Optimization III: Constrained Optimization

CS 205A:
Mathematical Methods for Robotics, Vision, and Graphics

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Constrained Problems

$$\begin{aligned} \text{minimize} \quad & f(\vec{x}) \\ \text{such that} \quad & g(\vec{x}) = \vec{0} \\ & h(\vec{x}) \geq \vec{0} \end{aligned}$$

Really Difficult!

Simultaneously:

- Minimizing f
- Finding roots of g
- Finding feasible points of h

Implicit Projection

Implicit surface: $g(\vec{x}) = 0$

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$$\begin{aligned} \text{minimize} \quad & \|\vec{x} - \vec{x}_0\|_2 \\ \text{such that} \quad & g(\vec{x}) = \vec{0} \end{aligned}$$

Nonnegative Least-Squares

$$\text{minimize}_{\vec{x}} \; \|A\vec{x} - \vec{b}\|_2^2$$
$$\text{such that} \;\; \vec{x} \geq \vec{0}$$

Manufacturing

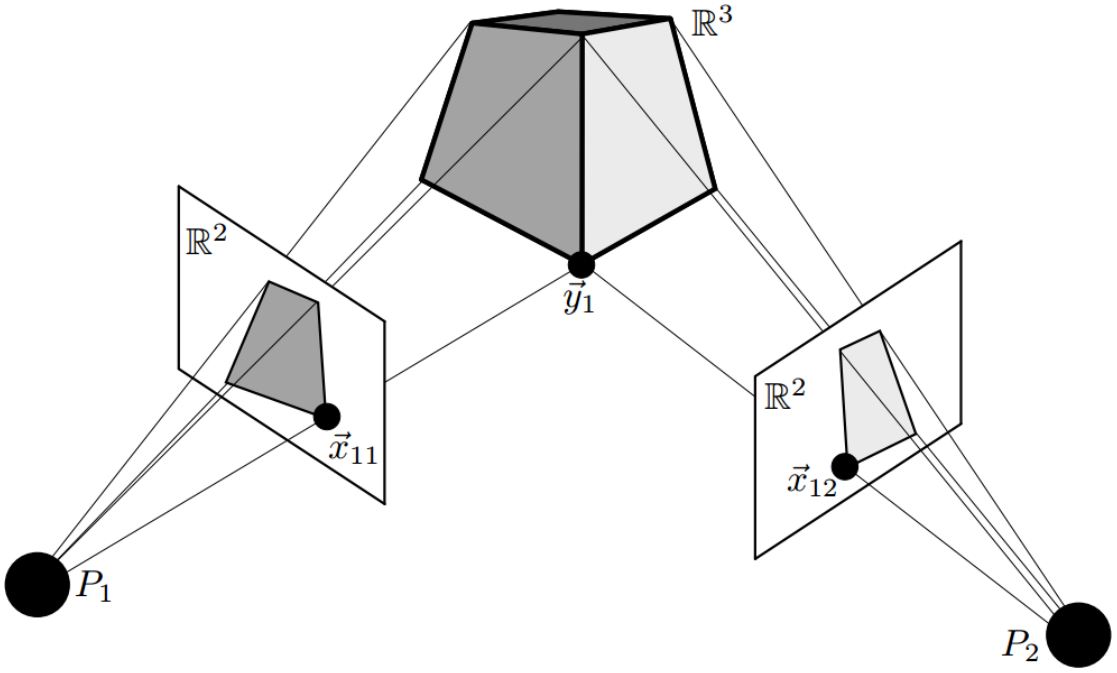
- m materials
- s_i units of material i in stock
- n products
- p_j profit for product j
- Product j uses c_{ij} units of material i

Manufacturing

$$\begin{array}{ll} \text{maximize}_{\vec{x}} & \sum_j p_j x_j \\ \text{such that} & x_j \geq 0 \ \forall j \\ & \sum_j c_{ij} x_j \leq s_i \ \forall i \end{array}$$

"Maximize profits where you make a positive amount of each product and use limited material."

Bundle Adjustment



$$\min_{\vec{y}_j, P_i} \sum_{ij} \|P_i \vec{y}_j - \vec{x}_{ij}\|_2^2$$
$$\text{s.t.} \quad P_i \text{ orthogonal } \forall i$$

Basic Definitions

Feasible point and feasible set

A *feasible point* is any point $\vec{x}$ satisfying $g(\vec{x}) = \vec{0}$ and $h(\vec{x}) \geq \vec{0}$. The *feasible set* is the set of all points $\vec{x}$ satisfying these constraints.

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Critical point of constrained optimization

A critical point is one satisfying the constraints that also is a local maximum, minimum, or saddle point of f within the feasible set.

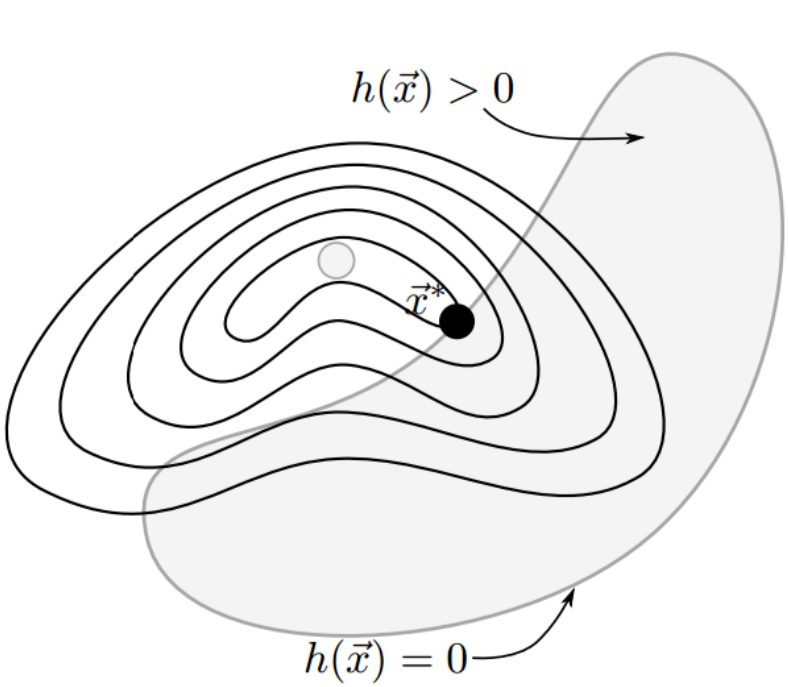
Differential Optimality

Without h:

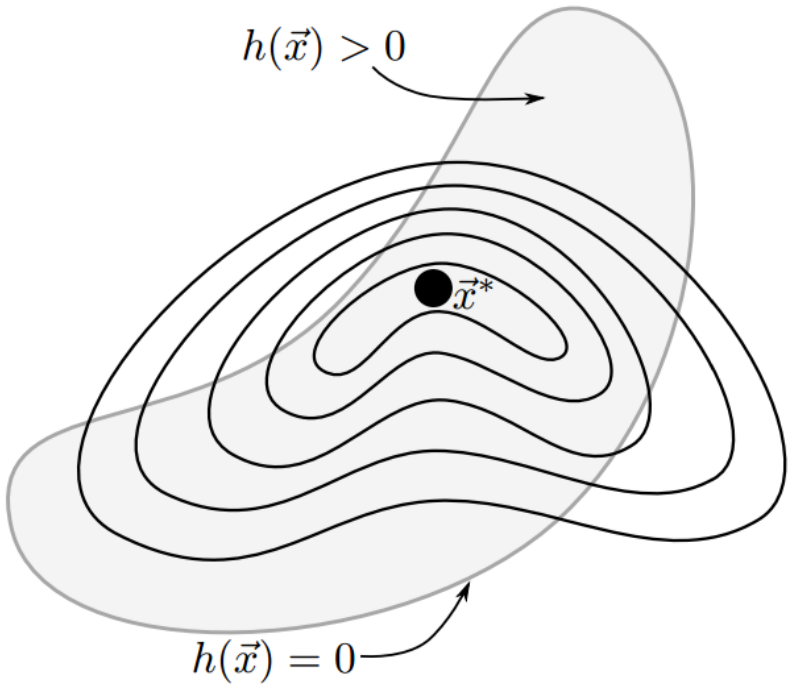
$$\Lambda(\vec{x}, \vec{\lambda}) \equiv f(\vec{x}) - \vec{\lambda} \cdot g(\vec{x})$$

Lagrange Multipliers

Inequality Constraints at $\vec{x}^*$



Active constraint
$h(\vec{x}^*) = 0$



Inactive constraint
$h(\vec{x}^*) > 0$

Inequality Constraints at $\vec{x}^*$

Two cases:

- **Active:** $h_i(\vec{x}^*) = 0$
  Optimum might change if constraint is removed
- **Inactive:** $h_i(\vec{x}^*) > 0$
  Removing constraint does not change $\vec{x}^*$ locally

Idea

Remove inactive constraints and make active constraints equality constraints.

Lagrange Multipliers

$$\Lambda(\vec{x}, \vec{\lambda}, \vec{\mu}) \equiv f(\vec{x}) - \vec{\lambda} \cdot g(\vec{x}) - \vec{\mu} \cdot h(\vec{x})$$

No longer a critical point! But if we ignore that:

$$\vec{0} = \nabla f(\vec{x}) - \sum_i \lambda_i \nabla g_i(\vec{x}) - \sum_j \mu_j \nabla h_j(\vec{x})$$

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$$\mu_j h_j(\vec{x}) = \vec{0}$$

Zero out inactive constraints!

Inequality Direction

So far: Have not distinguished between $h_j(\vec{x}) \geq 0$ and $h_j(\vec{x}) \leq 0$

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- Direction to decrease f: $-\nabla f(\vec{x}^*)$
- Direction to decrease h_j: $-\nabla h_j(\vec{x}^*)$

Inequality Direction

So far: Have not distinguished between

$$h_j(\vec{x}) \geq 0 \text{ and } h_j(\vec{x}) \leq 0$$

- Direction to decrease f: $-\nabla f(\vec{x}^*)$
- Direction to decrease h_j: $-\nabla h_j(\vec{x}^*)$

$$\nabla f(\vec{x}^*) \cdot \nabla h_j(\vec{x}^*) \geq 0$$

Dual Feasibility

$$\mu_j \geq 0$$

KKT Conditions

Theorem (Karush-Kuhn-Tucker (KKT) conditions)

$\vec{x}^* \in \mathbb{R}^n$ is a critical point when there exist $\vec{\lambda} \in \mathbb{R}^m$ and $\vec{\mu} \in \mathbb{R}^p$ such that:

- $\vec{0} = \nabla f(\vec{x}^*) - \sum_i \lambda_i \nabla g_i(\vec{x}^*) - \sum_j \mu_j \nabla h_j(\vec{x}^*)$ *("stationarity")*
- $g(\vec{x}^*) = \vec{0}$ *and* $h(\vec{x}) \geq \vec{0}$ *("primal feasibility")*
- $\mu_j h_j(\vec{x}^*) = 0$ *for all* j *("complementary slackness")*
- $\mu_j \geq 0$ *for all* j *("dual feasibility")*

Sequential Quadratic Programming (SQP)

$$\vec{x}_{k+1} \equiv \vec{x}_k + \arg\min_{\vec{d}} \left[\frac{1}{2} \vec{d}^\top H_f(\vec{x}_k) \vec{d} + \nabla f(\vec{x}_k) \cdot \vec{d} \right]$$

$$\text{such that } g_i(\vec{x}_k) + \nabla g_i(\vec{x}_k) \cdot \vec{d} = 0$$

$$h_i(\vec{x}_k) + \nabla h_i(\vec{x}_k) \cdot \vec{d} \geq 0$$

Equality Constraints Only

$$\begin{pmatrix} H_f(\vec{x}_k) & [Dg(\vec{x}_k)]^\top \\ Dg(\vec{x}_k) & 0 \end{pmatrix} \begin{pmatrix} \vec{d} \\ \vec{\lambda} \end{pmatrix} = \begin{pmatrix} -\nabla f(\vec{x}_k) \\ -g(\vec{x}_k) \end{pmatrix}$$

- Can approximate H_f
- Can limit distance along $\vec{d}$

Inequality Constraints

Active set methods:

Keep track of active constraints and enforce as equality, update based on gradient

Barrier Methods: Equality Case

$$f_\rho(\vec{x}) \equiv f(\vec{x}) + \rho \|g(\vec{x})\|_2^2$$

Unconstrained optimization, crank up ρ until $g(\vec{x}) \approx \vec{0}$

Caveat: H_{f_ρ} becomes poorly conditioned

Barrier Methods: Inequality Case

$$\text{Inverse barrier:} \quad \frac{1}{h_i(\vec{x})}$$

$$\text{Logarithmic barrier:} \quad -\log h_i(\vec{x})$$

To Read: Convex Programming

A ray of hope:

Minimizing convex functions with convex constraints

Next