Linear Systems and LU

CS 205A:
Mathematical Methods for Robotics, Vision, and Graphics

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Linear Systems

$$A\vec{x} = \vec{b}$$

$$A \in \mathbb{R}^{m \times n}$$

$$\vec{x} \in \mathbb{R}^{n}$$

$$\vec{b} \in \mathbb{R}^{m}$$

Case 1: Solvable

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} -1 \\ 1 \end{pmatrix}$$

Case 2: No Solution

$$\begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} -1 \\ 1 \end{pmatrix}$$

Case 3: Infinitely Many Solutions

$$\begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} -1 \\ -1 \end{pmatrix}$$

"Underdetermined"

No Other Cases

Proposition

If $A\vec{x} = \vec{b}$ has two distinct solutions $\vec{x}_0$ and $\vec{x}_1$, it has infinitely many solutions.

Common Misconception

Solvability *can* depend on $\vec{b}$!

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} -1 \\ 0 \end{pmatrix}$$

Dependence on Shape

Proposition

Tall matrices admit unsolvable right hand sides.

Proposition

Wide matrices admit right hand sides with infinite numbers of solutions.

For Now

All matrices will be:

- Square
- Invertible

Inverting Matrices

Do *not* compute A^{-1} if you do not need it.

- Not the same as solving $A\vec{x} = \vec{b}$
- Can be slow *and* poorly conditioned

Example

$$\begin{aligned} y - z &= -1 \\ 3x - y + z &= 4 \\ x + y - 2z &= -3 \end{aligned} \iff \left(\begin{array}{ccc|c} 0 & 1 & -1 & -1 \\ 3 & -1 & 1 & 4 \\ 1 & 1 & -2 & -3 \end{array}\right)$$

- Permute rows
- Row scaling
- Forward/back substitution

Row Operations: Permutation

$$\sigma : \{1, \ldots, m\} \to \{1, \ldots, m\}$$

$$P_\sigma \equiv \begin{pmatrix} - & \vec{e}_{\sigma(1)}^\top & - \\ - & \vec{e}_{\sigma(2)}^\top & - \\ & \cdots & \\ - & \vec{e}_{\sigma(m)}^\top & - \end{pmatrix}$$

Row Operations: Row Scaling

$$S_a \equiv \begin{pmatrix} a_1 & 0 & 0 & \cdots \\ 0 & a_2 & 0 & \cdots \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & a_m \end{pmatrix}$$

Row Operations: Elimination

"Scale row k by constant c and add result to row ℓ."

$$E \equiv I + c\vec{e}_\ell \vec{e}_k^{\top}$$

Solving via Elimination Matrices

$$\left(\begin{array}{ccc|c} 0 & 1 & -1 & -1 \\ 3 & -1 & 1 & 4 \\ 1 & 1 & -2 & -3 \end{array}\right)$$

Reverse order!

Introducing Gaussian Elimination

Big idea:

General strategy to solve linear systems via elimination matrices.[1]

[1] Implicitly.

Sequence of Invertible Steps

$$A\vec{x} = \vec{b}$$

$$E_1 A\vec{x} = E_1\vec{b}$$

$$E_2 E_1 A\vec{x} = E_2 E_1\vec{b}$$

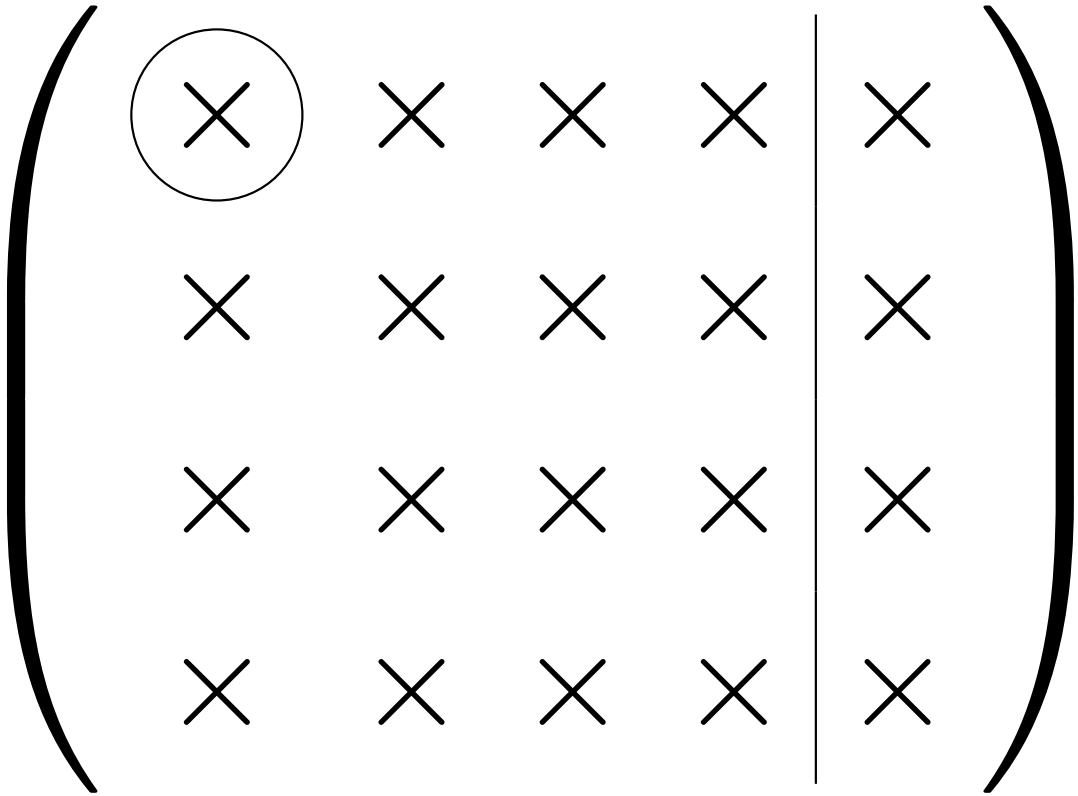
$$\vdots$$

$$E_k \cdots E_2 E_1 A\vec{x} = E_k \cdots E_2 E_1\vec{b}$$

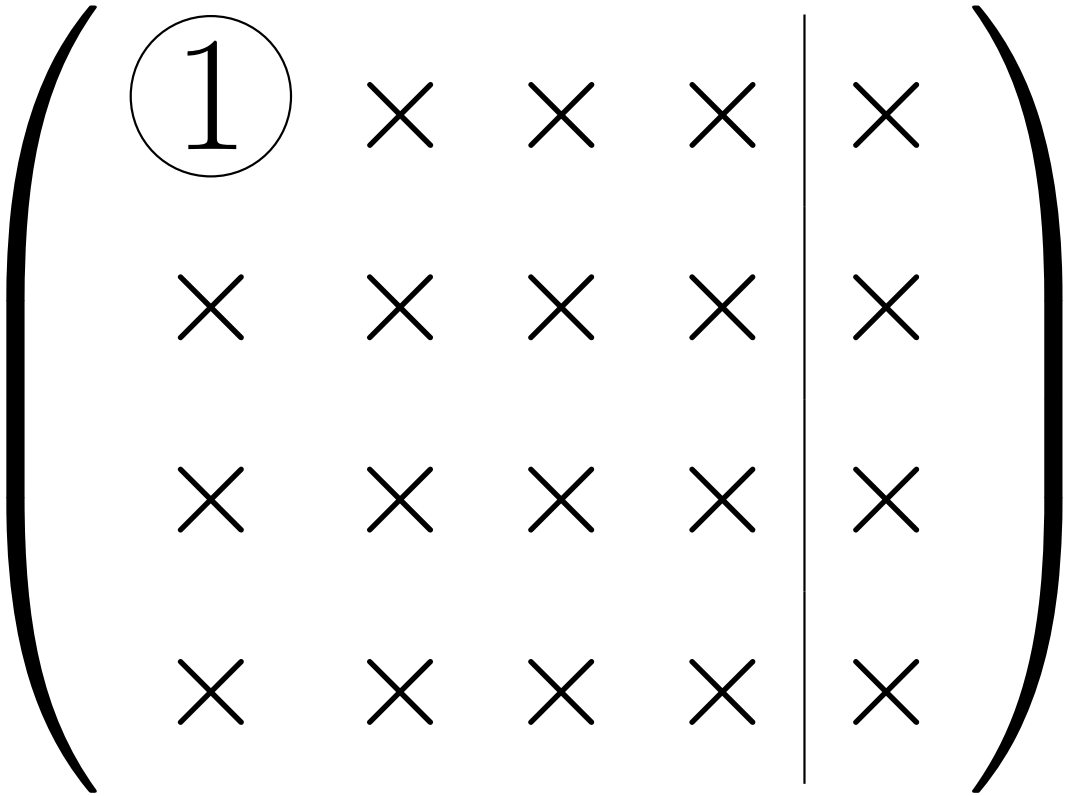
Shape of Systems

$$\left(A \middle| \vec{b} \right) = \left(\begin{array}{cccc|c} \times & \times & \times & \times & \times \\ \times & \times & \times & \times & \times \\ \times & \times & \times & \times & \times \\ \times & \times & \times & \times & \times \end{array} \right)$$

Pivot



Row Scaling



Forward Substitution

$$\left(\begin{array}{cccc|c} \textcircled{1} & \times & \times & \times & \times \\ 0 & \times & \times & \times & \times \\ 0 & \times & \times & \times & \times \\ 0 & \times & \times & \times & \times \end{array}\right)$$

Forward Substitution

$$\left(\begin{array}{cccc|c} 1 & \times & \times & \times & \times \\ 0 & \textcircled{1} & \times & \times & \times \\ 0 & 0 & \times & \times & \times \\ 0 & 0 & \times & \times & \times \end{array}\right)$$

Upper Triangular Form

$$\left(\begin{array}{cccc|c} 1 & \times & \times & \times & \times \\ 0 & 1 & \times & \times & \times \\ 0 & 0 & 1 & \times & \times \\ 0 & 0 & 0 & \textcircled{1} & \times \end{array}\right)$$

Back Substitution

$$\left(\begin{array}{cccc|c} 1 & \times & \times & 0 & \times \\ 0 & 1 & \times & 0 & \times \\ 0 & 0 & 1 & 0 & \times \\ 0 & 0 & 0 & \textcircled{1} & \times \end{array}\right)$$

Back Substitution

$$\left(\begin{array}{cccc|c} 1 & \times & 0 & 0 & \times \\ 0 & 1 & 0 & 0 & \times \\ 0 & 0 & \textcircled{1} & 0 & \times \\ 0 & 0 & 0 & 1 & \times \end{array}\right)$$

Back Substitution

$$\left(\begin{array}{cccc|c} \textcircled{1} & 0 & 0 & 0 & \times \\ 0 & 1 & 0 & 0 & \times \\ 0 & 0 & 1 & 0 & \times \\ 0 & 0 & 0 & 1 & \times \end{array}\right)$$

Steps of Gaussian Elimination

1. Forward substitution: For each row $i = 1, 2, \ldots, m$
   - Scale row to get pivot 1
   - For each $j > i$, subtract multiple of row i from row j to zero out pivot column
2. Backward substitution: For each row $i = m, m-1, \ldots, 1$
   - For each $j < i$, zero out rest of column

Total Running Time

$$O\left(n^3\right)$$

Problem

$$A = \begin{pmatrix} \textcircled{0} & 1 \\ 1 & 0 \end{pmatrix}$$

Even Worse

$$A = \begin{pmatrix} \varepsilon & 1 \\ 1 & 0 \end{pmatrix}$$

Pivoting

Pivoting

Permuting rows and/or columns to avoid dividing by small numbers or zero.

- *Partial* pivoting
- *Full* pivoting

$$\begin{pmatrix} 1 & 10 & -10 \\ 0 & \textcircled{0.1} & 9 \\ 0 & 4 & 6.2 \end{pmatrix}$$

Reasonable Use Case

$$\begin{aligned} A\vec{x}_1 &= \vec{b}_1 \\ A\vec{x}_2 &= \vec{b}_2 \\ &\vdots \end{aligned}$$

Can we do *precomputation* on A to avoid redoing the same work?

Do we have to do $O(n^3)$ steps each time?

Observation

Steps of Gaussian elimination depend only on structure of A.

Solving Upper Triangular Systems

$$\left(\begin{array}{cccc|c} 1 & \times & \times & \times & \times \\ 0 & 1 & \times & \times & \times \\ 0 & 0 & 1 & \times & \times \\ 0 & 0 & 0 & \textcircled{1} & \times \end{array}\right)$$

After Back Substitution

$$\left(\begin{array}{cccc|c} 1 & \times & \times & 0 & \times \\ 0 & 1 & \times & 0 & \times \\ 0 & 0 & 1 & 0 & \times \\ \textcircled{0} & \textcircled{0} & \textcircled{0} & 1 & \times \end{array}\right)$$

No need to subtract the 0's explicitly!

$O(n)$ time

Next Pivot: Same Observation

$$\left(\begin{array}{cccc|c} 1 & \times & 0 & 0 & \times \\ 0 & 1 & 0 & 0 & \times \\ \textcircled{0} & \textcircled{0} & 1 & \textcircled{0} & \times \\ 0 & 0 & 0 & 1 & \times \end{array}\right)$$

Observation

Triangular systems can be solved in $O(n^2)$ time.

Upper Triangular Part of A

$$A\vec{x} = \vec{b}$$

$$\vdots$$

$$M_k \cdots M_1 A\vec{x} = M_k \cdots M_1 \vec{b}$$

Define:

$$U \equiv M_k \cdots M_1 A$$

Lower Triangular Part

$$
\begin{aligned}
U &= M_k \cdots M_1 A \\
\Rightarrow A &= \left(M_1^{-1} \cdots M_k^{-1}\right) U \\
&\equiv LU
\end{aligned}
$$

Why Is L Triangular?

$$S \equiv \mathrm{diag}(a_1, a_2, \ldots)$$

$$E \equiv I + c\vec{e}_\ell \vec{e}_k^{\top}$$

Proposition

The product of triangular matrices is triangular.

Solving Systems Using LU

$$A\vec{x} = \vec{b}$$

$$\Rightarrow LU\vec{x} = \vec{b}$$

1. Solve $L\vec{y} = \vec{b}$ using forward substitution.
2. Solve $U\vec{x} = \vec{y}$ using backward substitution.

$O(n^2)$ (given LU factorization)

LU: Compact Storage

$$\begin{pmatrix} U & U & U & U \\ L & U & U & U \\ L & L & U & U \\ L & L & L & U \end{pmatrix}$$

Assumption: Diagonal elements of L are 1.

Warning: Do not multiply this matrix!

Computing LU Factorization

Small modification of forward substitution step to keep track of L.[2]

[2] See lecture notes for pseudocode.

LU with Pivoting

$$PAQ = LU$$

Next