarrow \dots$$



Cycle Group

- Let c be a k -chain
- If it has no boundary, it is a k -cycle (zycle?)
- $\partial_k c = \emptyset$, so $c \in \ker \partial_k$
- The k th cycle group is

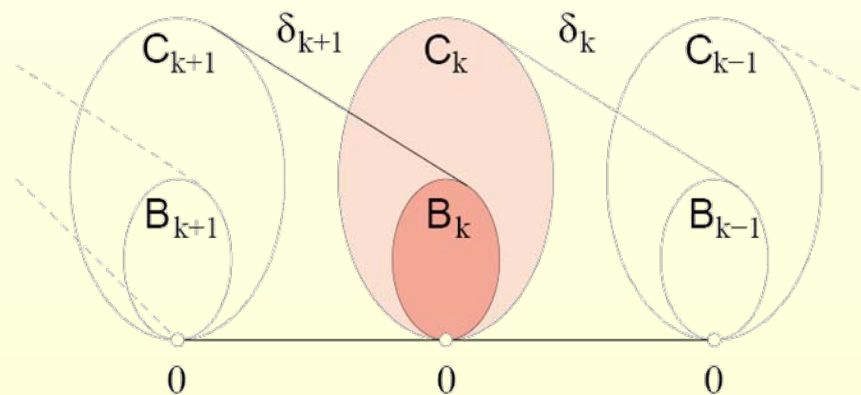
$$Z_k = \ker \partial_k = \{c \in C_k \mid \partial_k c = \emptyset\}.$$



Boundary Group

- Let b be a k -chain
- If b is a boundary of something, it is a k -boundary.
- The k th boundary group is

$$\mathbf{B}_k = \text{im } \partial_{k+1} = \{c \in \mathbf{C}_k \mid \exists d \in \mathbf{C}_{k+1} : c = \partial_{k+1}d\}.$$



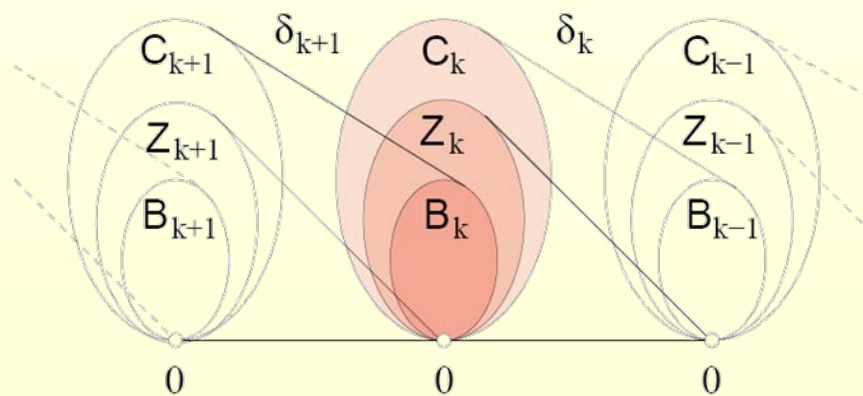
Nesting Property

- Let b be a k -boundary.
- Then, $\exists c \in \mathbf{C}_{k+1}$, such that $b = \partial_{k+1}c$.
- What is the boundary of b ?

$$\partial_k b = \partial_k \partial_{k+1} c = \emptyset,$$

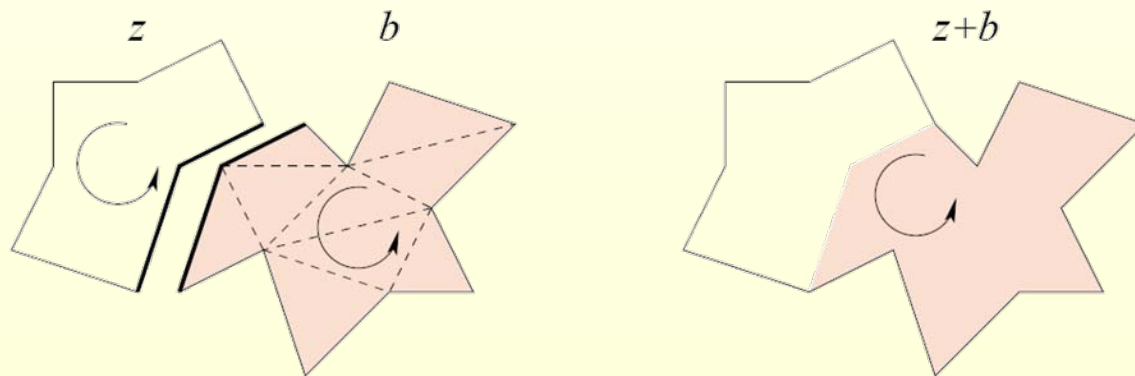
by the boundary theorem.

- That is, every boundary is a cycle!
- $\mathbf{B}_k \subseteq \mathbf{Z}_k \subseteq \mathbf{C}_k$



Cycle Equivalence

- z is a k -cycle
- b is a k -boundary
- We would like to have $z + b$ be equivalent to z
- That is, if $z_1 - z_2 = b$ where b is a boundary, then $z_1 \sim z_2$
- Any boundary would do!

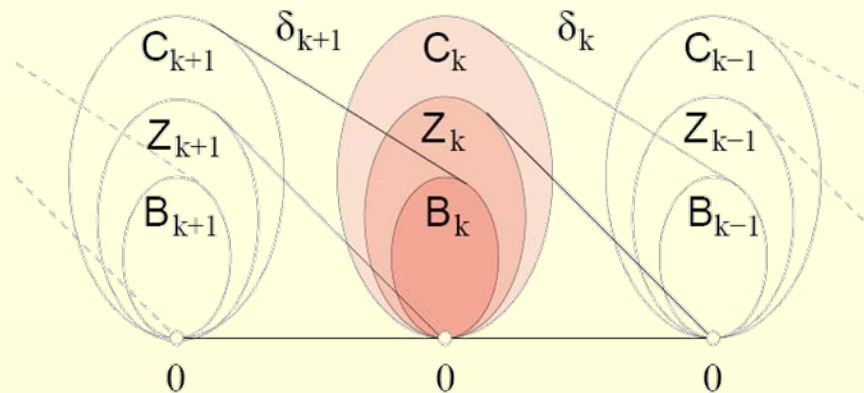


Simplicial Homology

- The k th homology group is

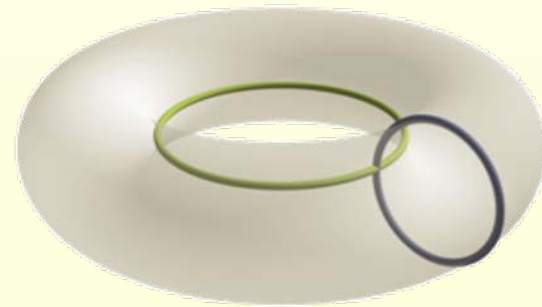
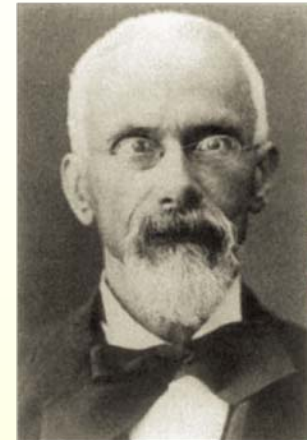
$$H_k = Z_k / B_k = \ker \partial_k / \text{im } \partial_{k+1}.$$

- If $z_1 = z_2 + B_k$, $z_1, z_2 \in Z_k$, we say z_1 and z_2 are **homologous**
- $z_1 \sim z_2$.



Z_2 Homology

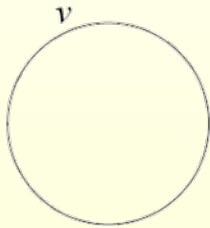
- H_k is a vector space
- k th Betti number $\beta_k = \text{rank } H_k$
 $= \text{rank } Z_k - \text{rank } B_k$
- Enrico Betti (1823 – 1892)
- Geometric interpretation in R^3
 - β_0 is number of components
 - β_1 is rank of a basis for tunnels
 - β_2 is number of voids



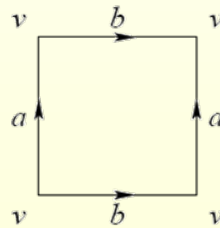
1, 2, 1

Homology of 2-Manifolds

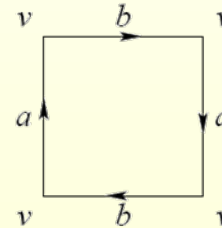
2-manifold	H_0	H_1	H_2
sphere	\mathbb{Z}	$\{0\}$	\mathbb{Z}
torus	\mathbb{Z}	$\mathbb{Z} \times \mathbb{Z}$	\mathbb{Z}
projective plane	\mathbb{Z}	\mathbb{Z}_2	$\{0\}$
Klein bottle	\mathbb{Z}	$\mathbb{Z} \times \mathbb{Z}_2$	$\{0\}$



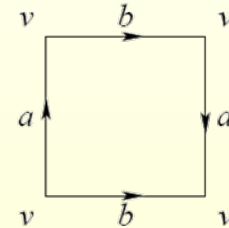
(a) Sphere



(b) Torus



(c) Projective
plane



(d) Klein bottle

Again, homology is independent of the triangulation

Euler Revisited

- Let K be a simplicial complex and $s_i = |\{\sigma \in K \mid \dim \sigma = i\}|$. The Euler characteristic $\chi(K)$ is

$$\chi(K) = \sum_{i=0}^{\dim K} (-1)^i s_i = \sum_{\sigma \in K - \{\emptyset\}} (-1)^{\dim \sigma}.$$

- We have new language!
- Let \mathbf{C}_* be the chain complex on K
- $\text{rank}(\mathbf{C}_i) = |\{\sigma \in K \mid \dim \sigma = i\}|$
- $\chi(K) = \chi(\mathbf{C}_*) = \sum_i (-1)^i \text{rank}(\mathbf{C}_i)$.

Euler-Poincaré

- Homology functors H_*
- $H_*(C_*)$ is a chain complex:

$$\dots \rightarrow H_{k+1} \xrightarrow{\partial_{k+1}} H_k \xrightarrow{\partial_k} H_{k-1} \rightarrow \dots$$

- What is its Euler characteristic?
- (Theorem) $\chi(K) = \chi(C_*) = \chi(H_*(C_*))$.

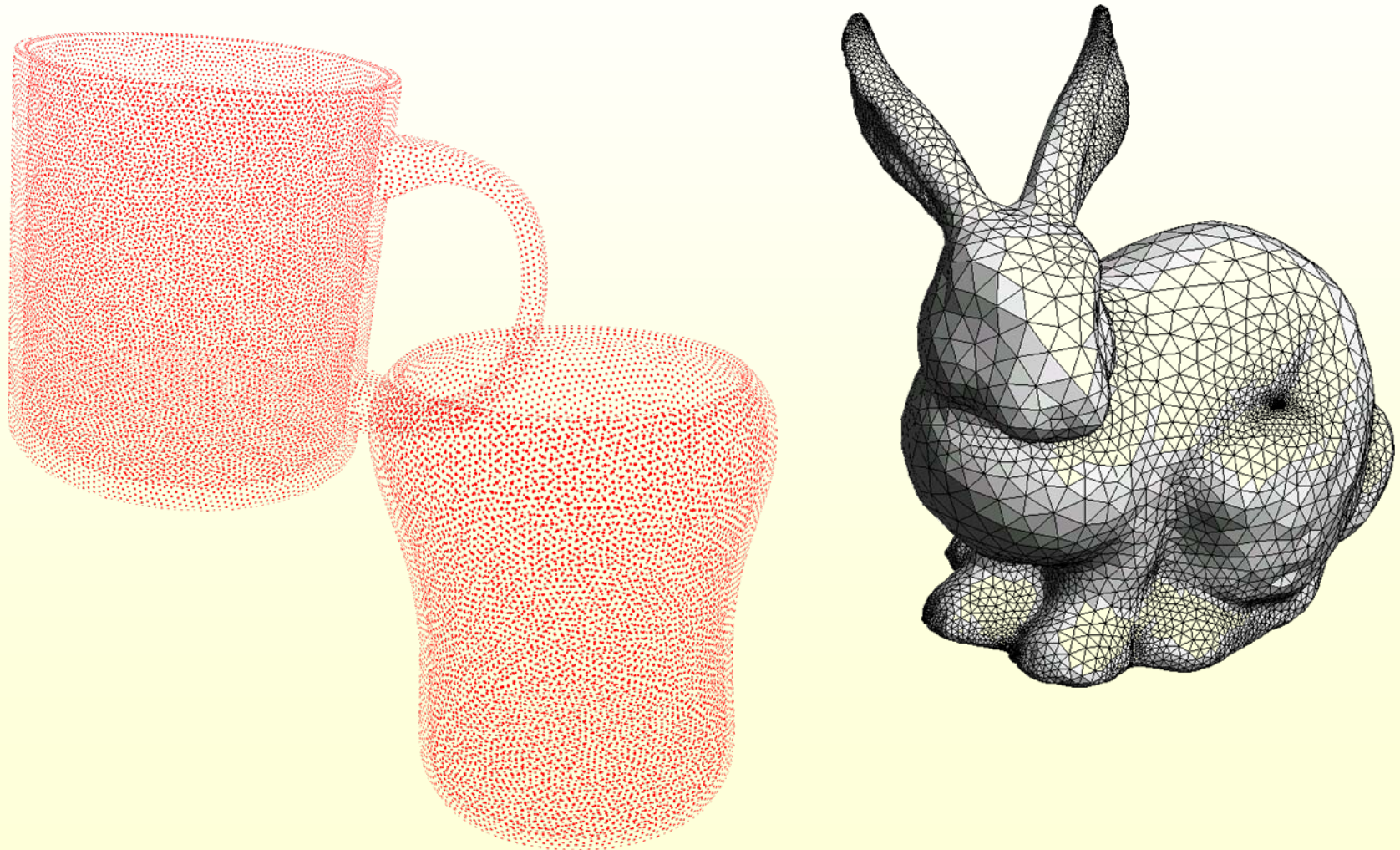
$$\sum_i (-1)^i s_i = \sum_i (-1)^i \text{rank}(H_i) = \sum_i (-1)^i \beta_i$$

- Sphere: $2 = 1 - 0 + 1$
- Torus: $0 = 1 - 2 + 1$

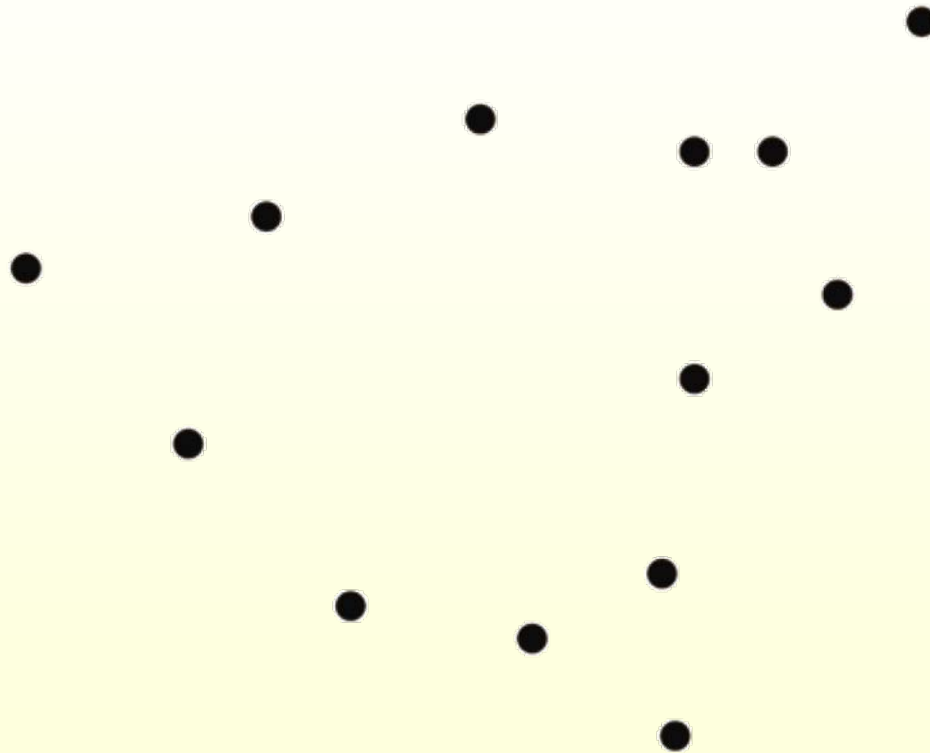
Point Clouds and the Complex Zoo



What Does This All Mean for Point Clouds or Meshes?

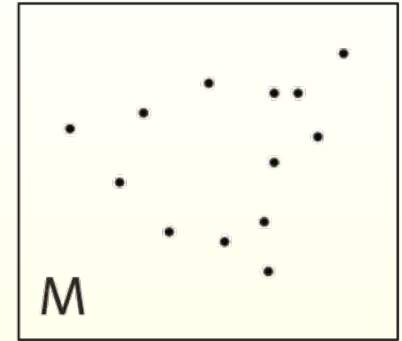


Topology of Points



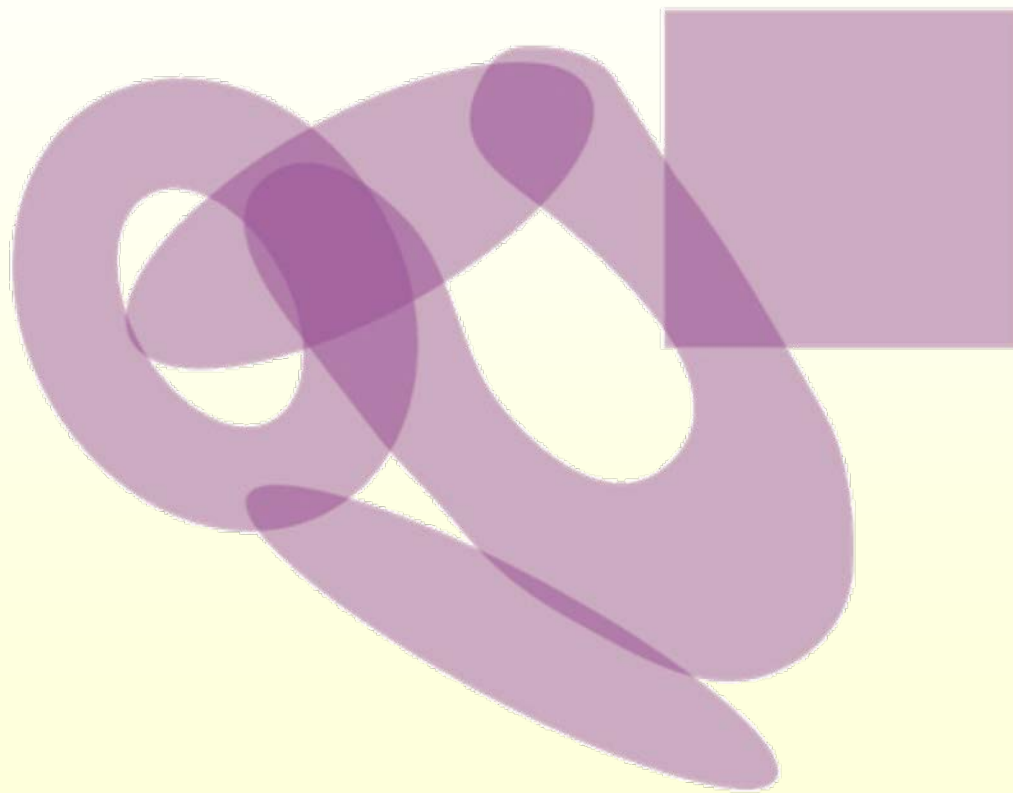
A Hidden Space X

- Topological space X
- Underlying space
- Given: set of sample points M from X



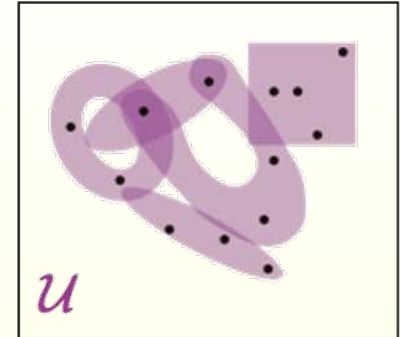
- Question: How can we recover the topology of X from M ?
- Problem: M has no interesting topology.

Open Cover

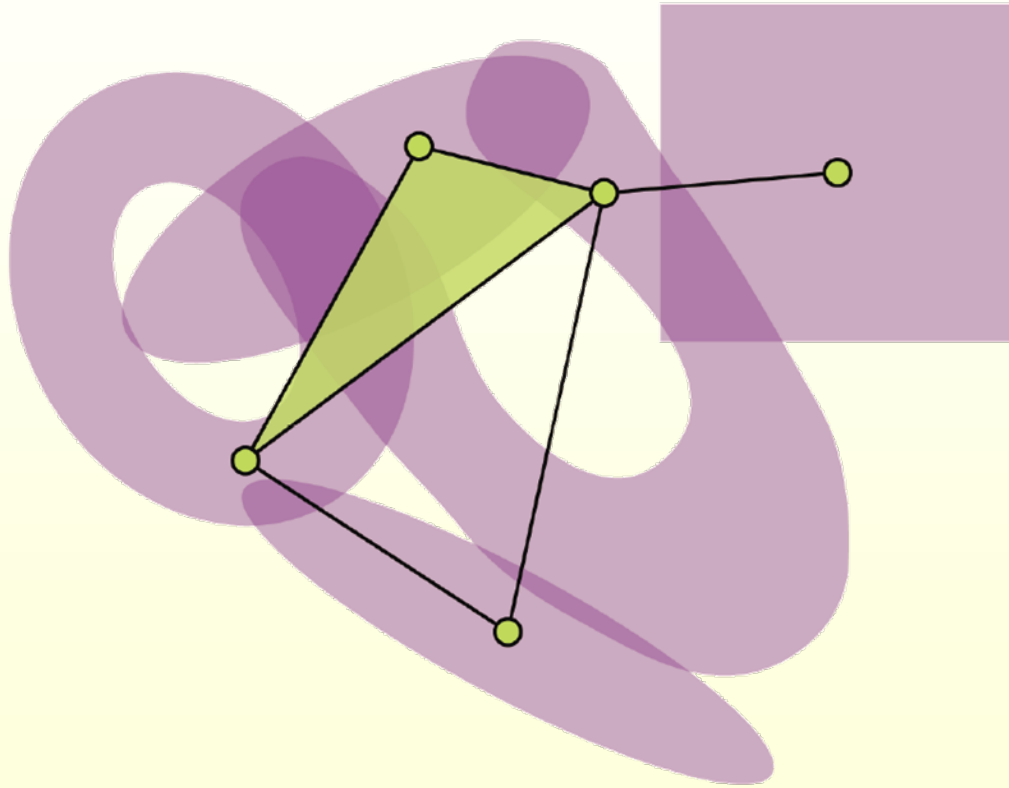


Formally

- Cover $\mathcal{U} = \{U_i\}_{i \in I}$
 - U_i , open
 - $M \subseteq \bigcup_{i \in I} U_i$
- Idea: The cover approximates the underlying space \mathbb{X}
- Question': What is the topology of \mathcal{U} ?
- Problem: \mathcal{U} is an infinite point set



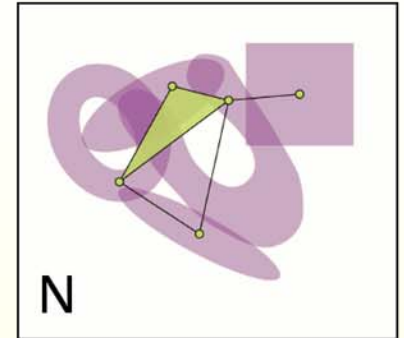
The Nerve of the Cover



An **abstract** simplicial complex

Formally

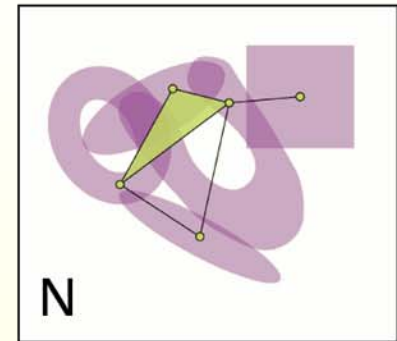
- \mathbb{X} : topological space
- $\mathcal{U} = \bigcup_{i \in I} U_i$: open cover of \mathbb{X}
- The **nerve** N of \mathcal{U} is
 - $\emptyset \in N$
 - If $\bigcap_{j \in J} U_j \neq \emptyset$ for $J \subseteq I$, then $J \in N$
- Dual structure
- (Abstract) Simplicial complex



The Nerve Lemma

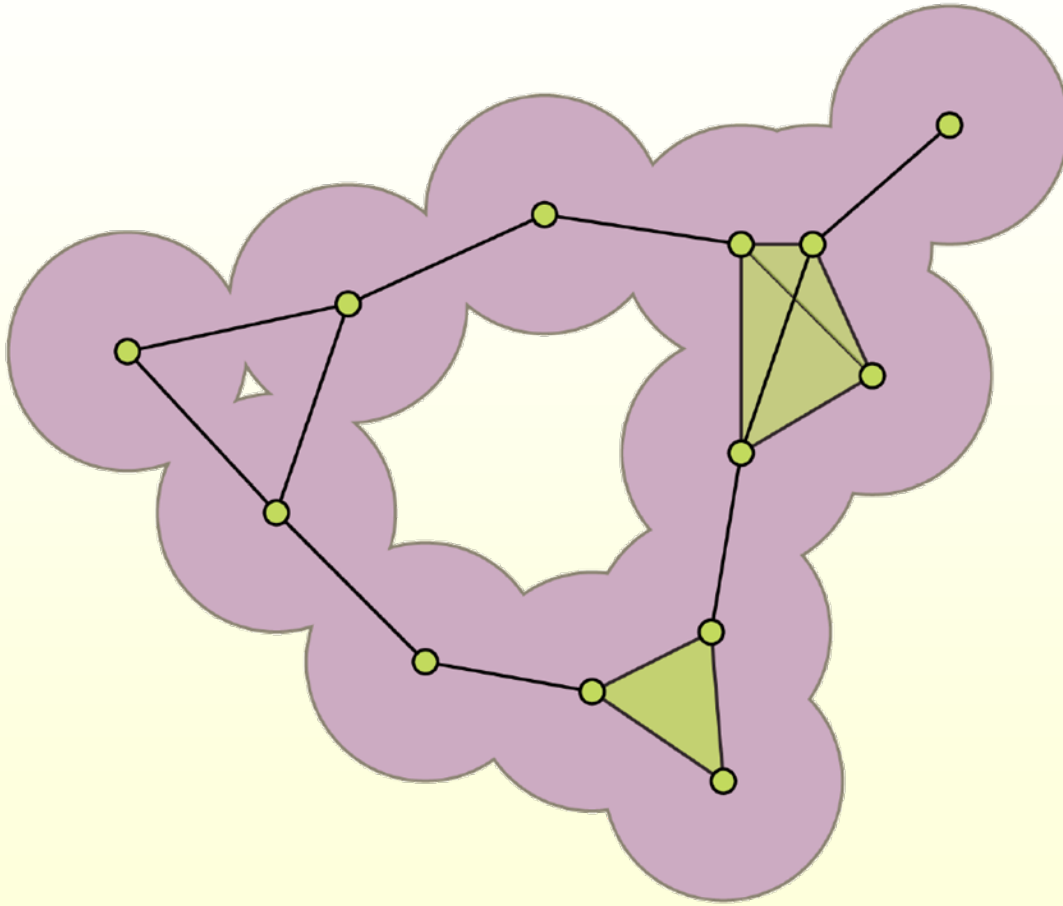
- **(Lemma [Leray])**

If sets in the cover are contractible, and their finite ~~unions~~ intersections are contractible, then $N \simeq \mathcal{U}$.



- *The cover should not introduce or eliminate topological structure*
- Idea: Use “nice” sets for covering
 - contractible
 - convex
- Dual (abstract) simplicial complex will be our representation

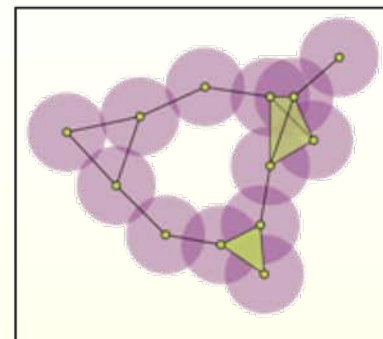
The Čech Complex



Ball radius is a parameter ϵ

Formally

- Set: Ball of radius ε
 $B_\varepsilon(x) = \{y \mid d(x, y) < \varepsilon\}$
- Cover: B_ε at every point in M

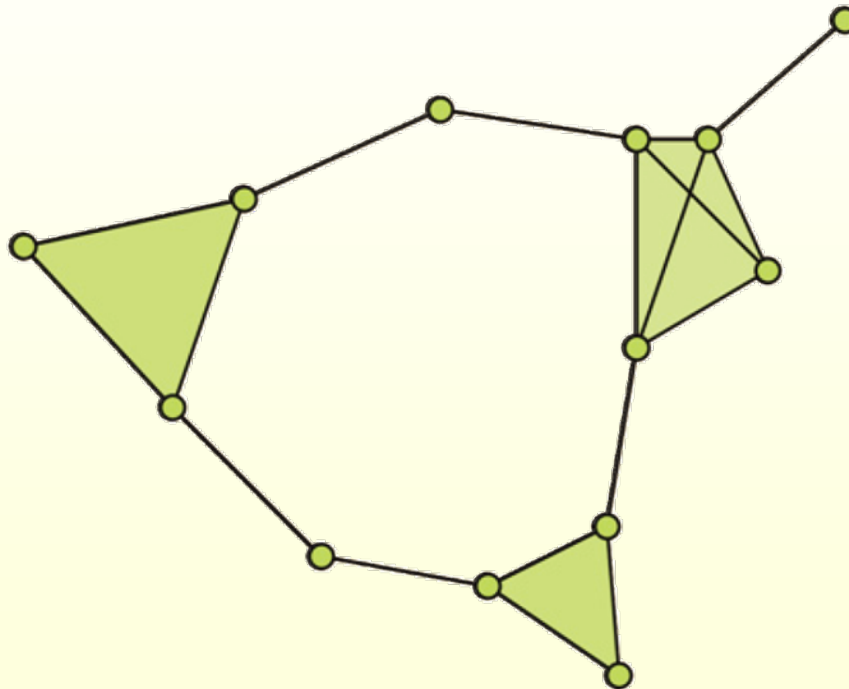


- **Cech complex** is nerve of the union of ε -balls

$$C_\varepsilon(M) = \left\{ \text{conv } T \mid T \subseteq M, \bigcup_{m \in T} B_\varepsilon(m) \neq \emptyset \right\}$$

- Cover satisfies Nerve Lemma
- Eduard Cech (1893 – 1960)

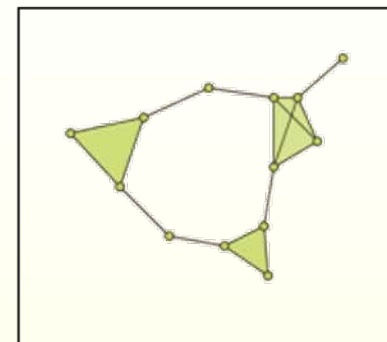
Vietoris-Rips Complex



Distance is a parameter ϵ

Formally

1. Construct ε -graph
2. Expand by add a simplex whenever all its faces are in the complex
 - Note: We expand by dimension

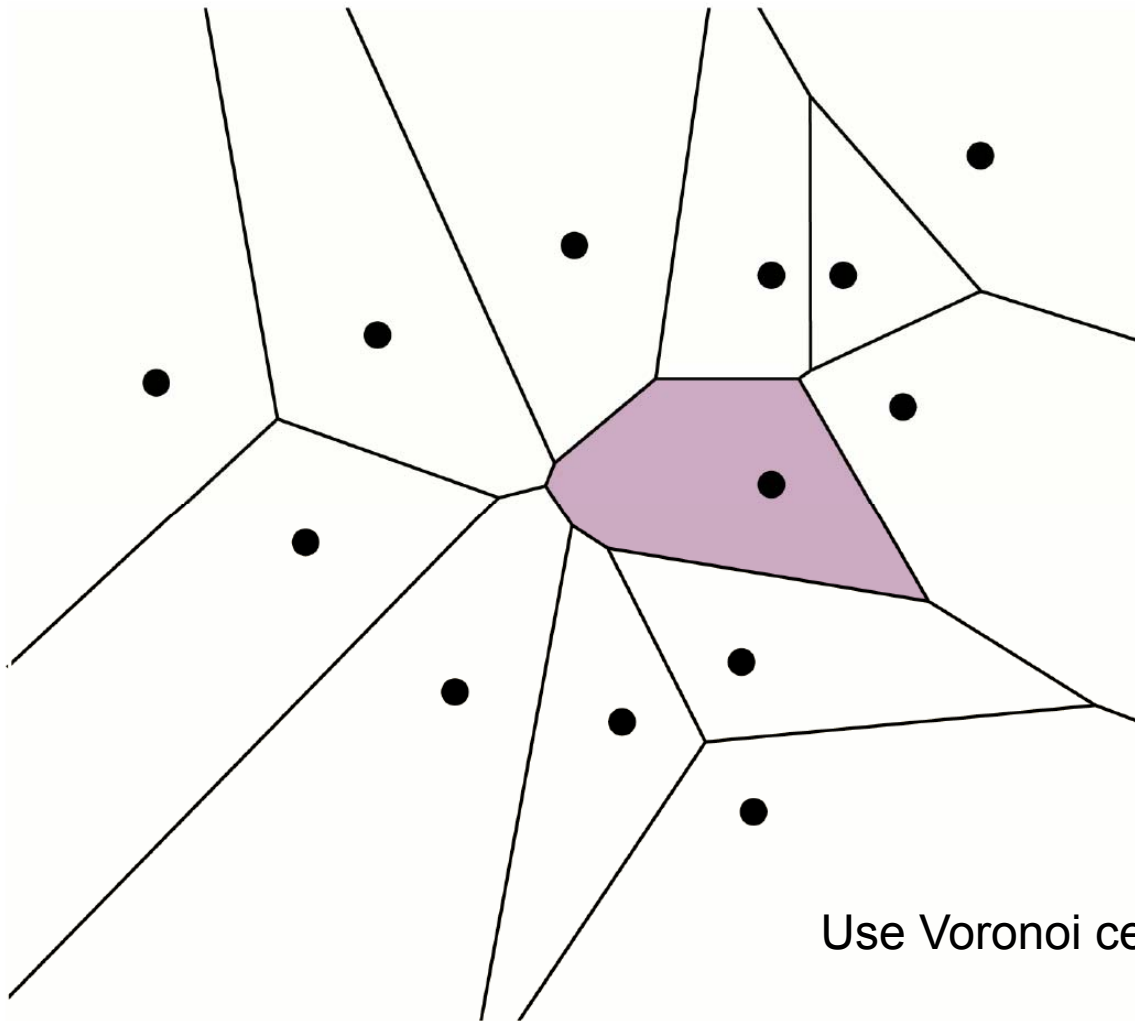


$$V_\epsilon(M) = \{\text{conv } T \mid T \subseteq M, d(x, y) < \epsilon, \forall x, y \in T\}$$

- $V_{2\epsilon}(M) \supseteq C_\epsilon(M)$
- Not homotopic to union of balls
- Leopold Vietoris (1891 – 2002)
- Eliyahu Rips (1948 –)

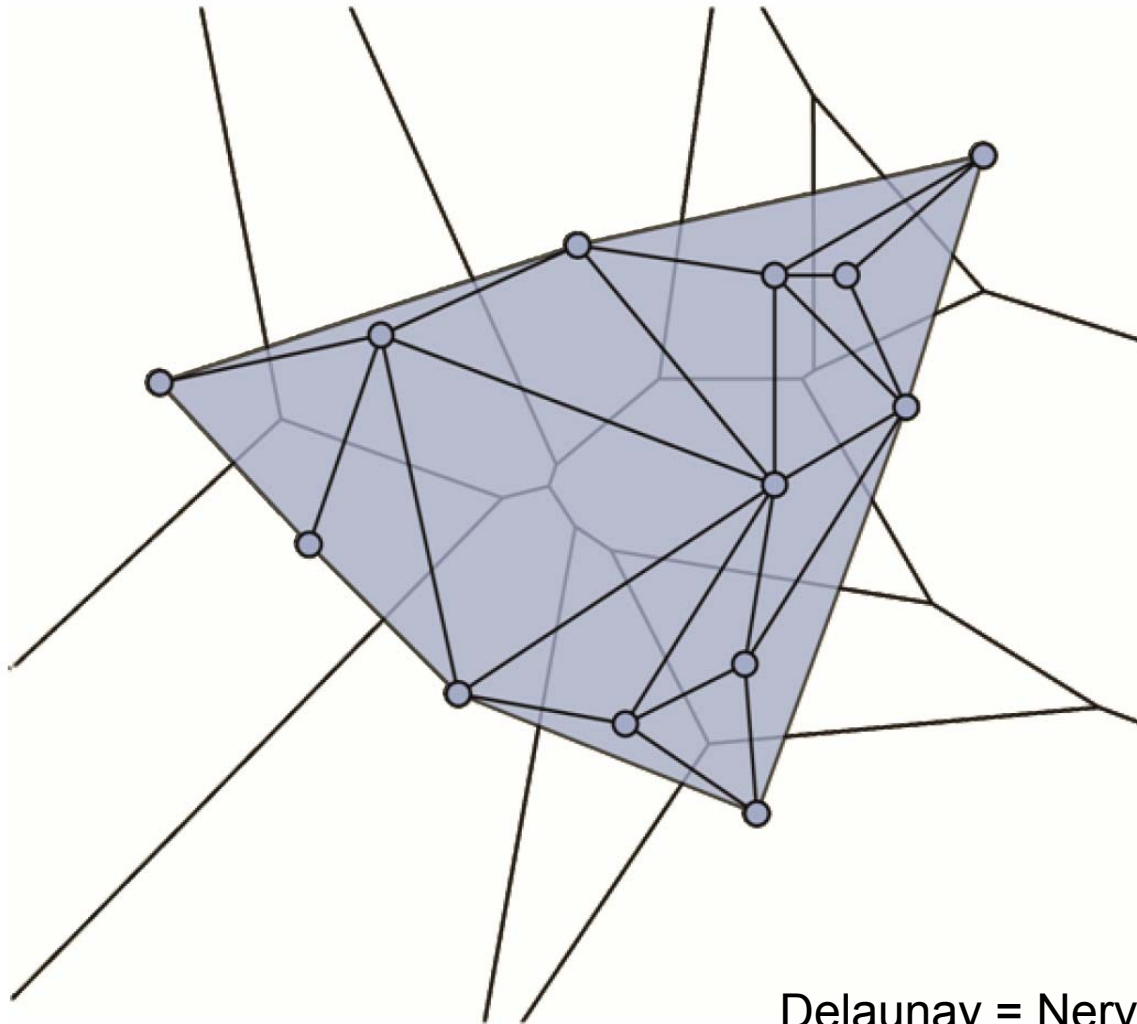


Geometric Complexes: Voronoi



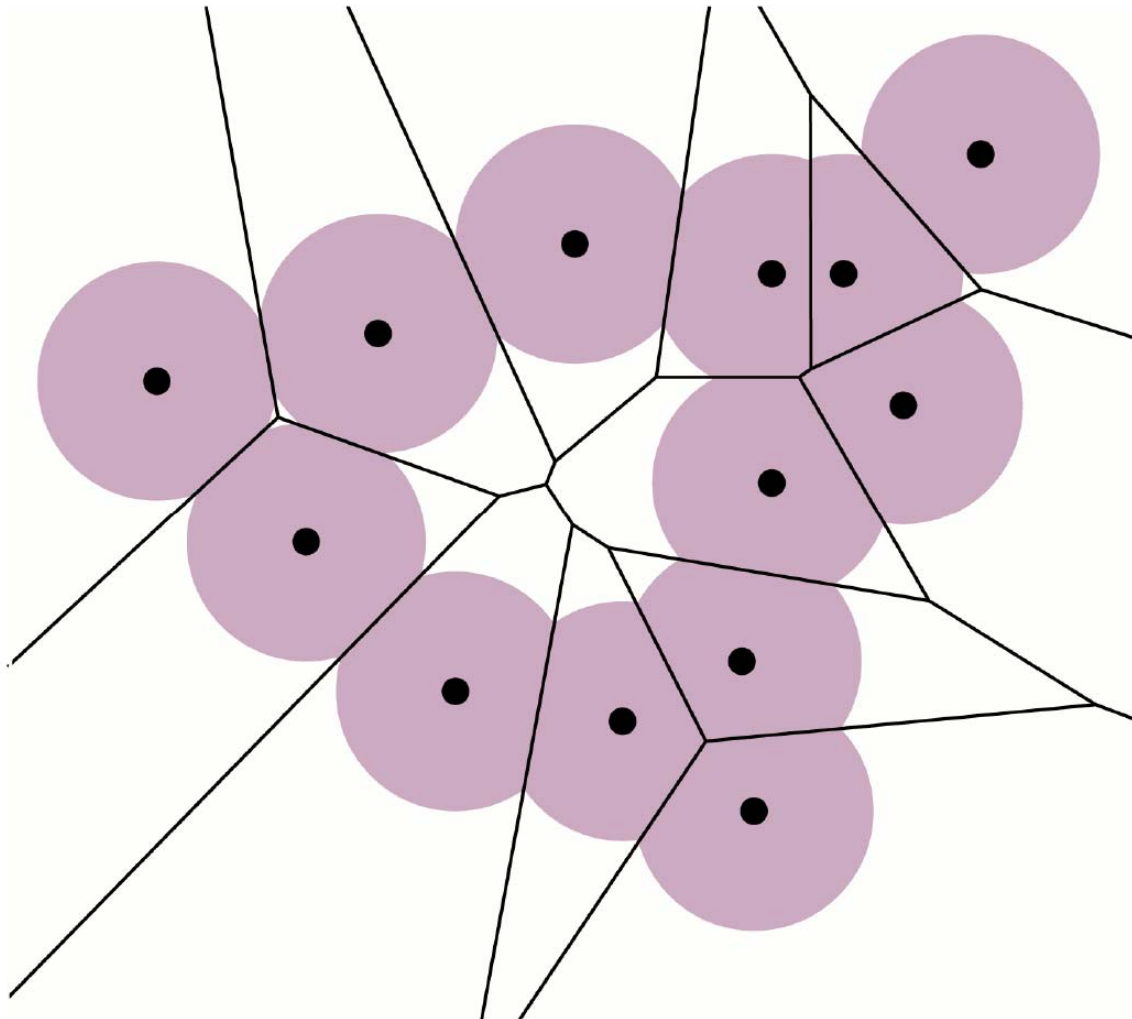
Use Voronoi cells as open cover

Dual Complex: Delaunay



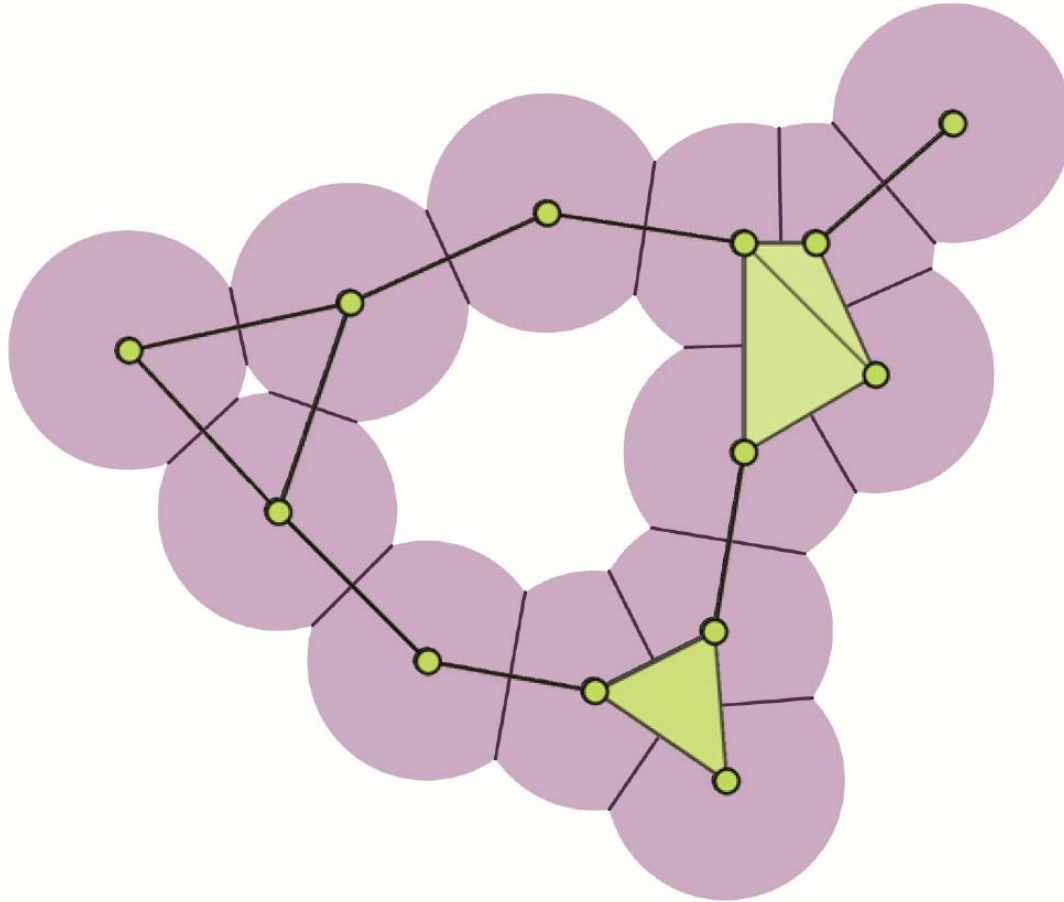
Delaunay = Nerve of Voronoi

Restricted Voronoi



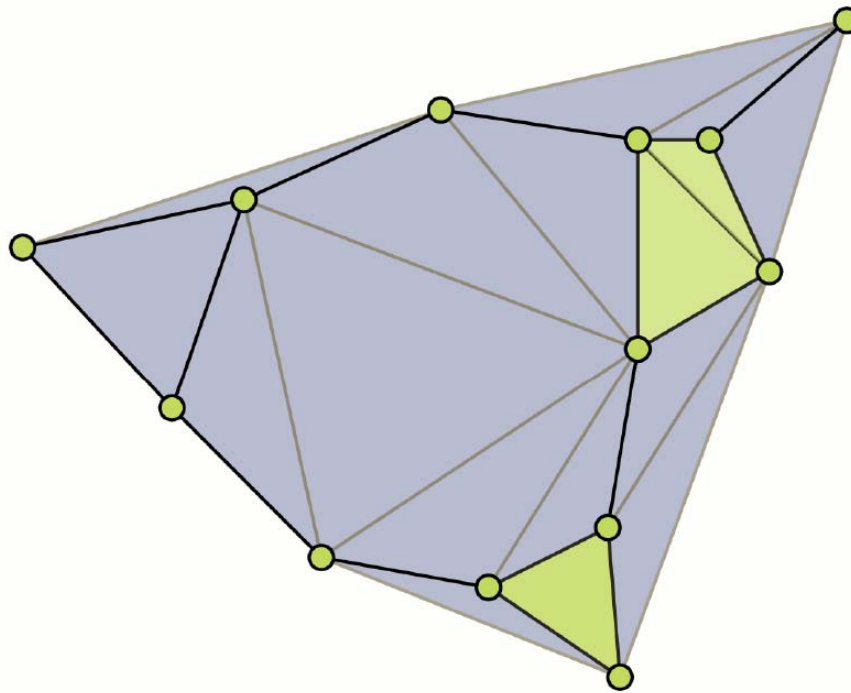
Ball radius is a parameter ϵ

Alpha Complex



Ball radius is a parameter ϵ

Subcomplex of Delaunay



Formally

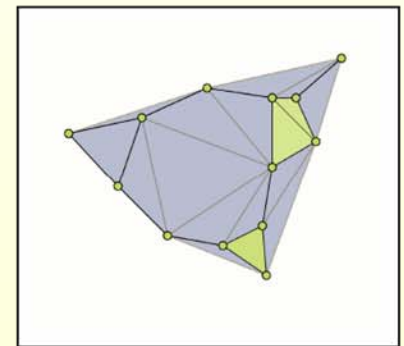
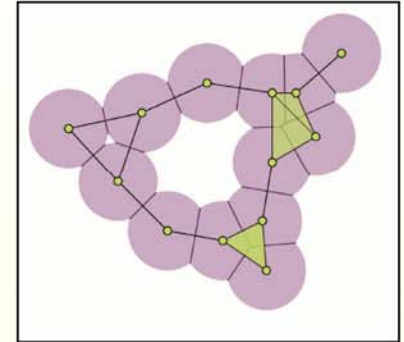
- Alpha cell: $A_\varepsilon(p) = B_\varepsilon(p) \cap V(p)$
- Alpha shape: union of alpha cells
- Alpha complex: nerve of alpha shape

$$A_\varepsilon(M) = \left\{ \text{conv} T \mid T \subseteq M, \bigcap_{p \in T} A_\varepsilon(p) \neq \emptyset \right\}$$

- Let D be the Delaunay triangulation
 - $A_0 = \emptyset$
 - $A_\varepsilon \subseteq D$
 - $A_\infty = D$

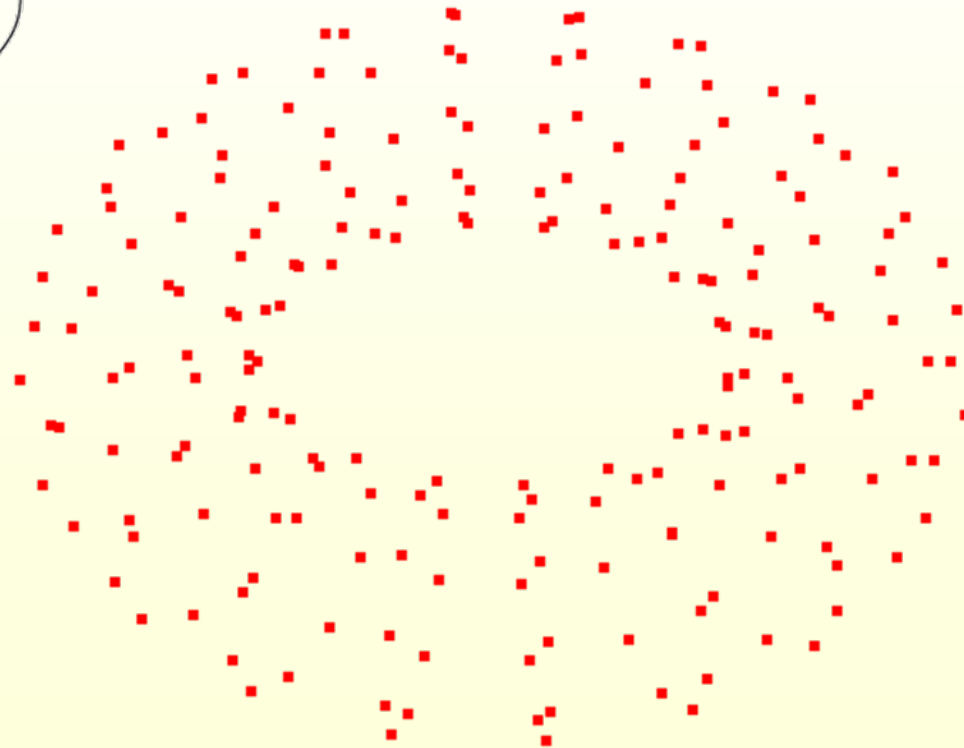
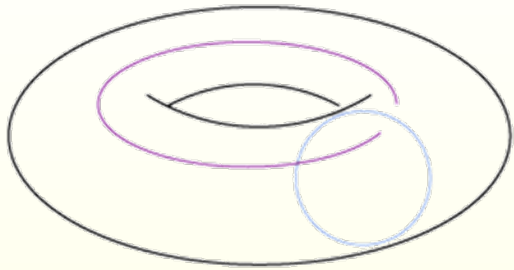
- $A_\varepsilon \simeq C_\varepsilon$

- [Edelsbrunner, Kirkpatrick, and Seidel '83], et al.



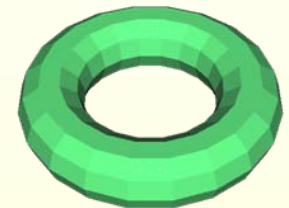
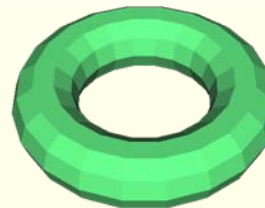
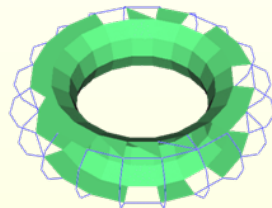
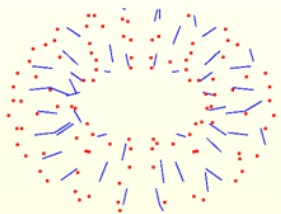
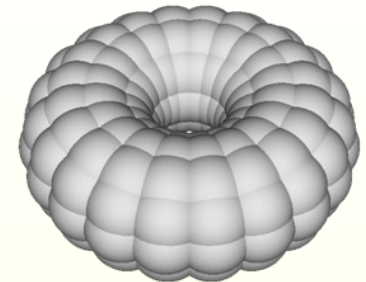
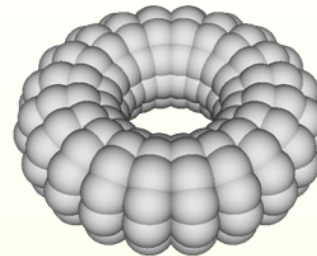
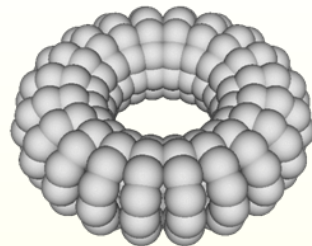
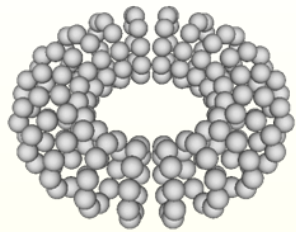
Persistent Homology

Detecting a Torus



PCD

Question of Scale: A Filtration



$$\beta_0 = 150$$

$$\beta_1 = 0$$

$$\beta_0 = 1$$

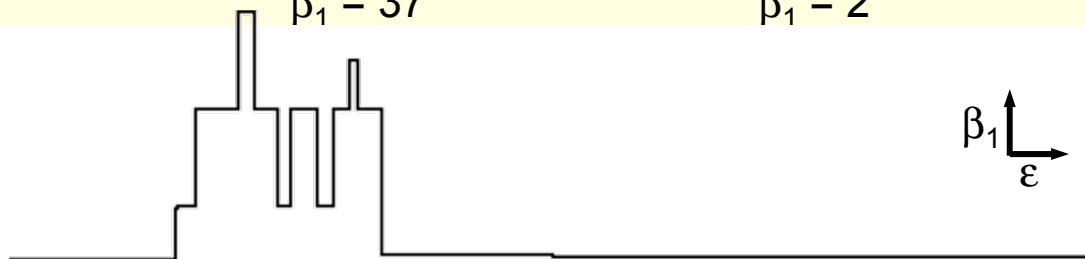
$$\beta_1 = 37$$

$$\beta_0 = 1$$

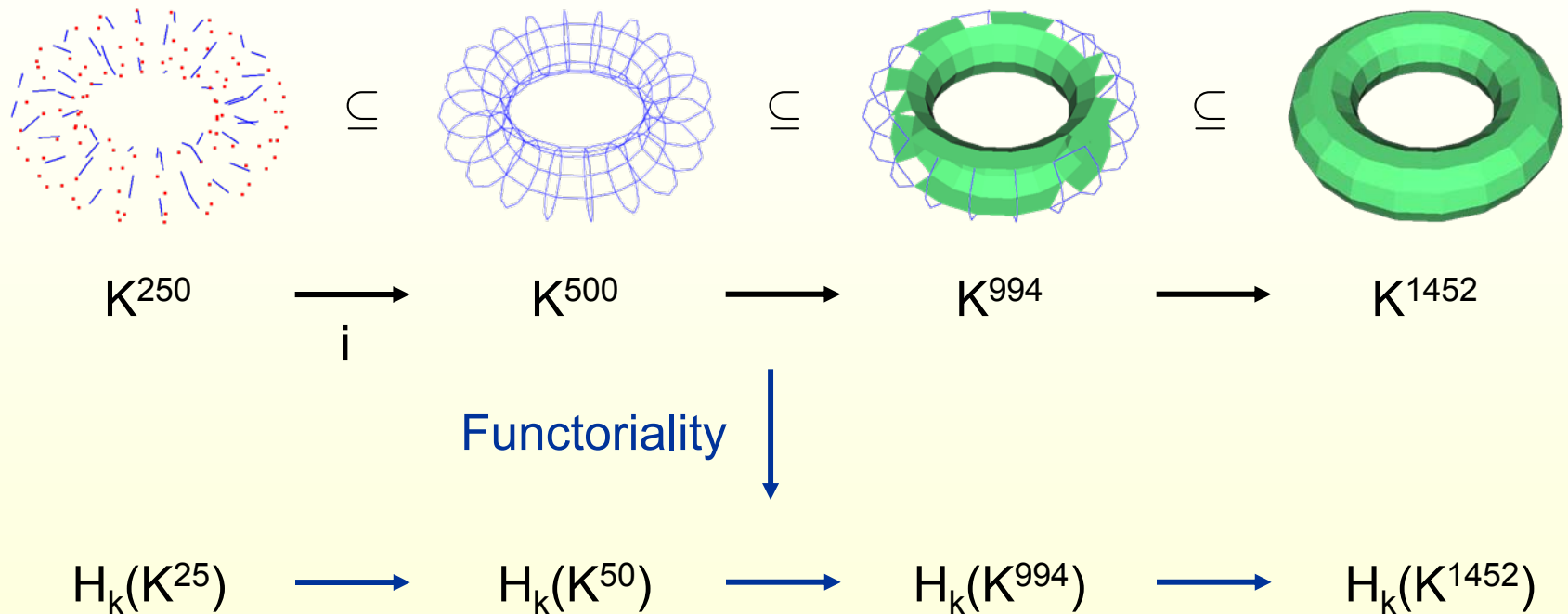
$$\beta_1 = 2$$

$$\beta_0 = 1$$

$$\beta_1 = 1$$

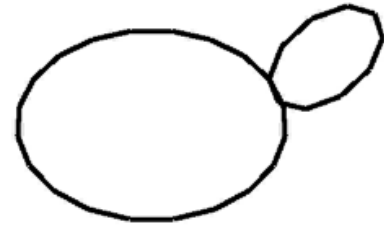
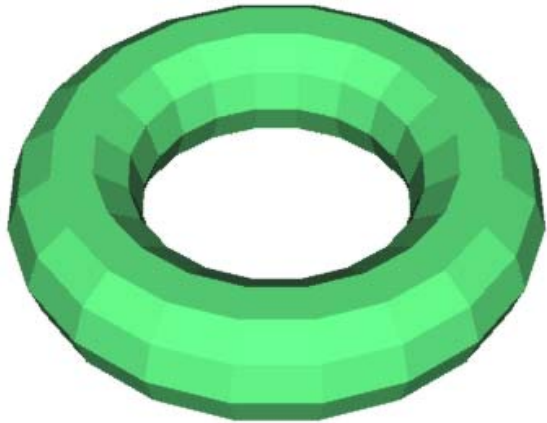


Inductive Systems on Complexes



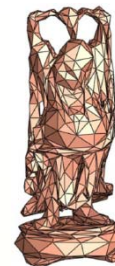
Idea: Follow basis elements from **birth** to **death** while maintaining **compatible bases**

Consistent Bases Exist



Persistent Homology

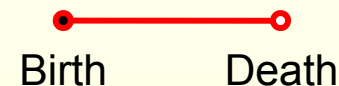
[Zomorodian, Edelsbrunner, Letcher 2002]



● Homology: $H_k(K^I) = Z_k(K^I) / B_k(K^I)$

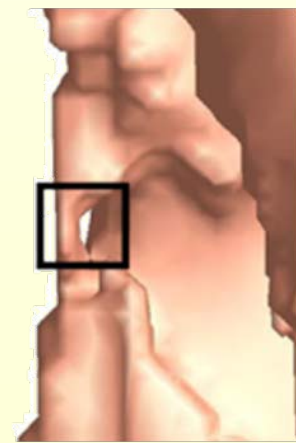
● The p -persistent k -th Homology group

$$H_k^{I, p} = Z_k^I / (B_k^{I+p} \cap Z_k^I)$$

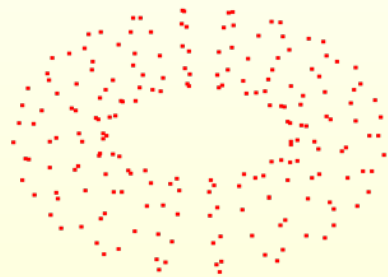


Persistent topological features are part of the shape; transient ones may be noise.

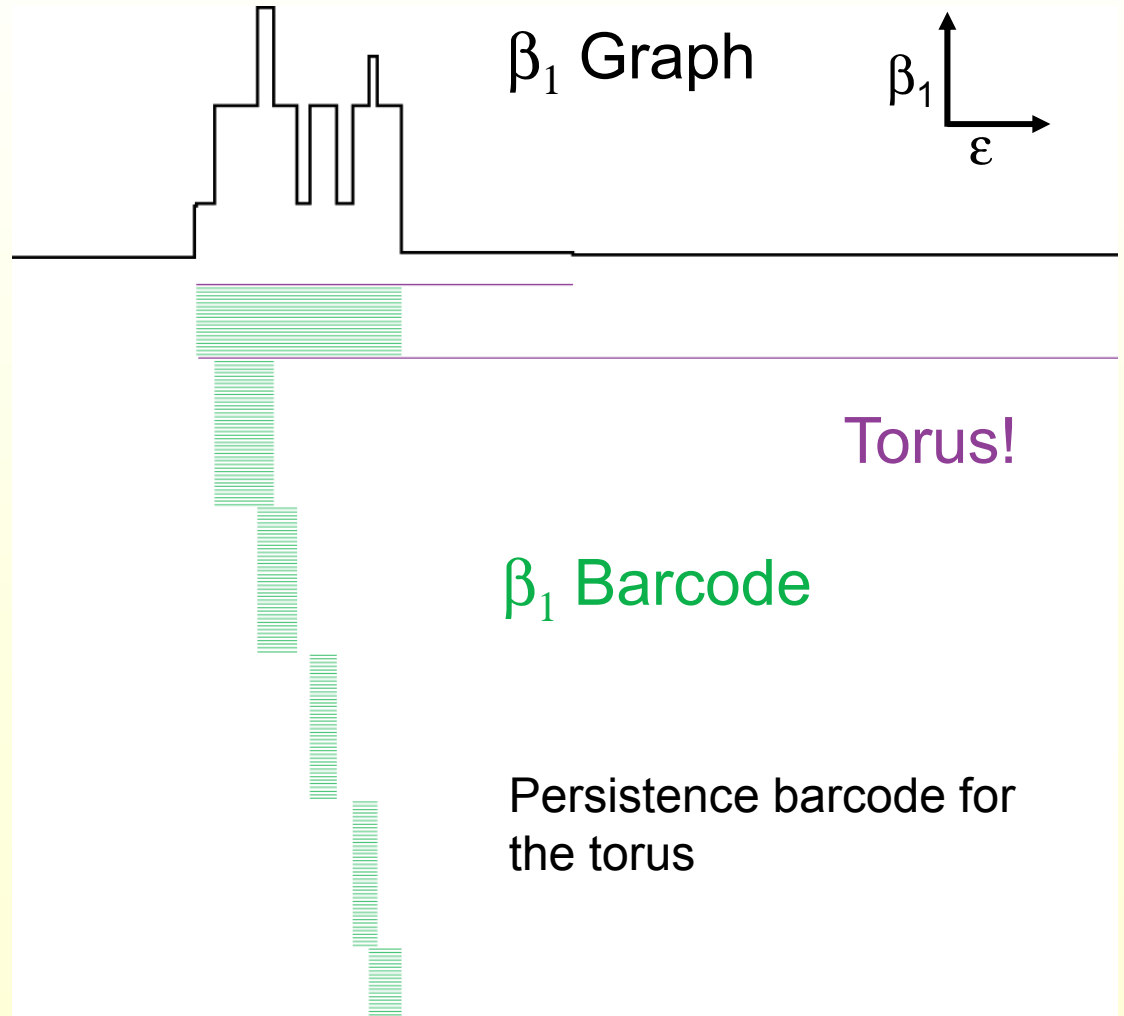
● Persistence Barcode: multiset of intervals



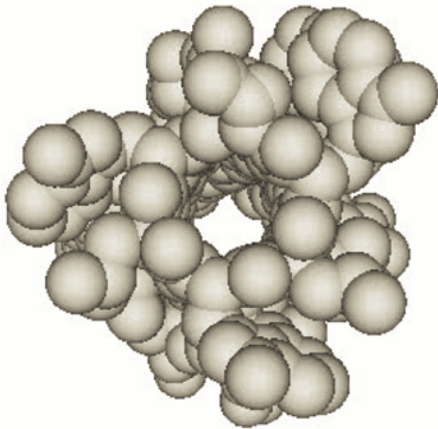
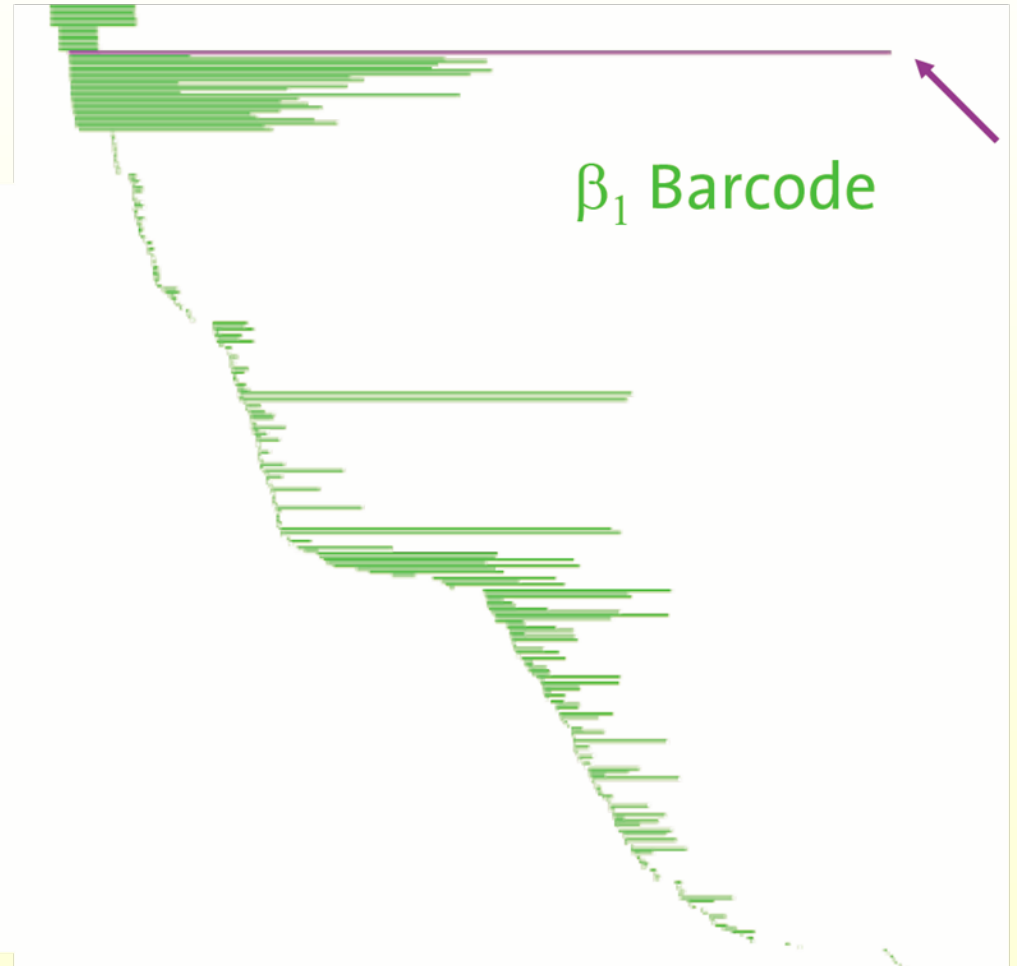
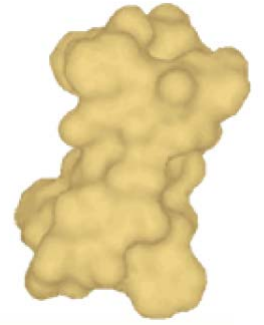
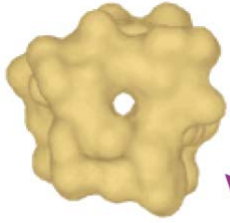
Deconstructing the Graph



PCD



3D Structure Discovery: Gramicidin A



Making Topology a Finer Tool

Geometry
discriminating

Topology
classifying

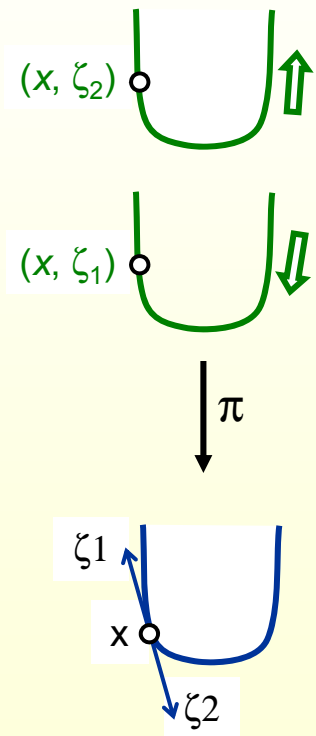
- Topology: connectivity of a space
- Key Idea: no reason to look at the original space only
 - Add geometry \Rightarrow look at **derived space(s)**
 - **Compute topology of derived space(s)**

1. Find filtration
2. Compute persistence

via the **tangent complex**

Our recipe

2-D Curve Tangent Complex

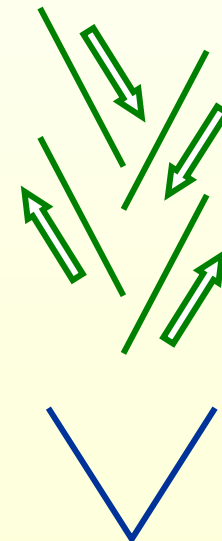


$T(X)$ has **two** components:
 $\beta_0(T(X)) = 2$

A corner point has four tangent directions:
 $\beta_0(T(X)) = 4$

There are **two** points in its fiber $\pi^{-1}(x)$

Every point x on a smooth curve X has **two** tangent directions.

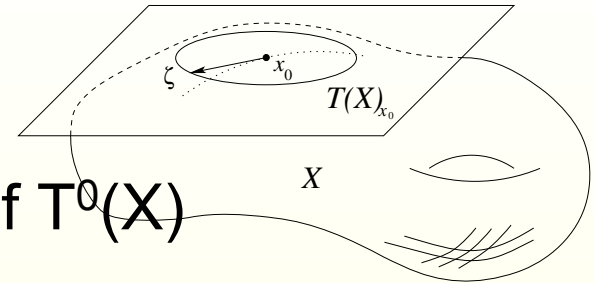


3-D Curvature-Filtered Tangent Complex

- Derived space

- $T^0(X)$: space of (point, tangent)

- Tangent complex $T(X)$** : closure of $T^0(X)$



- Filtration by increasing curvature

- Let $\rho(x, \zeta)$ be the radius of the circle of second order contact

- $T_\delta^0(X)$: points of $T^0(X)$ with $1/\rho \leq \delta$.

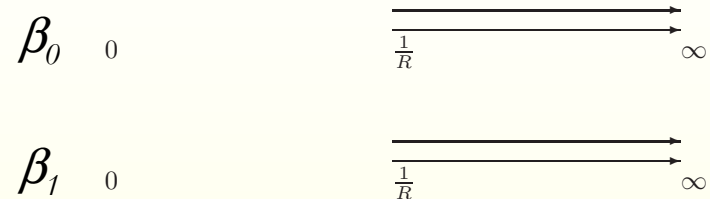
- $T_\delta(X)$: closure of $T_\delta^0(X)$

- Filtered tangent complex $T^{filt}(X)$** is the family

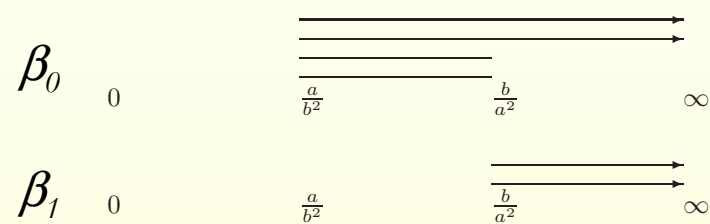
$$\{T_\delta(X)\}_{\delta \geq 0}$$

Persistence Barcodes: Circle vs. Ellipse

T^{filt} (circle of radius R) is simple:
the entire complex (2 copies of circle)
appears at once, at $\delta = 1/R$.

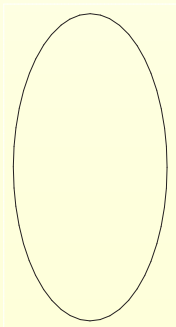


T^{filt} (ellipse) evolves through **four stages**: points at *lower* curvature appear earlier.

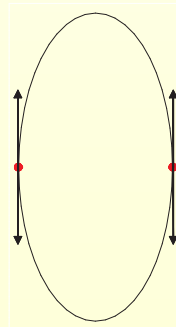


Persistence Barcodes

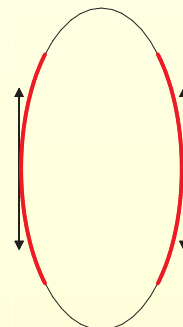
$$\frac{x^2}{a^2} + \frac{y^2}{b^2} = 1$$



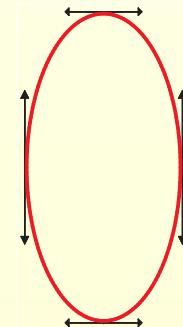
$$0 \leq \delta < \frac{a}{b^2}$$



$$\delta = \frac{a}{b^2}$$

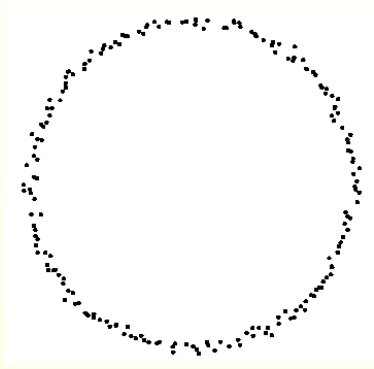


$$\frac{a}{b^2} < \delta < \frac{b}{a^2}$$

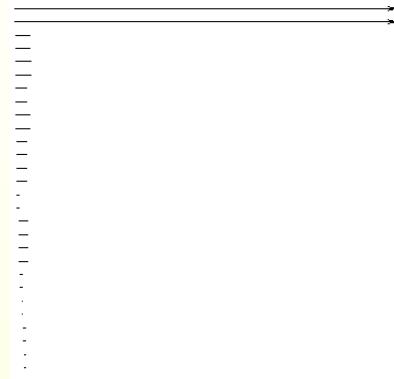


$$\delta \geq \frac{b}{a^2}$$

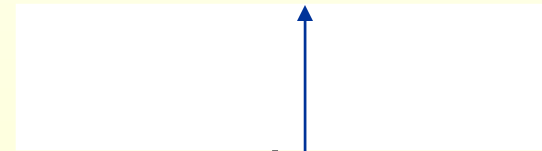
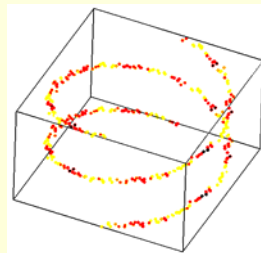
Applying Barcodes to 2D PCDs



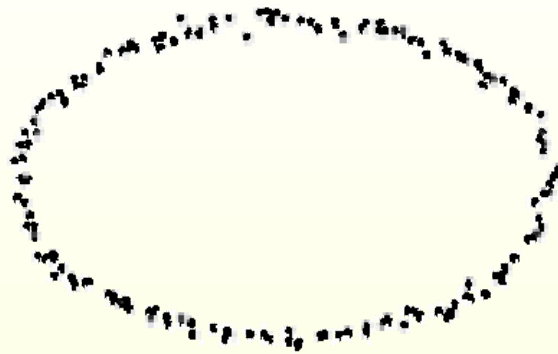
Input: Shape



Output: Descriptor

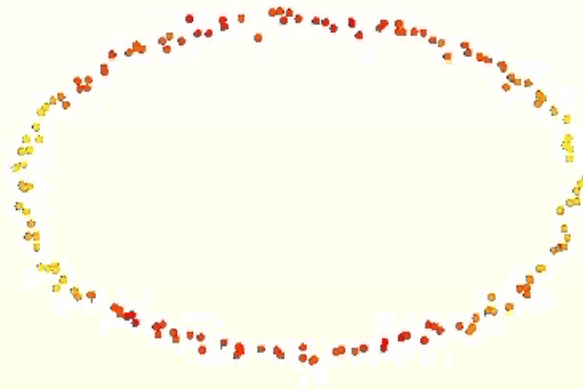


Fibers



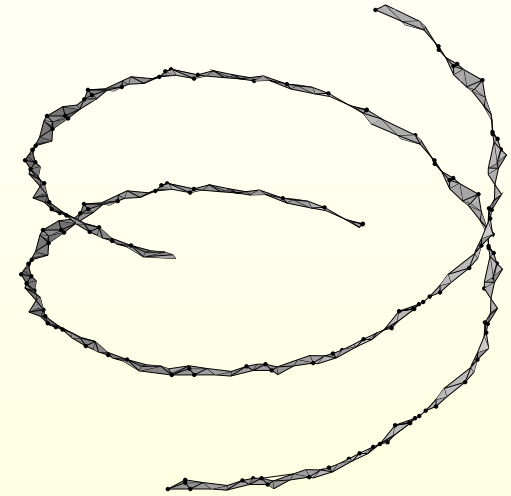
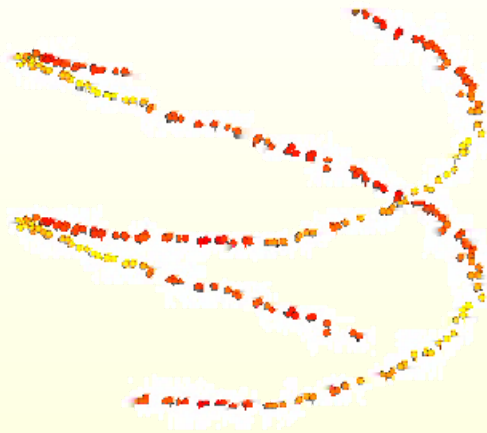
- PCD $P \subset X$, sampled from smooth closed 1-manifold
- We compute tangent fibers $\pi^{-1}(P)$ by normal estimation at each point

Filtering by Curvature



- Construct tangent complex incrementally
- Transform points to coordinate frame provided by tangent computation
- Fit osculating parabola to estimate curvature (more robust integral methods possible)

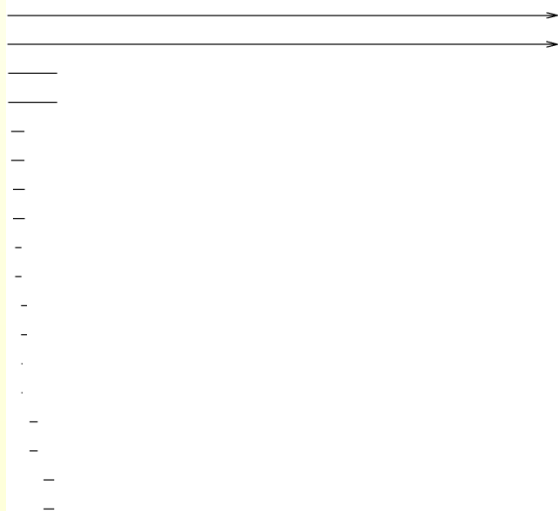
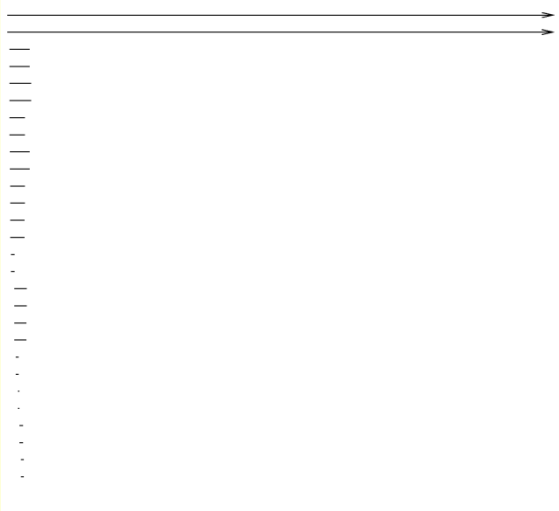
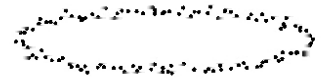
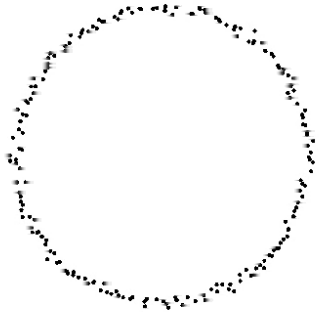
Approximating $T(X)$



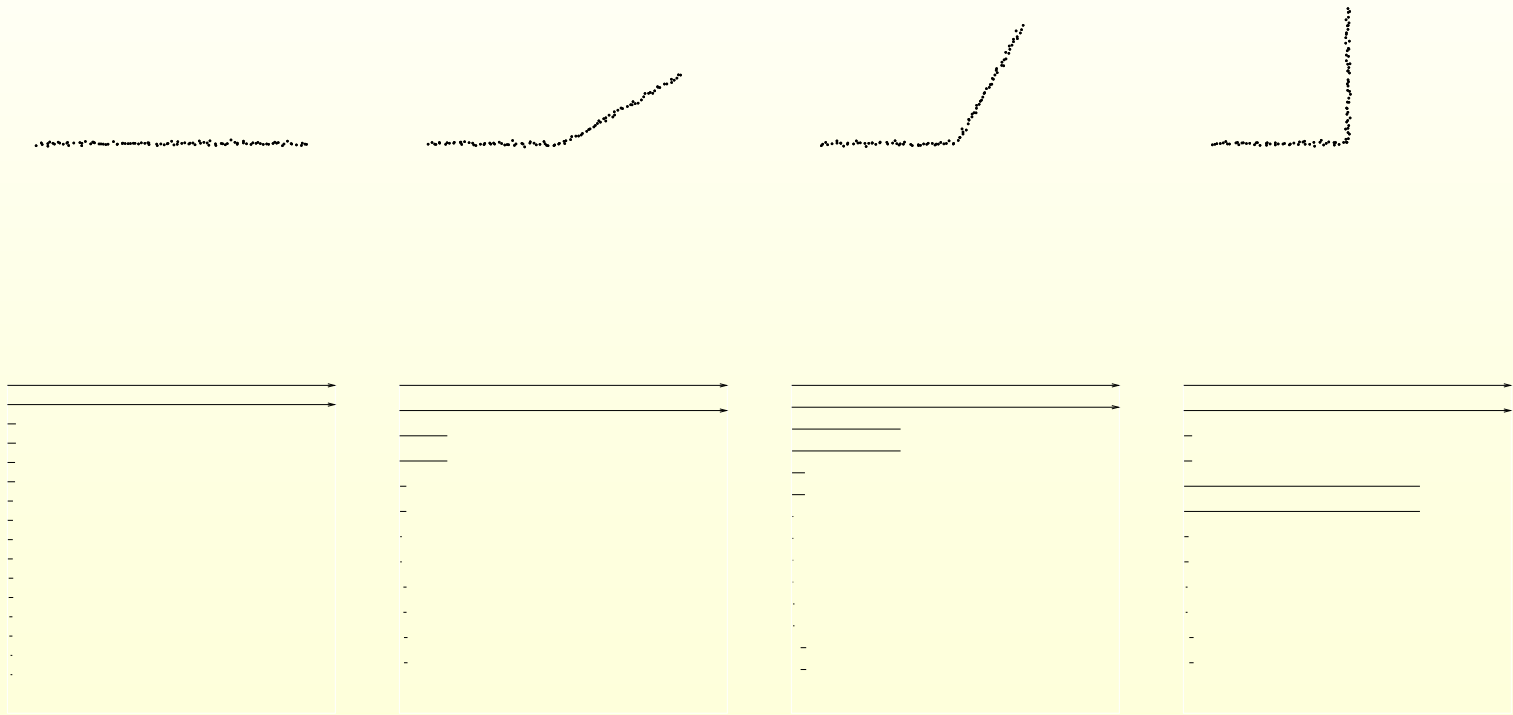
● $\mathbb{R}^n \times \mathbb{S}^{n-1}$ with $ds^2 = dx^2 + \omega^2 d\zeta^2$

● $T(X) \approx \bigcup_{p \in \pi^{-1}(P)} B_\varepsilon(p)$

Family of Ellipses



Articulated Arm Parametrization



Summary

- We are flooded by point set *data* and need to find structure in them
- *Topology* studies connectivity of spaces
- *Topological analysis* may be viewed as generalization of clustering
- To analyze point sets, we require a *combinatorial representation* approximating the original space
- *Homology* focuses on the structure of cycles
- *Persistent homology* analyzes the relationship of structures at multiple scales