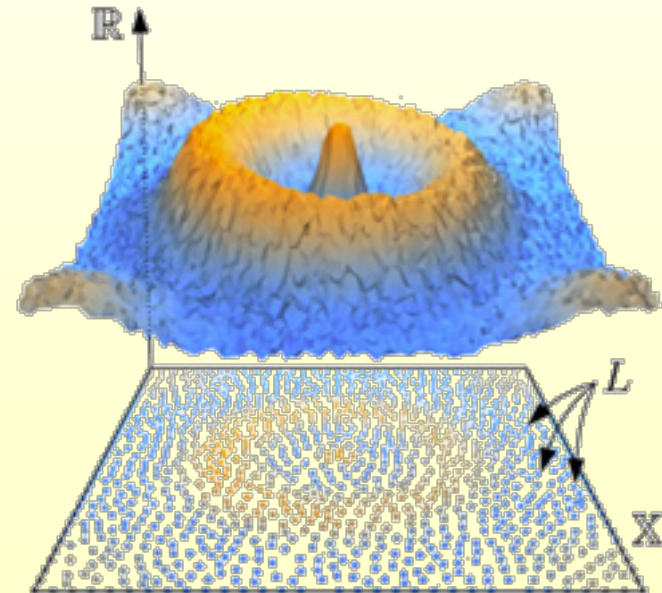


# CS268: Computational Topology and Topological Data Analysis, I

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Leonidas J. Guibas



# Computational Topology



Herbert Edelsbrunner



Gunnar Carlsson



Afra Zomorodian

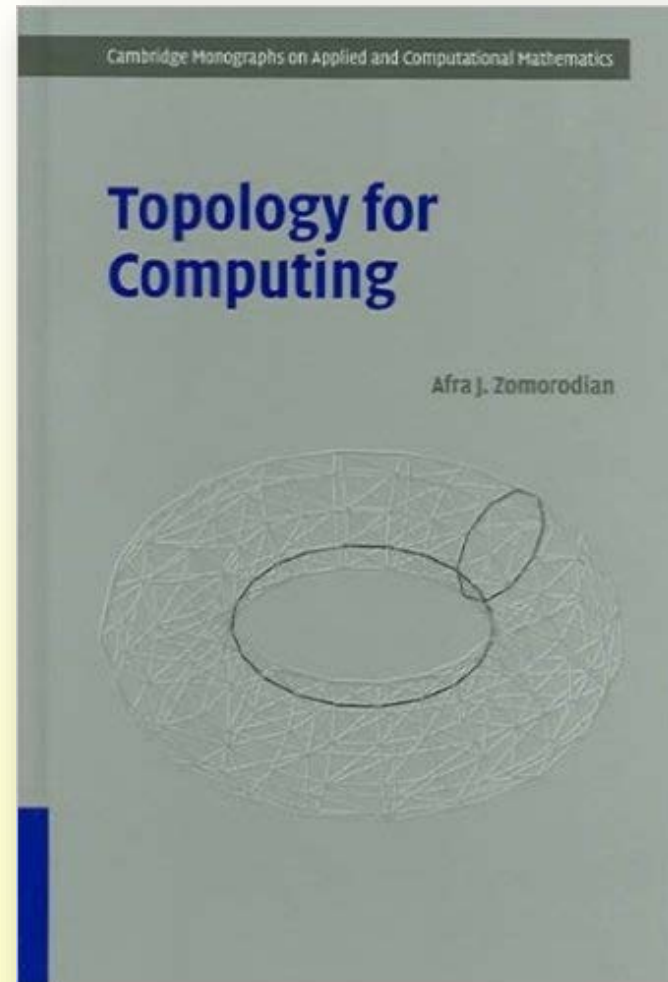
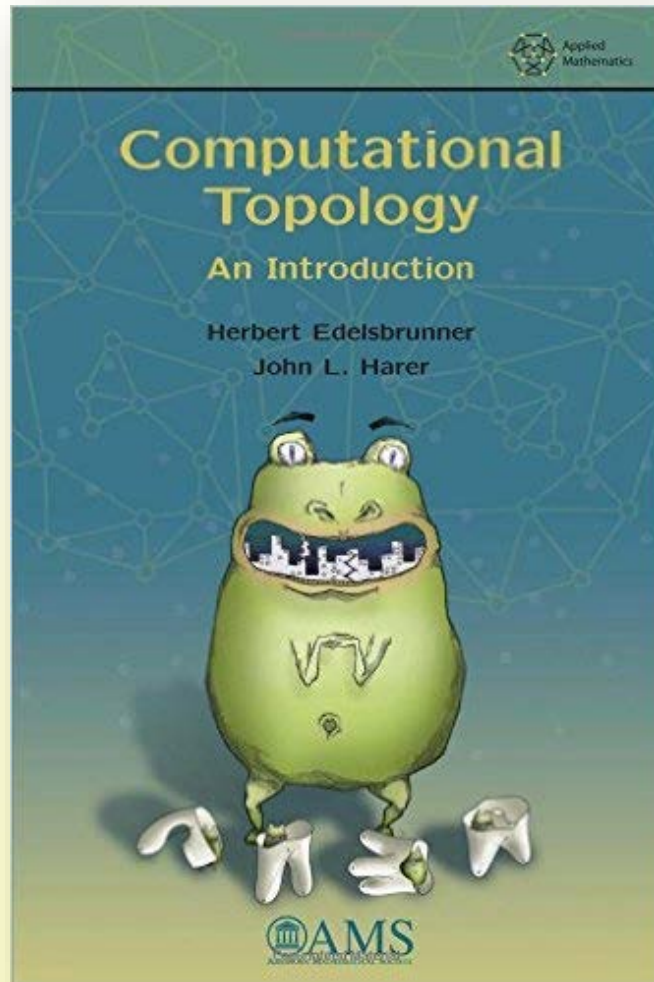


Frederic Chazal



Robert Ghrist

# Some Textbooks

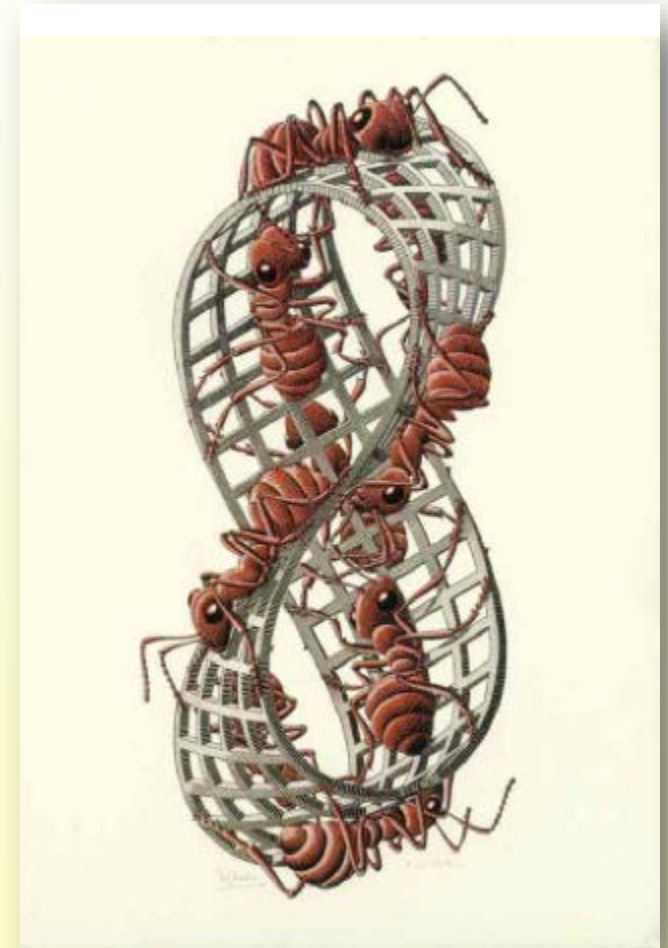


# Topological Data Analysis (TDA)

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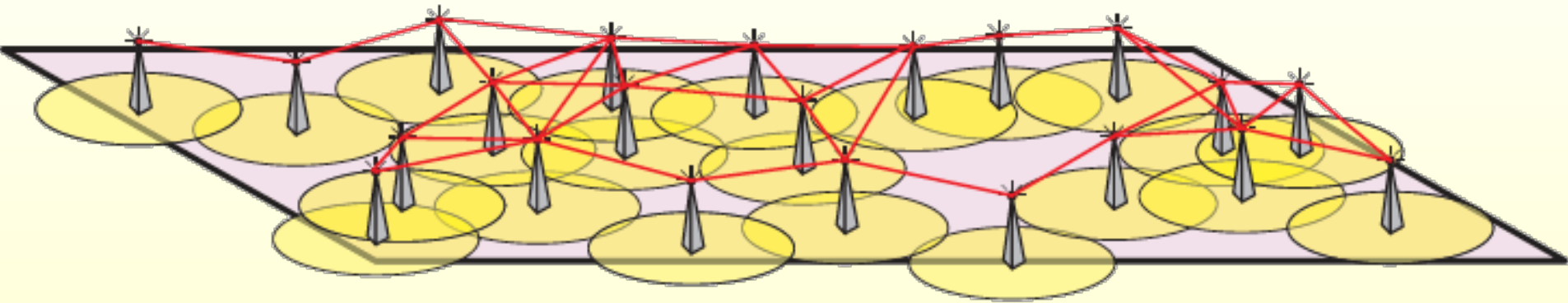
Topology — Computational  
Topology

Homology — Persistent Homology



Slides ack: Afra Zomorodian, Ryan Lewis,  
Robert Ghrist

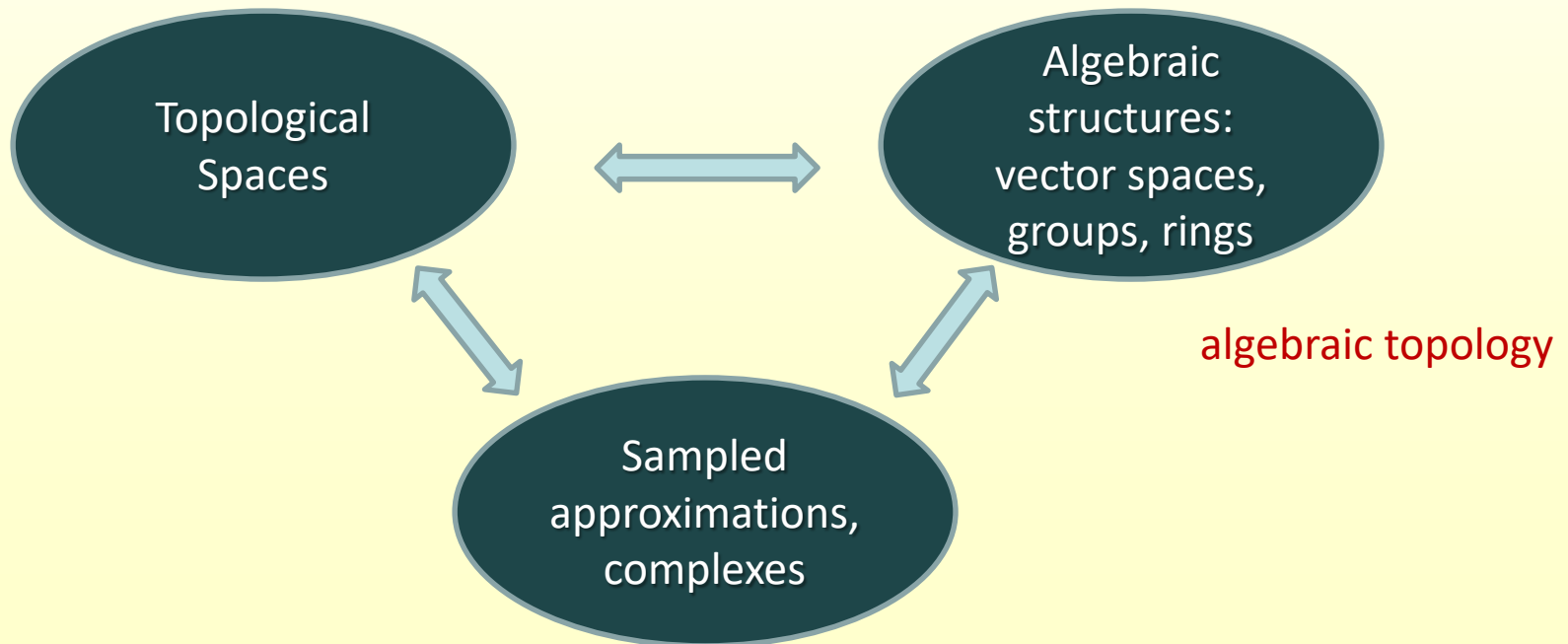
# Homological Sensor Networks



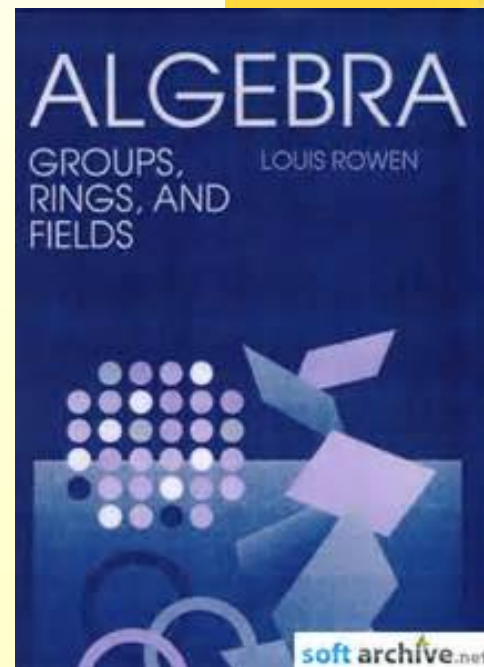
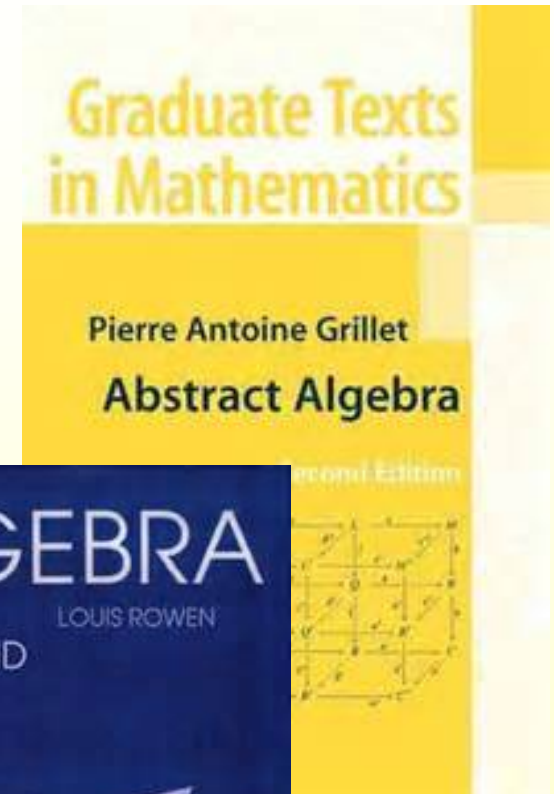
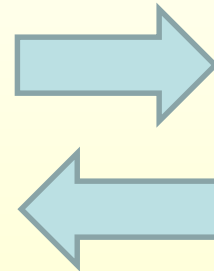
It's all linear algebra ...

# Topology and Topology Inference

- ◆ Topology is the branch of mathematics that studies the connectivity of spaces, and the obstructions to such connectivity
- ◆ Topology studies **global** structure



# From Data to Algebraic Objects



Homology groups as  
data descriptors

# Topology



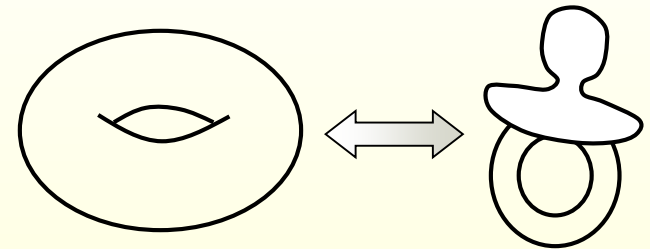
The bridges of Königsberg



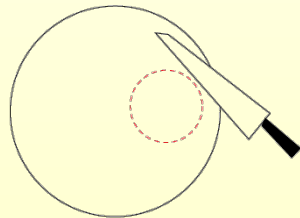
# Connectivity for 2-Manifolds: Ordinary Surfaces

- ◆ Topology does not take distances too seriously – we are allowed to stretch and shrink

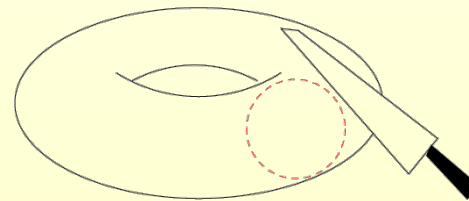
- ◆ Homeomorphism: 1-1, onto, bi-continuous



- ◆ But we care about cutting, puncturing, stitching, gluing ...



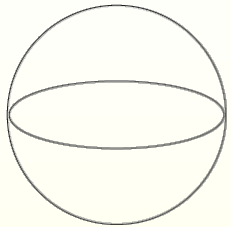
(a) No matter where we cut the sphere, we get two pieces



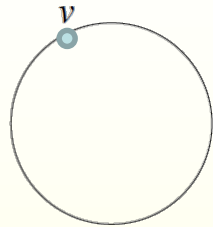
(b) If we're careful, we can cut the torus and still leave it in one piece.

- ◆ Note: connectivity information is indexed by dimension

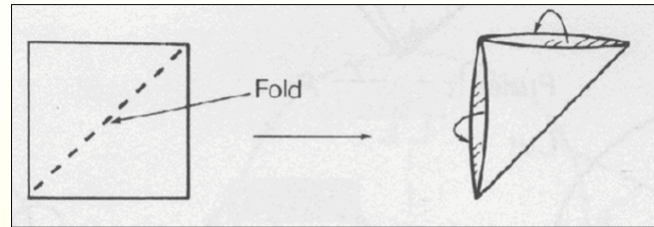
# 2-Manifold Zoo



(a)  $\{x \in \mathbb{R}^3 \mid |x| = 1\}$

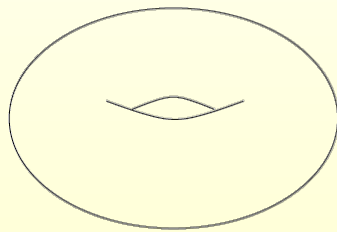


(b) Identify boundary to  $v$

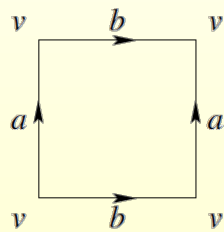


(c) Instructions for a flat sphere

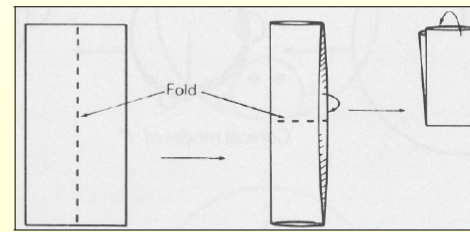
Sphere



(a) Donut surface



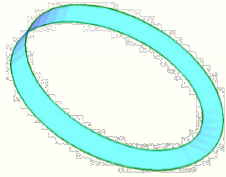
(b) Diagram



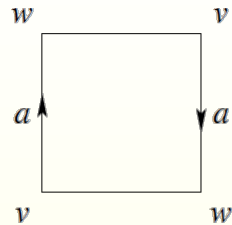
(c) Instructions for a flat torus

Torus

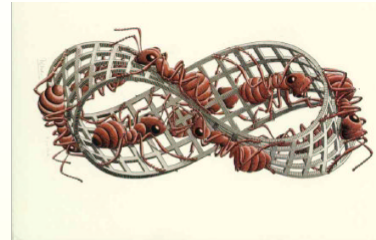
# More Exotic Animals



(a) Embedded

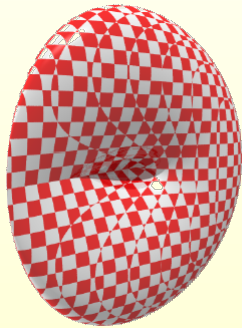


(b) Diagram

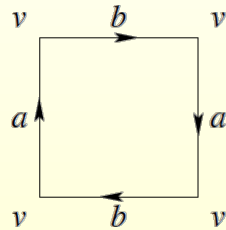


(c) Escher's *Möbius Strip II* (on its side)

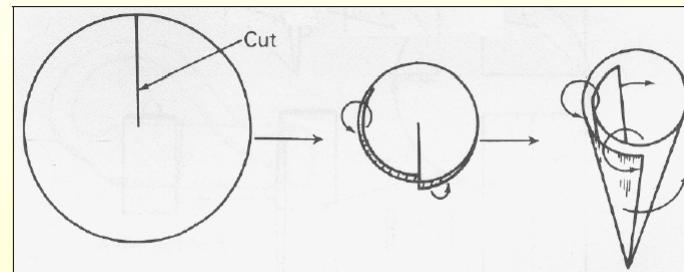
Möbius strip



Cross cap + disk

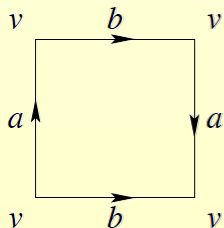


(a) Diagram

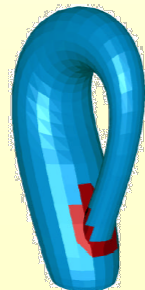


(b) Instructions for a flat  $\mathbb{RP}^2$

Projective plane



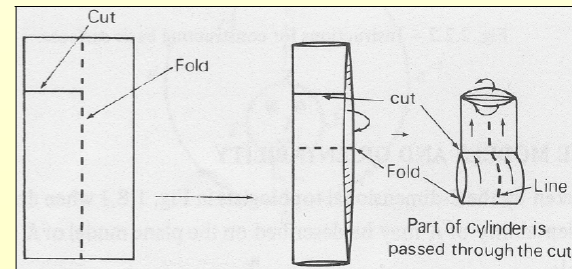
(a) Diagram



(b) An immersion



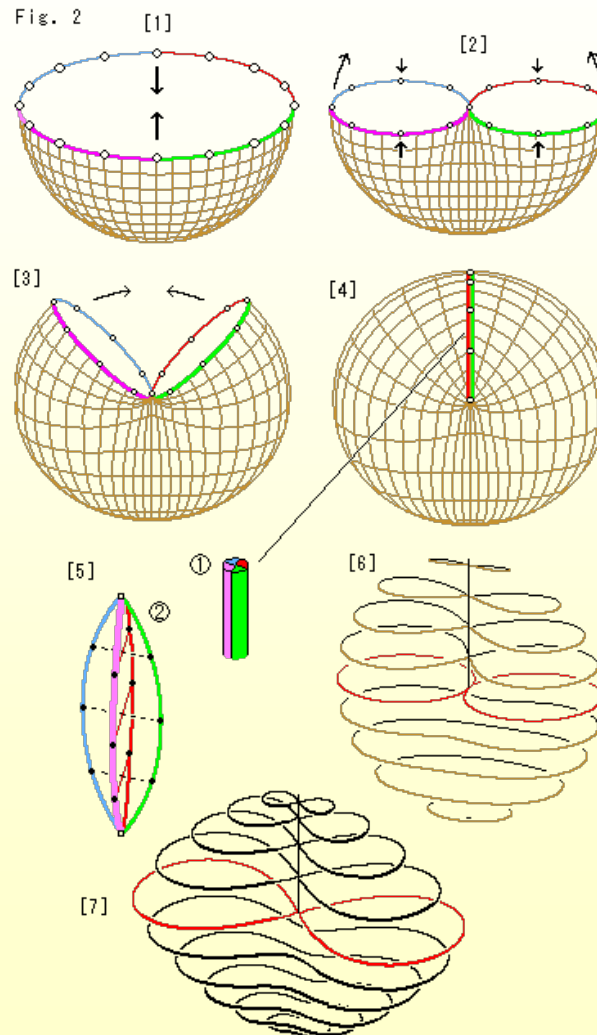
(c) Cut in half (a Möbius strip)



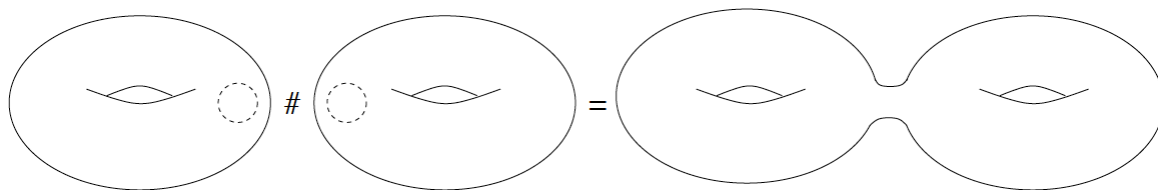
(d) Instructions for a flat  $\mathbb{K}^2$

Klein bottle

# Projective Plane



# Connected Sums

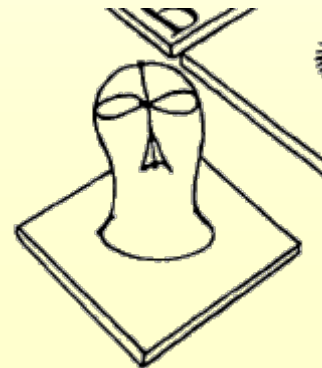


- ◆ **Classification Theorem of 2-Manifolds (1860):** Every closed connected compact surface is a connected sum of a sphere with a number of tori and projective planes (sphere + handles + cross cups)



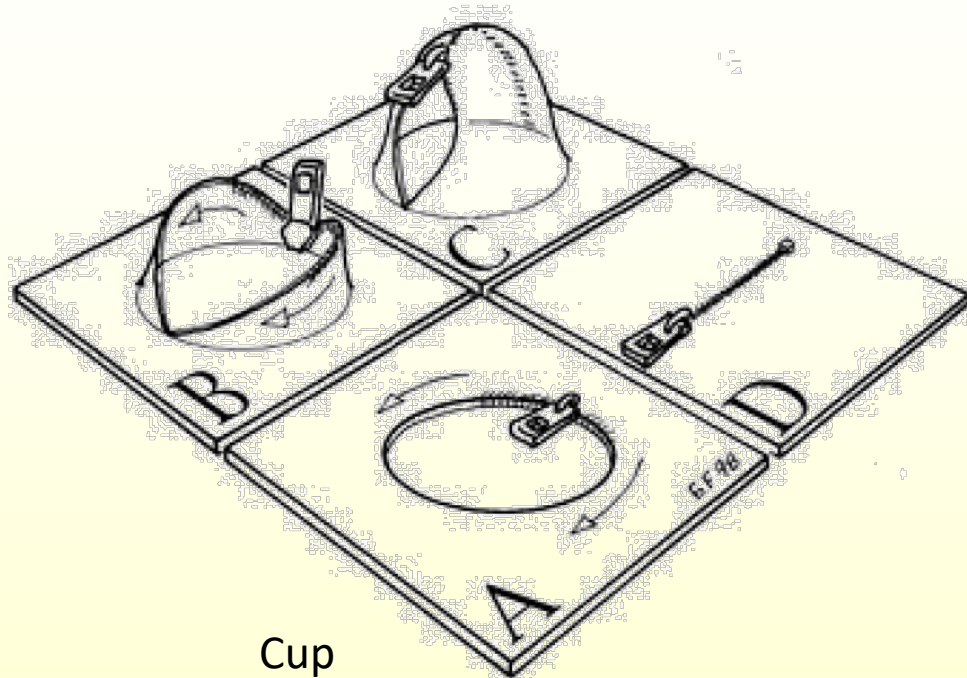
Handle

Klein bottle = sphere + 2 cross cups

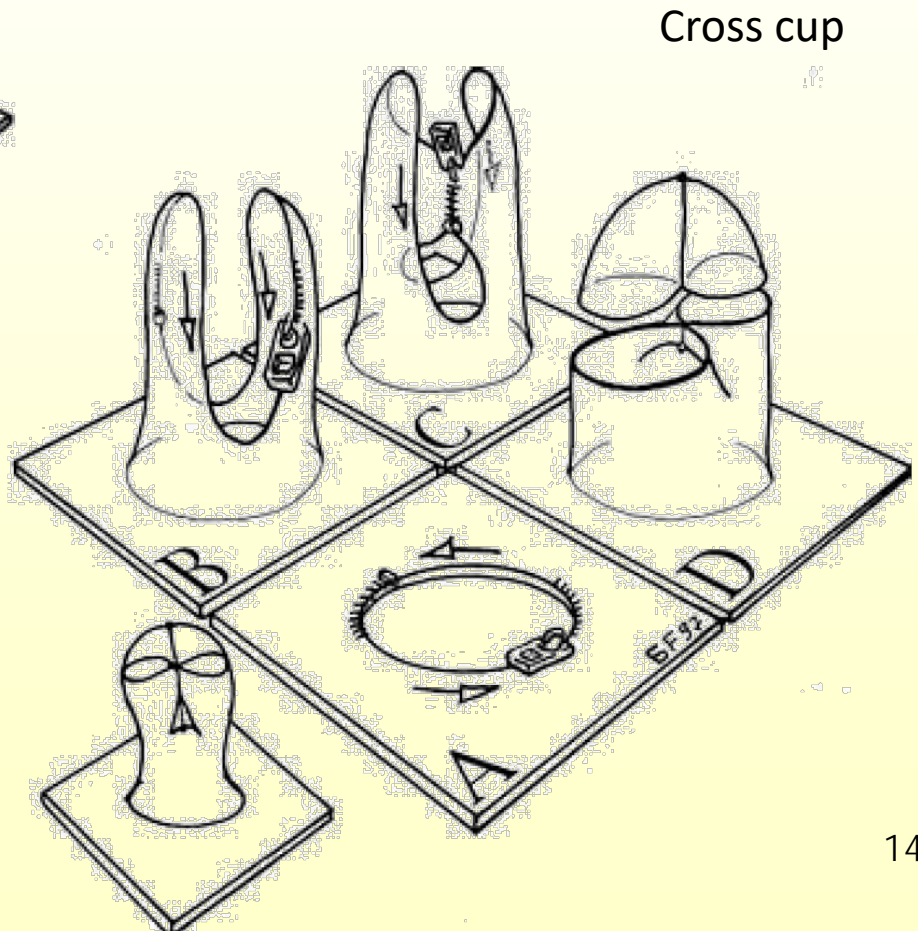


Cross cup

# J. Conway's ZIP Proof (Zero Irrelevancy Proof)



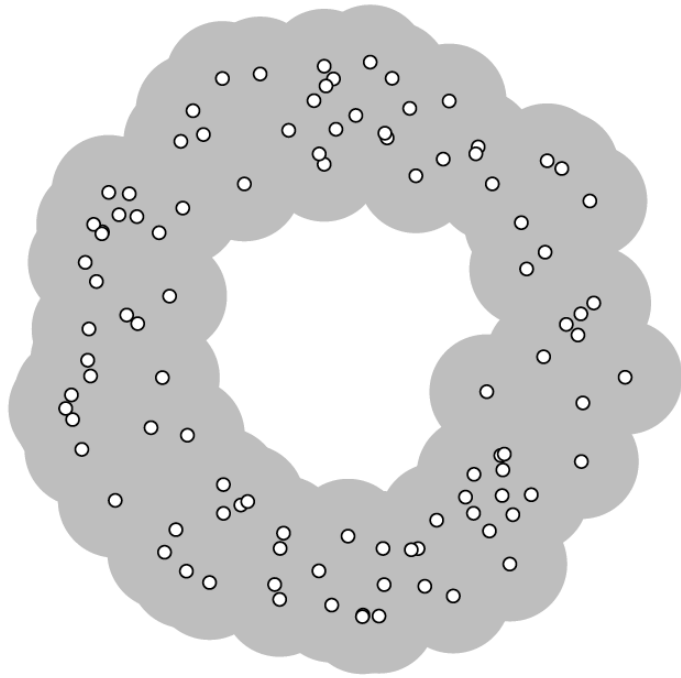
Cup



Cross cup

# Sampled Spaces

# Recovering “Shape” from Sampled Data



1. Set of points in  $\mathbb{R}^2$
2. Looks like an annulus.

What is this?

What does it look like ?

**Aim:** recover the topology of the underlying space from which the data was sampled



# Example: The Space of Natural Images

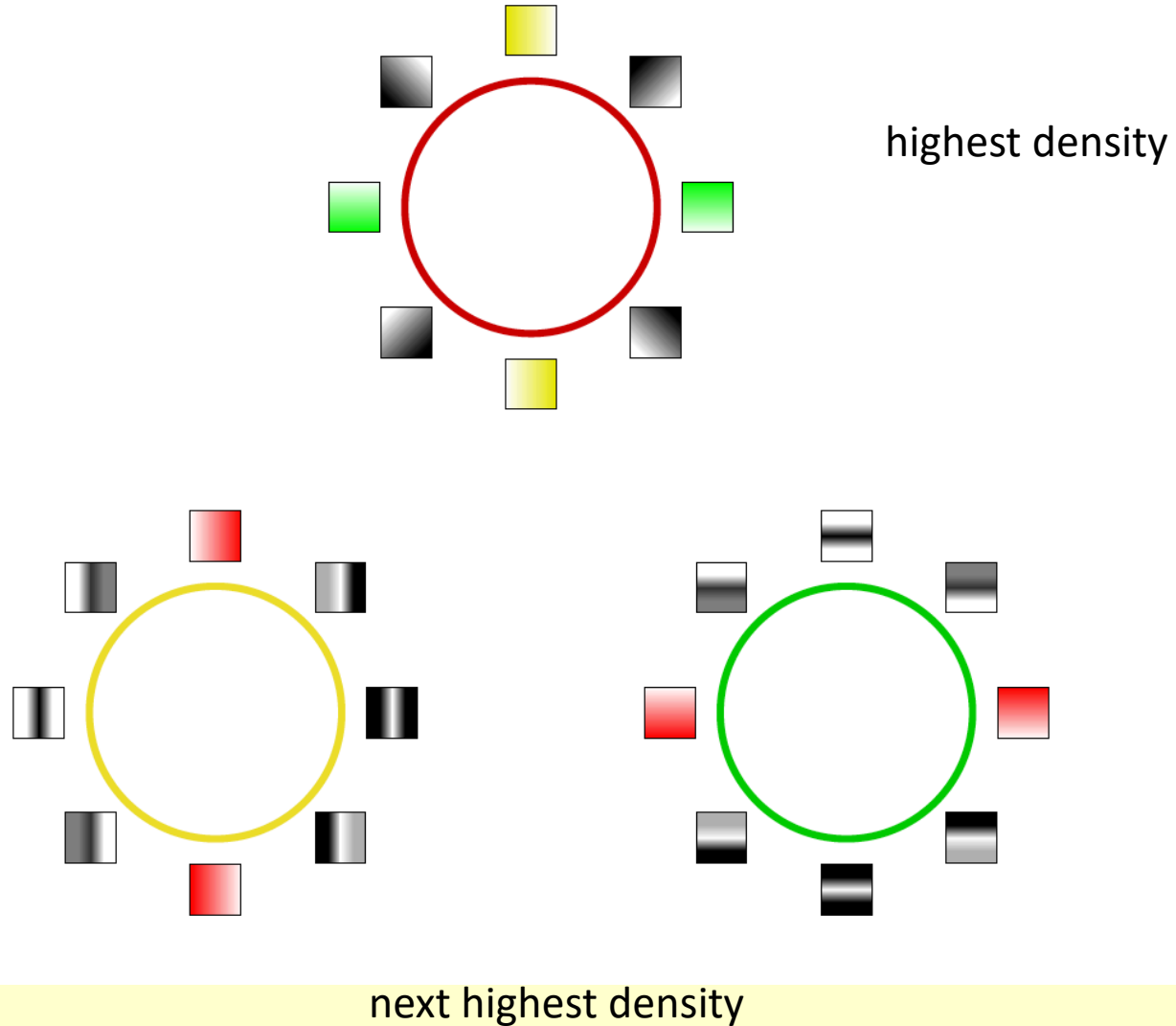
(Carlsson, Ishkanov, de Silva, Zomorodian IJCV 2008)

- ◆ Lee-Mumford-Pedersen investigated whether a statistically significant difference exists between natural and random images
- ◆ Natural images form a “subspace” of all images. Dimension of ambient space e.g.  $640 \times 480 = 307\,200$
- ◆ This space of natural images should have:
  - ◆ high dimension: there are many different images
  - ◆ even higher co-dimension: random images look nothing like natural ones
- ◆ Data is a collection of black-and-white images used in cognitive science research

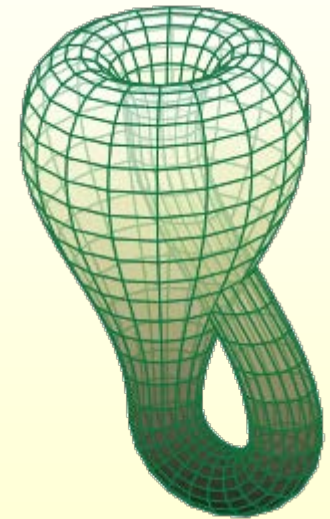
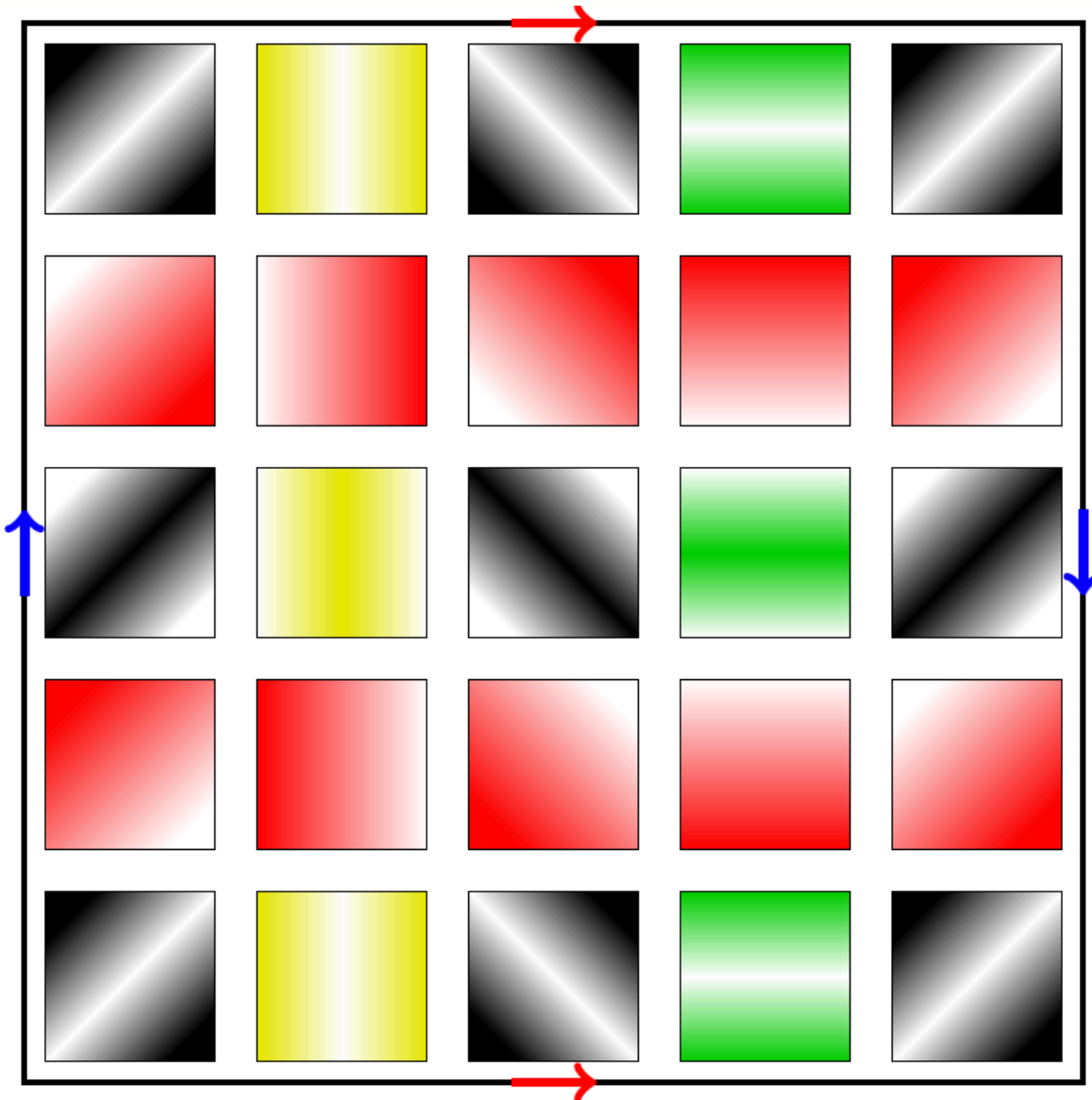
# Natural 3 x 3 Patches

- ◆ Instead of studying entire images, we consider the distribution of 3 x 3 pixel patches
- ◆ Most of these are roughly constant in natural images -- they drown out structure
- ◆ L.M.P. chose 8,500,000 patches with high contrast
- ◆ Each 3 x 3-patch is considered a vector in  $\mathbb{R}^9$
- ◆ Normalize brightness:  $\mathbb{R}^9 \rightarrow \mathbb{R}^8$
- ◆ Normalize contrast:  $\mathbb{R}^8 \rightarrow S^7$

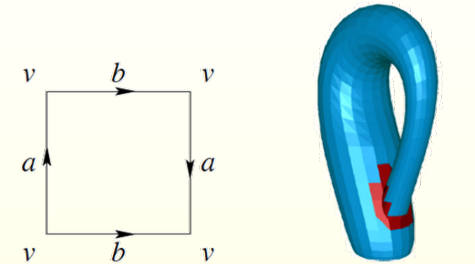
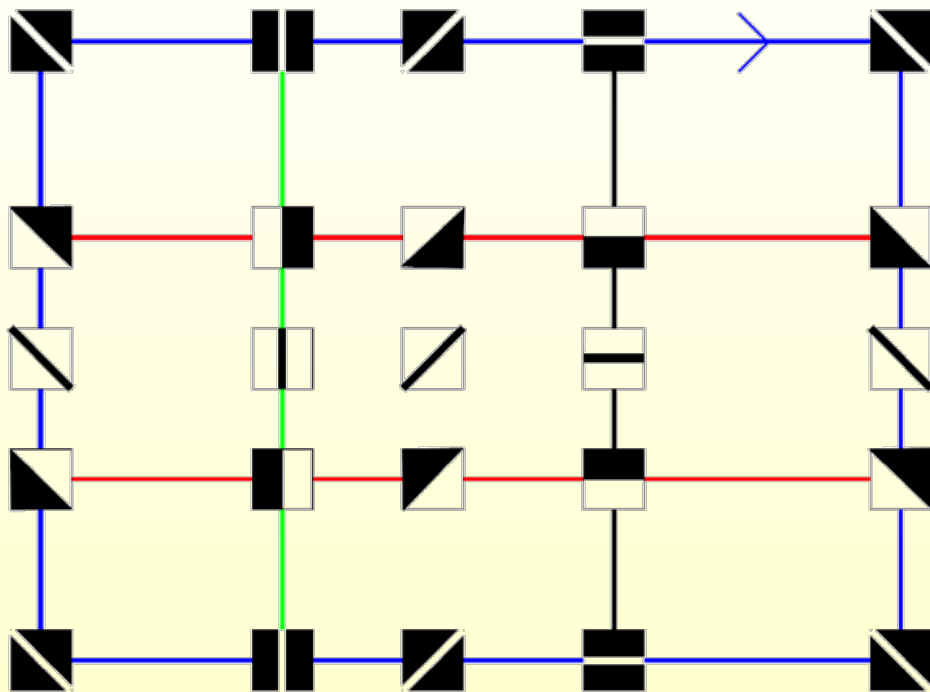
# High-Density Areas



# Klein Bottle of Pixel Patches

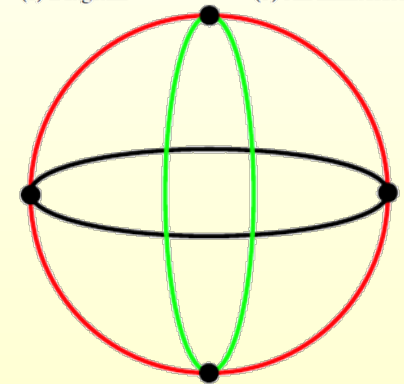


# Klein Bottle Structure



(a) Diagram

(b) An immersion



$$(\beta_1 = 5)$$

(source: [Carlsson, Ishkhanov, de Silva, Zomorodian 2008])

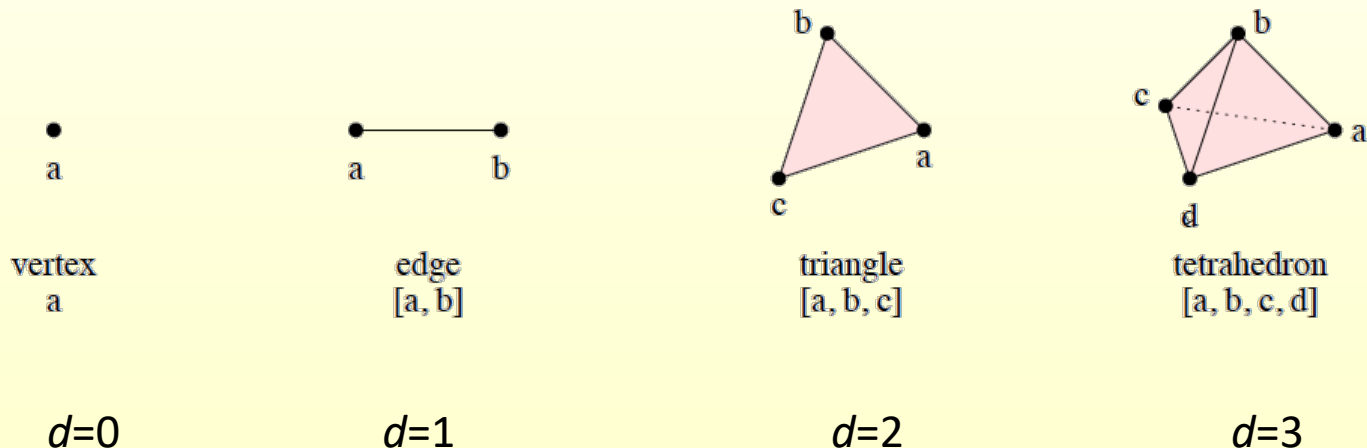
# Applications of the Analysis

- ◆ An efficient way to parametrize image patches
- ◆ **Image compression:** a 3 x 3-cluster may be described using 4 values
  - ◆ Position of its projection onto the Klein bottle
  - ◆ Original brightness
  - ◆ Original contrast
- ◆ **Texture analysis:** textures yield distributions of occurring patches on the Klein bottle. Rotating the texture corresponds to translating the distribution.

# Simplicial Complexes: Combinatorial Topology

# A Simplex

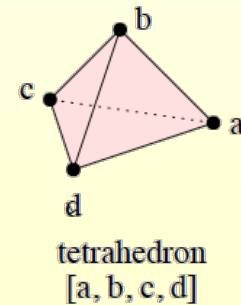
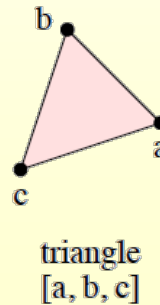
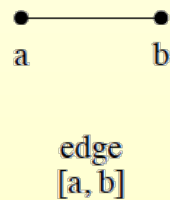
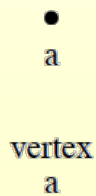
- A  **$k$ -simplex** is the convex hull of  $k + 1$  affinely independent points  $S = \{v_0, v_1, \dots, v_k\}$ . The points in  $S$  are the **vertices** of the simplex.
- A  $k$ -simplex is a  $k$ -dimensional subspace of  $\mathbb{R}^d$ ,  $\dim \sigma = k$ .





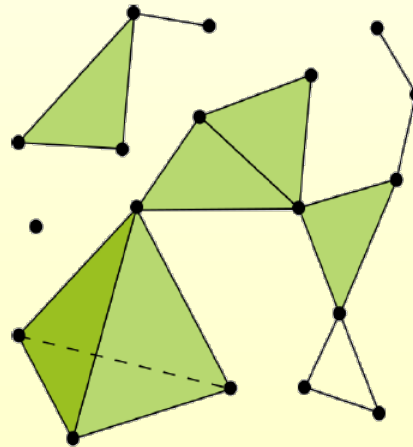
# Faces / Subsimplices

- $\sigma$ : a  $k$ -simplex defined by  $S = \{v_0, v_1, \dots, v_k\}$ .
- $\tau$  defined by  $T \subseteq S$  is a **face** of  $\sigma$
- $\sigma$  is its **coface**.
- $\sigma \geq \tau$  and  $\tau \leq \sigma$ .
- $\sigma \leq \sigma$  and  $\sigma \geq \sigma$ .

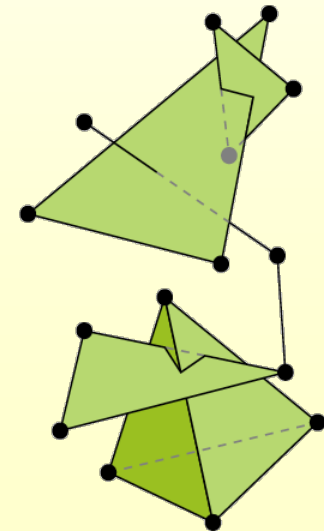


# Simplicial Complexes

- A **simplicial complex**  $K$  is a finite set of simplices such that
  1.  $\sigma \in K, \tau \leq \sigma \Rightarrow \tau \in K$ ,
  2.  $\sigma, \sigma' \in K \Rightarrow \sigma \cap \sigma' \leq \sigma, \sigma'$  or  $\sigma \cap \sigma' = \emptyset$ .
- The **dimension** of  $K$  is  $\dim K = \max\{\dim \sigma \mid \sigma \in K\}$ .
- The **vertices** of  $K$  are the zero-simplices in  $K$ .
- A simplex is **principal** if it has no proper coface in  $K$ .

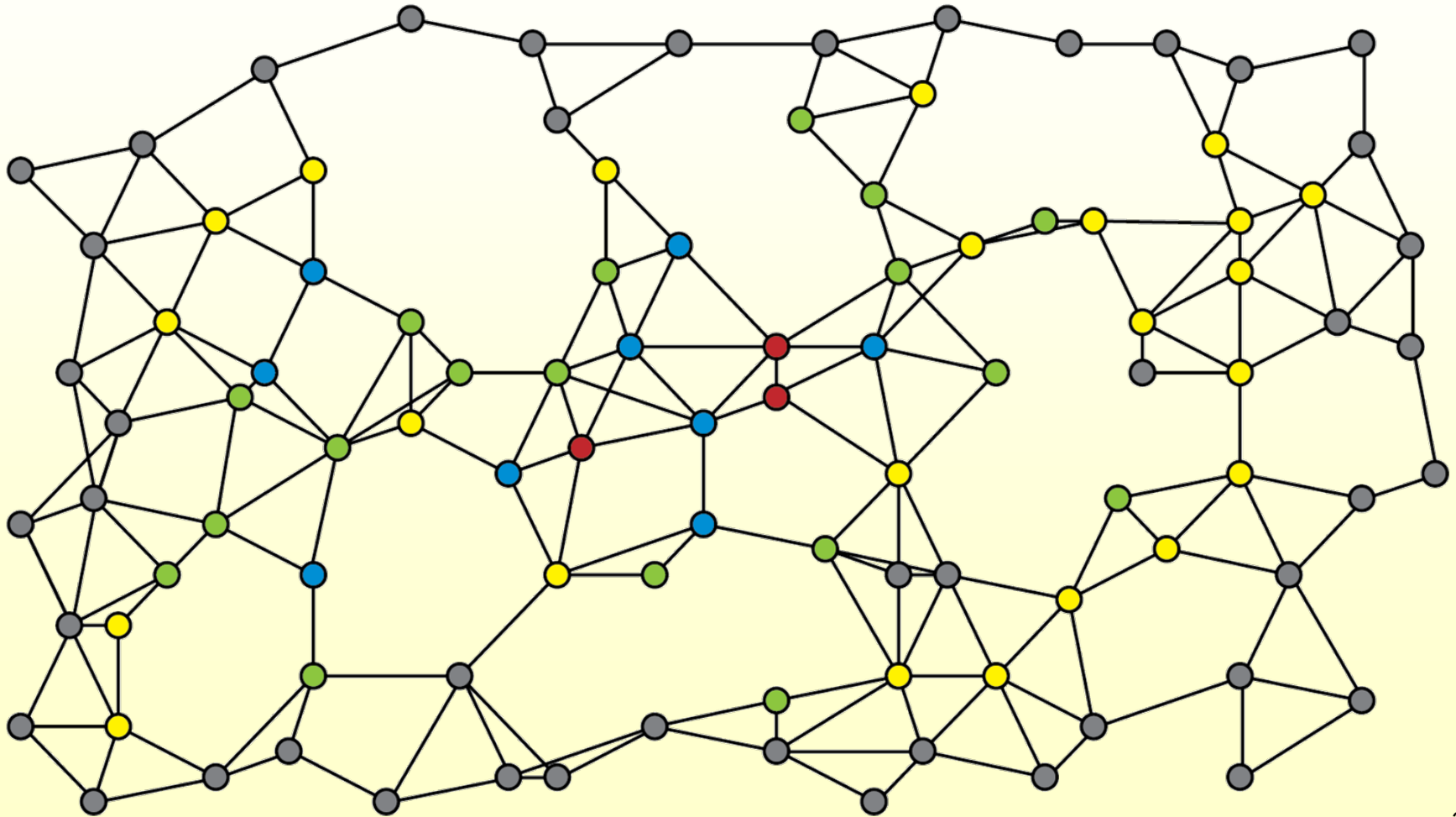


(left) an example

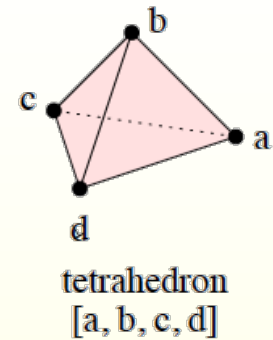
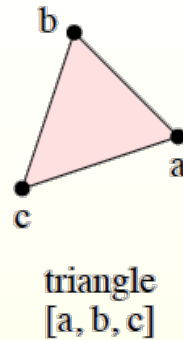
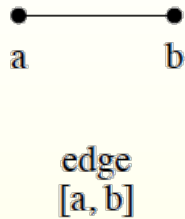
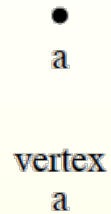


(right) a non example

# Good Models for Sensor Networks



# Size of a Simplex



- $\emptyset$  is the  $(-1)$ -simplex.
- A  $k$ -simplex has  $\binom{k+1}{l+1}$  faces of dimension  $l$
- Total size is:

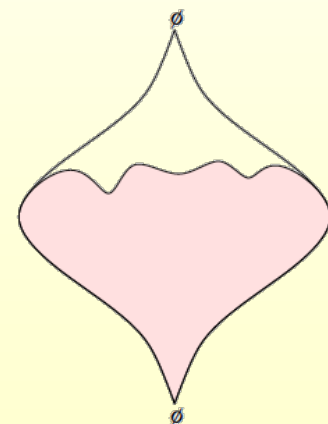
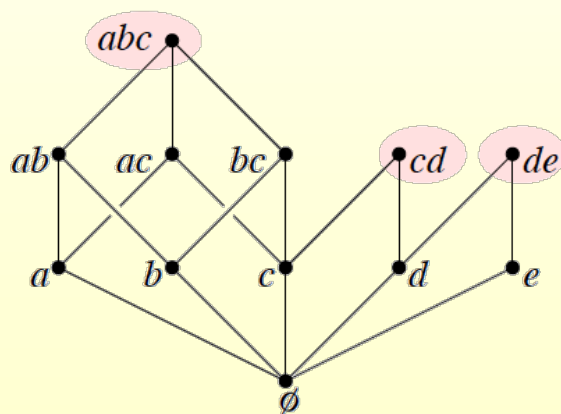
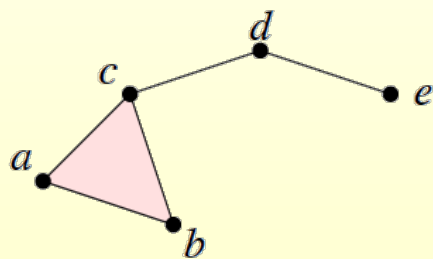
$$\sum_{l=-1}^k \binom{k+1}{l+1} = 2^{k+1}$$

$k/l$	0	1	2	3
0	1	0	0	0
1	2	1	0	0
2	3	3	1	0
3	4	6	4	1
4	?	?	?	?

Binomial coefficients

# Abstract Simplicial Complexes

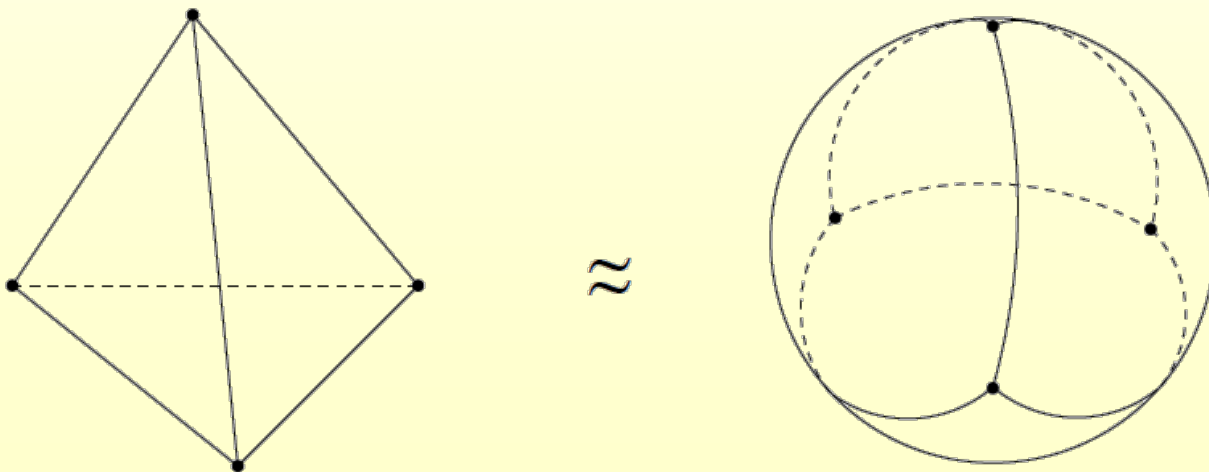
- An **abstract simplicial complex** is a set  $K$ , together with a collection  $\mathcal{S}$  of subsets of  $K$  called **(abstract) simplices** such that:
  1. For all  $v \in K$ ,  $\{v\} \in \mathcal{S}$ . We call the sets  $\{v\}$  the **vertices** of  $K$ .
  2. If  $\tau \subseteq \sigma \in \mathcal{S}$ , then  $\tau \in \mathcal{S}$ .
- We call  $\mathcal{S}$  the complex.



Natural partial order structure

# Continuous to Discrete Link: Triangulations

- The **underlying space**  $|K|$  of a simplicial complex  $K$  is  $|K| = \cup_{\sigma \in K} \sigma$ .
- $|K|$  is a topological space.
- A **triangulation** of a topological space  $\mathbb{X}$  is a simplicial complex  $K$  such that  $|K| \approx \mathbb{X}$ .



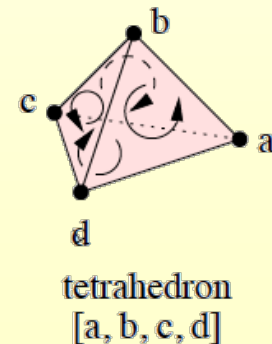
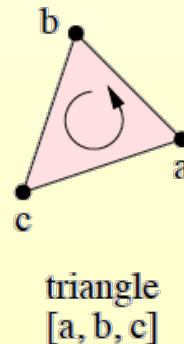
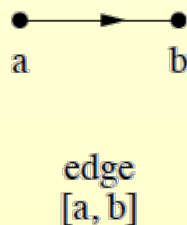
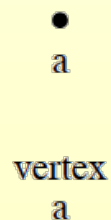
# Orientability

- An **orientation** of a  $k$ -simplex  $\sigma \in K$ ,  $\sigma = \{v_0, v_1, \dots, v_k\}$ ,  $v_i \in K$  is an equivalence class of orderings of the vertices of  $\sigma$ , where

$$(v_0, v_1, \dots, v_k) \sim (v_{\tau(0)}, v_{\tau(1)}, \dots, v_{\tau(k)})$$

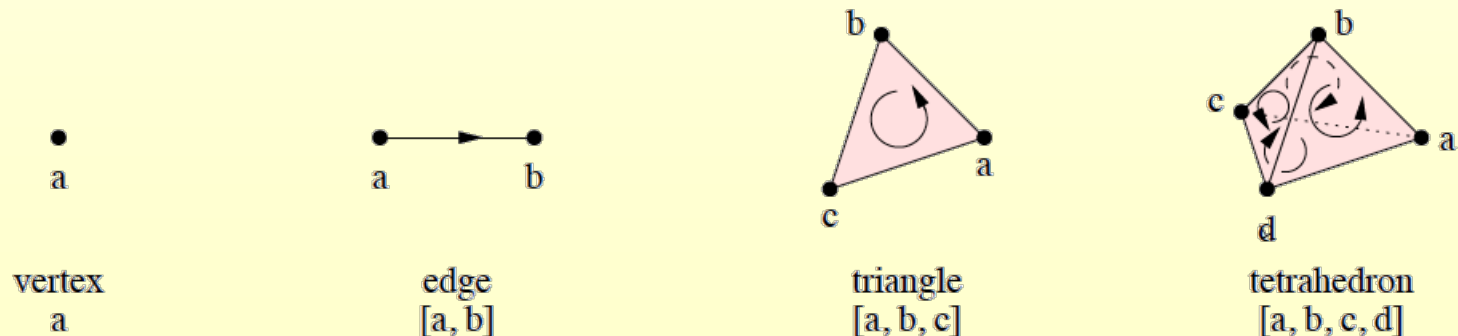
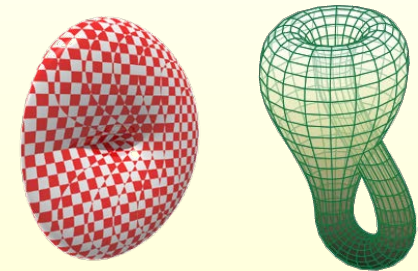
are equivalent orderings if the parity of the permutation  $\tau$  is even.

- We denote an **oriented simplex**, a simplex with an equivalence class of orderings, by  $[\sigma]$ .



# Orientability

- Two  $k$ -simplices sharing a  $(k - 1)$ -face  $\sigma$  are **consistently oriented** if they induce different orientations on  $\sigma$ .
- A triangulable  $d$ -manifold is **orientable** if all  $d$ -simplices can be oriented consistently.
- Otherwise, the  $d$ -manifold is **non-orientable**





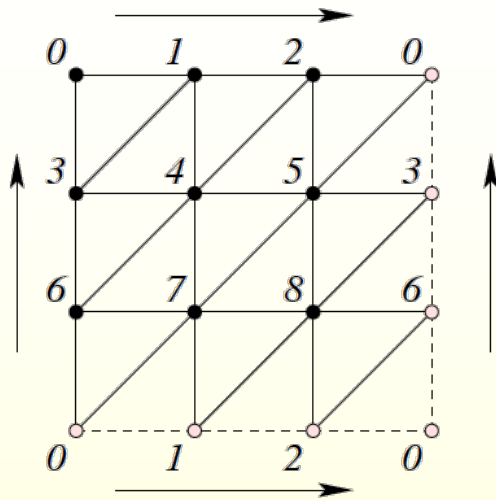
# Euler Characteristic: A Topological Invariant

- $K$  a simplicial complex with  $s_k$   $k$ -simplices.
- The **Euler characteristic**  $\chi(K)$  is

$$\chi(K) = \sum_{i=0}^{\dim K} (-1)^i s_i = \sum_{\sigma \in K - \{\emptyset\}} (-1)^{\dim \sigma}.$$

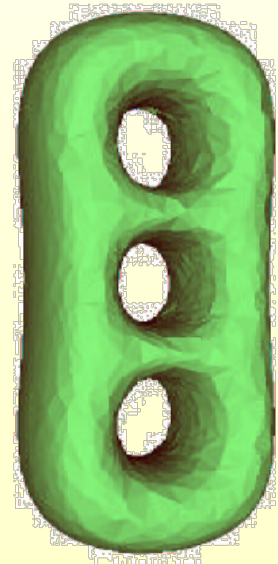
- $v - e + f = 1$  (Graph Theory)
- Invariant for  $|K|$
- **Any** triangulation gives the same answer!
- Intrinsic property

# More on Euler



2-Manifold	$\chi$
Sphere $S^2$	2
Torus $T^2$	0
Klein bottle $\mathbb{K}^2$	0
Projective plane $\mathbb{RP}^2$	1

- (Theorem) For compact surfaces  $M_1, M_2$ ,  
 $\chi(M_1 \# M_2) = \chi(M_1) + \chi(M_2) - 2$ .
- $\chi(gT^2) = 2 - 2g$
- $\chi(g\mathbb{RP}^2) = 2 - g$
- The connected sum of  $g$  tori is called a surface with **genus**  $g$ .



# denotes connected sum

# Topological Classification via Invariants

- (Theorem) Closed compact surfaces  $M_1$  and  $M_2$  are homeomorphic,  $M_1 \approx M_2$  iff
  1.  $\chi(M_1) = \chi(M_2)$  and
  2. either both surfaces are orientable or both are non-orientable.
- “iff” so full answer. We’re done!
- Higher dimensions?

This is what classical topology tries to do

# Algebraic Structures: Groups, Vector Spaces

# Groups

- A **group**  $\langle G, * \rangle$  is a set  $G$ , together with a binary operation  $*$  on  $G$ , such that the following axioms are satisfied:
  - (a)  $*$  is associative.
  - (b)  $G$  has an **identity**  $e$  element for  $*$  such that  $e * x = x * e = x$  for all  $x \in G$ .
  - (c) any element  $a$  has an **inverse**  $a'$  with respect to the operation  $*$ , i.e.  $\forall a \in G, \exists a' \in G$  such that  $a' * a = a * a' = e$ .
- If  $G$  is finite, the **order** of  $G$  is  $|G|$ .
- We often omit the operation and refer to  $G$  as the group.
- $\langle \mathbb{Z}, + \rangle, \langle \mathbb{R}, \cdot \rangle, \langle \mathbb{R}, + \rangle$ , are all groups.
- A group  $G$  is **abelian** if its binary operation  $*$  is commutative.

# Subgroups

- Let  $\langle G, * \rangle$  be a group and  $S \subseteq G$ . If  $S$  is closed under  $*$ , then  $*$  is the **induced operation on  $S$  from  $G$** .
- A subset  $H \subseteq G$  of group  $\langle G, * \rangle$  is a **subgroup of  $G$**  if  $H$  is a group and is closed under  $*$ . The subgroup consisting of the identity element of  $G$ ,  $\{e\}$  is the **trivial subgroup** of  $G$ . All other subgroups are **nontrivial**.
- (Theorem)  $H \subseteq G$  of a group  $\langle G, * \rangle$  is a subgroup of  $G$  iff:
  1.  $H$  is closed under  $*$ ,
  2. the identity  $e$  of  $G$  is in  $H$ ,
  3. for all  $a \in H$ ,  $a^{-1} \in H$ .
- Example: subgroups of  $\mathbb{Z}_4$

# Cosets

- Let  $H$  be a subgroup of  $G$ . Let the relation  $\sim_L$  be defined on  $G$  by:  
 $a \sim_L b$  iff  $a^{-1}b \in H$ . Let  $\sim_R$  be defined by:  $a \sim_R b$  iff  $ab^{-1} \in H$ .  
Then  $\sim_L$  and  $\sim_R$  are both equivalence relations on  $G$ .
- Let  $H$  be a subgroup of group  $G$ . For  $a \in G$ , the subset  
 $aH = \{ah \mid h \in H\}$  of  $G$  is the **left coset** of  $H$  containing  $a$ , and  
 $Ha = \{ha \mid h \in H\}$  is the **right coset** of  $H$  containing  $a$ .
- If left and right cosets match, the subgroup is **normal**.
- All subgroups  $H$  of an abelian group  $G$  are normal, as  
 $ah = ha, \forall a \in G, h \in H$
- $\{0, 2\}$  is a subgroup of  $\mathbb{Z}_4$ . It is normal. The coset of 1 is  
 $1 + \{0, 2\} = \{1, 3\}$ . That's all folks!

# Factor / Quotient Groups

- Let  $H$  be a normal subgroup of group  $G$ .
- Left coset multiplication is well-defined by the equation
$$(aH)(bH) = (ab)H$$
- The cosets of  $H$  form a group  $G/H$  under left multiplication
- $G/H$  is the **factor group** (or **quotient group**) of  $G$  modulo  $H$ .
- The elements in the same coset of  $H$  are **congruent modulo  $H$** .



# Example

$\mathbb{Z}_6$	0	2	4	1	3	5
0	0	2	4	1	3	5
2	2	4	0	3	5	1
4	4	0	2	5	1	3
1	1	3	5	2	4	0
3	3	5	1	4	0	2
5	5	1	3	0	2	4

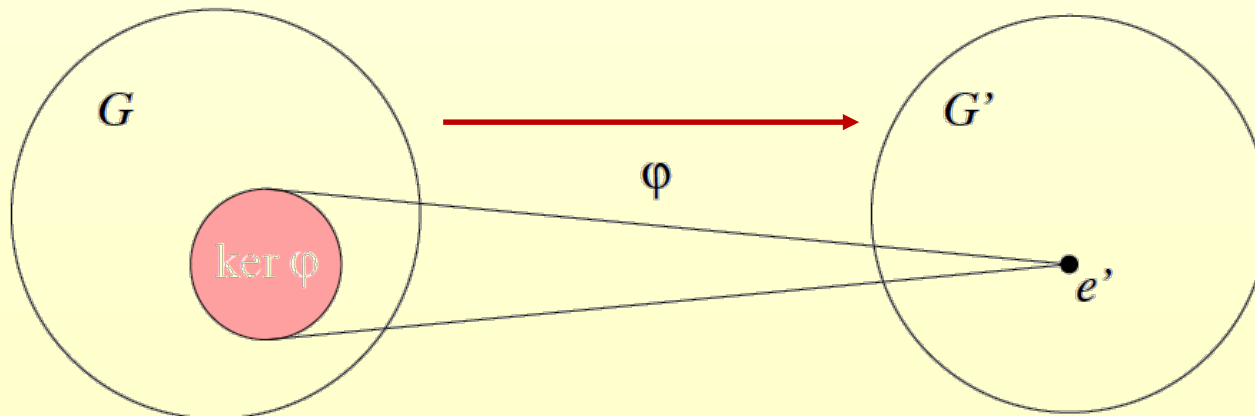
  

	*		

- $\{0, 2, 4\}$  is a normal subgroup
- Cosets  $\{0, 2, 4\}, \{1, 3, 5\}$
- $\mathbb{Z}_6 / \{0, 2, 4\} \cong \mathbb{Z}_2$

# Group Homomorphisms

- A map  $\varphi$  of a group  $G$  into a group  $G'$  is a *homomorphism* if  $\varphi(ab) = \varphi(a)\varphi(b)$  for all  $a, b \in G$ .
- If  $e$  is the identity in  $G$ , then  $\varphi(e)$  is the identity  $e'$  in  $G'$ .
- If  $a \in G$ , then  $\varphi(a^{-1}) = \varphi(a)^{-1}$ .
- If  $H$  is a subgroup of  $G$ , then  $\varphi(H)$  is a subgroup of  $G'$ .
- If  $K'$  is a subgroup of  $G'$ , then  $\varphi^{-1}(K')$  is a subgroup of  $G$ .
- The normal subgroup  $\ker \varphi = \varphi^{-1}(\{e'\}) \subseteq G$ , is the **kernel of  $\varphi$** .



# Decompositions for Finitely Generated Abelian Groups

- Let  $G_1, G_2, \dots, G_n$  be groups.
- The set is  $\prod_{i=1}^n G_i$  (Cartesian product)
- Binary operation:  
 $(a_1, a_2, \dots, a_n) \times (b_1, b_2, \dots, b_n) = (a_1 b_1, a_2 b_2, \dots, a_n b_n)$ .
- Then  $\langle \prod_{i=1}^n G_i, \times \rangle$  is a group.
- We call it the **direct product** of the groups  $G_i$ .
- Sometimes called **direct sum** with  $\oplus$ .

- (Theorem) Every finitely generated abelian group is isomorphic to product of cyclic groups of the form

$$\mathbb{Z}_{m_1} \times \mathbb{Z}_{m_2} \times \dots \times \mathbb{Z}_{m_r} \times \mathbb{Z} \times \mathbb{Z} \times \dots \times \mathbb{Z},$$

where  $m_i$  divides  $m_{i+1}$  for  $i = 1, \dots, r - 1$ .

- The direct product is unique: the number of factors of  $\mathbb{Z}$  is unique and the cyclic group orders  $m_i$  are unique.
- Free: basis, rank, vector space
- Torsion: module

# Homological Algebra: Functors and Categories

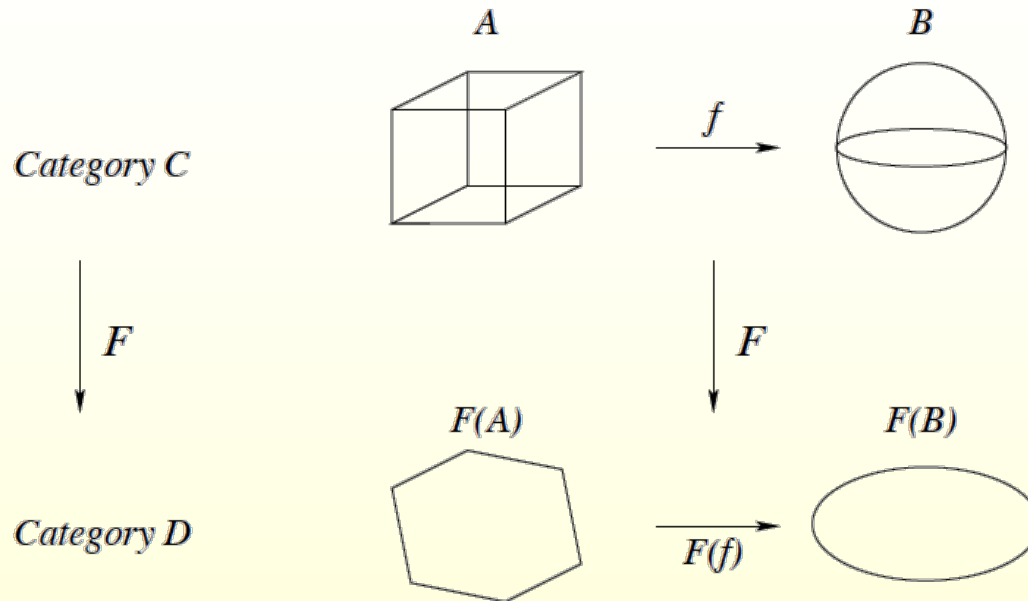
# Categories

- A collection  $\text{Ob}(\mathcal{C})$  of **objects**
- Sets  $\text{Mor}(X, Y)$  of **morphisms** for each pair  $X, Y \in \text{Ob}(\mathcal{C})$
- An identity morphism  $1 = 1_X \in \text{Mor}(X, X)$  for each  $X$ .
- a composition of morphisms function
  - $\circ : \text{Mor}(X, Y) \times \text{Mor}(Y, Z) \rightarrow \text{Mor}(X, Z)$  for each triple  $X, Y, Z \in \text{Ob}(\mathcal{C})$ , satisfying  $f \circ 1 = 1 \circ f = f$ , and  $(f \circ g) \circ h = f \circ (g \circ h)$ .
- A **category**  $\mathcal{C}$

# Example Categories

category	morphisms
sets	arbitrary functions
groups	homomorphisms
topological spaces	continuous maps
topological spaces	homotopy classes of maps

# Functors



- $X \in \mathcal{C}, F(X) \in \mathcal{D},$
- $f \in \text{Mor}(X, Y), F(f) \in \text{Mor}(F(X), F(Y))$
- $F(1) = 1$  and  $F(f \circ g) = F(f) \circ F(g)$
- $F$  is a **(covariant) functor**

# Functoriality

transformation of input  $\Rightarrow$  transformation of output  
Specifically, this is a commutative diagram:

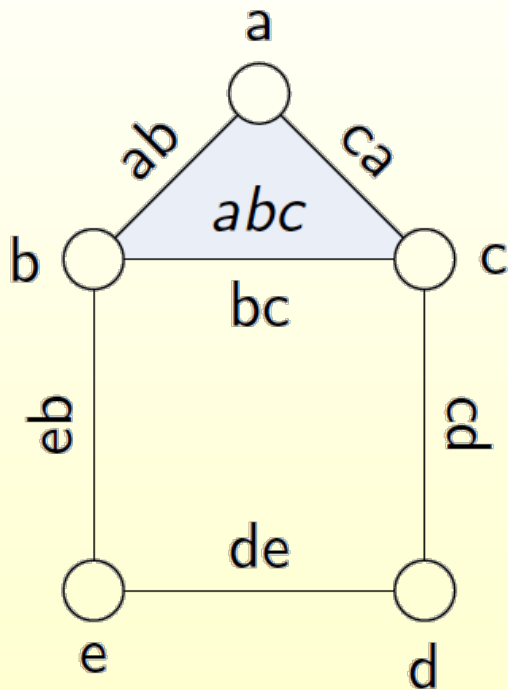
$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ H_* \downarrow & & \downarrow H_* \\ H_*(X) & \xrightarrow{H_*(f)} & H_*(Y) \end{array}$$

**Moral:** Invariants are not artifacts of arbitrary choices!



# Algebraic Topology: Homology

# Topology of Simplicial Complexes



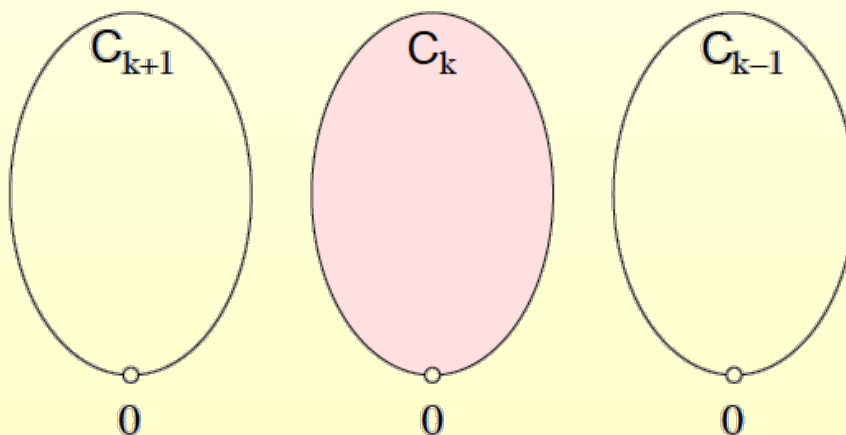
A simplicial complex is a collection of simplices

- ▶ Each simplex has a dimension.
- ▶ Collection is closed under subset relation.
- ▶ Simplices of dimension  $d$  have  $d + 1$  vertices
- ▶ Each simplex represented by an ordered list of vertices

# Chain Groups

Other coefficient fields/rings also OK

- Simplicial complex  $K$
- **$k$ -chain**:  $c = \sum_i n_i [\sigma_i]$ ,  $n_i \in \mathbb{Z}$ ,  $\sigma_i \in K$  (like a path)
- $[\sigma] = -[\tau]$  if  $\sigma = \tau$  and  $\sigma$  and  $\tau$  have different orientations.
- The  **$k$ th chain group  $\mathbf{C}_k$**  of  $K$  is the free abelian group on its set of oriented  $k$ -simplices
- $\text{rank } \mathbf{C}_k = ?$



# Boundary Operator

- The boundary operator  $\partial_k : \mathbf{C}_k \rightarrow \mathbf{C}_{k-1}$  is a homomorphism defined linearly on a chain  $c$  by its action on any simplex

$$\sigma = [v_0, v_1, \dots, v_k] \in c,$$

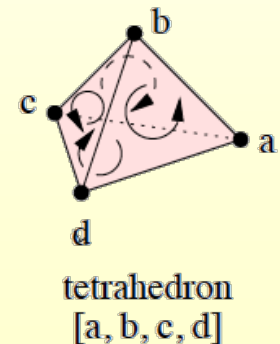
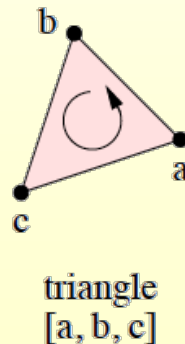
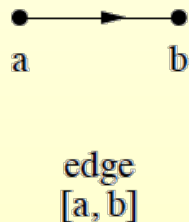
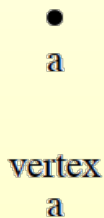
$$\partial_k \sigma = \sum_i (-1)^i [v_0, v_1, \dots, \hat{v}_i, \dots, v_k],$$

where  $\hat{v}_i$  indicates that  $v_i$  is deleted from the sequence.

- $\partial_1[a, b] = b - a.$
- $\partial_2[a, b, c] = [b, c] - [a, c] + [a, b] = [b, c] + [c, a] + [a, b].$
- $\partial_3[a, b, c, d] = [b, c, d] - [a, c, d] + [a, b, d] - [a, b, c].$

# Boundary Examples

- $\partial_1[a, b] = b - a.$
- $\partial_2[a, b, c] = [b, c] - [a, c] + [a, b] = [b, c] + [c, a] + [a, b].$
- $\partial_3[a, b, c, d] = [b, c, d] - [a, c, d] + [a, b, d] - [a, b, c].$
- $\partial_1\partial_2[a, b, c] = [c] - [b] - [c] + [a] + [b] - [a] = 0.$



# Boundary Theorem

- (Theorem)  $\partial_{k-1}\partial_k = 0$ , for all  $k$ .

- Proof:

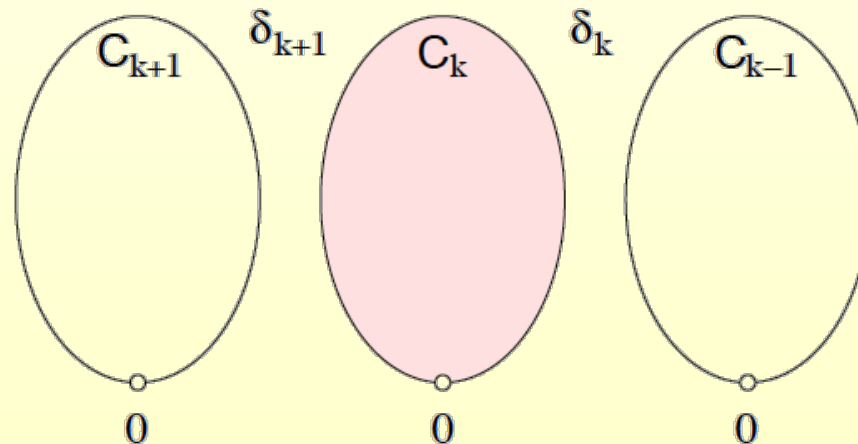
$$\begin{aligned}\partial_{k-1}\partial_k[v_0, v_1, \dots, v_k] &= \\ &= \partial_{k-1} \sum_i (-1)^i [v_0, v_1, \dots, \hat{v}_i, \dots, v_k] \\ &= \sum_{j < i} (-1)^i (-1)^j [v_0, \dots, \hat{v}_j, \dots, \hat{v}_i, \dots, v_k] \\ &\quad + \sum_{j > i} (-1)^i (-1)^{j-1} [v_0, \dots, \hat{v}_i, \dots, \hat{v}_j, \dots, v_k] \\ &= 0,\end{aligned}$$

as switching  $i$  and  $j$  in the second sum negates the first sum.

# Chain Complex

- The boundary operator connects the chain groups into a **chain complex**  $C_*$ :

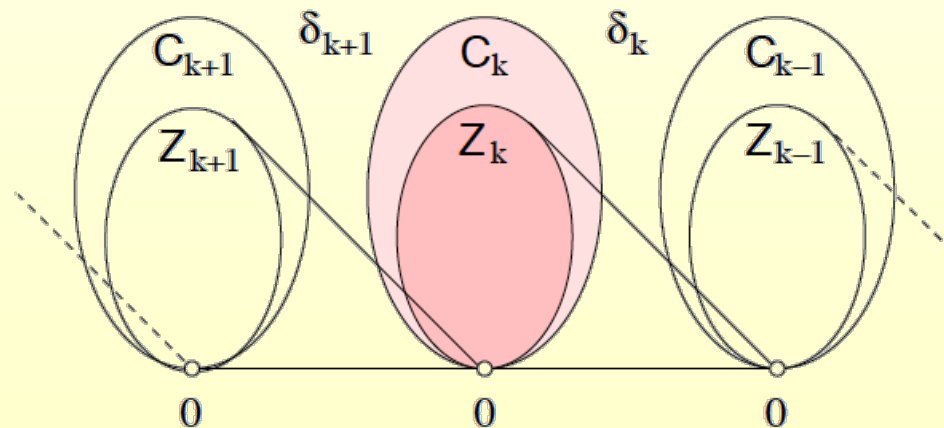
$$\dots \rightarrow C_{k+1} \xrightarrow{\partial_{k+1}} C_k \xrightarrow{\partial_k} C_{k-1} \rightarrow \dots$$



# Cycle Group

- Let  $c$  be a  $k$ -chain
- If it has no boundary, it is a  $k$ -cycle (zycle?)
- $\partial_k c = \emptyset$ , so  $c \in \ker \partial_k$
- The  $k$ th cycle group is

$$Z_k = \ker \partial_k = \{c \in C_k \mid \partial_k c = \emptyset\}.$$

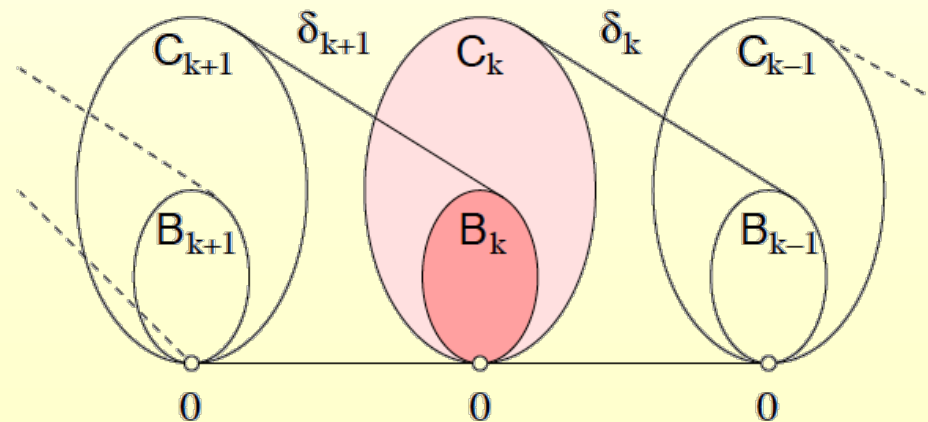




# Boundary Group

- Let  $b$  be a  $k$ -chain
- If  $b$  is a boundary of something, it is a  $k$ -boundary.
- The  $k$ th boundary group is

$$\mathbf{B}_k = \text{im } \partial_{k+1} = \{c \in \mathbf{C}_k \mid \exists d \in \mathbf{C}_{k+1} : c = \partial_{k+1}d\}.$$



# Boundaries are Cycles!

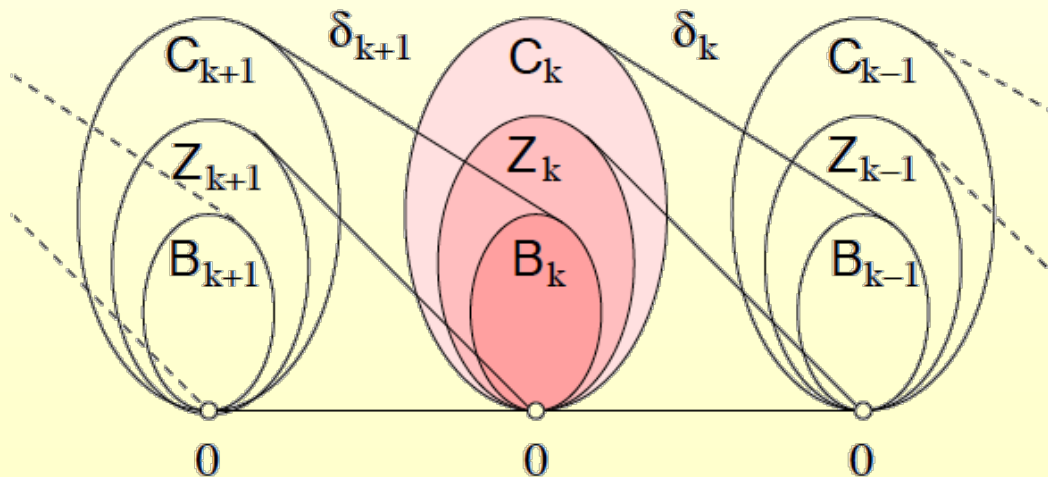
- Let  $b$  be a  $k$ -boundary.
- Then,  $\exists c \in \mathbf{C}_{k+1}$ , such that  $b = \partial_{k+1}c$ .
- What is the boundary of  $b$ ?

$$\partial_k b = \partial_k \partial_{k+1} c = \emptyset,$$

- $\mathbf{B}_k \subseteq \mathbf{Z}_k \subseteq \mathbf{C}_k$

by the boundary theorem.

- That is, every boundary is a cycle!

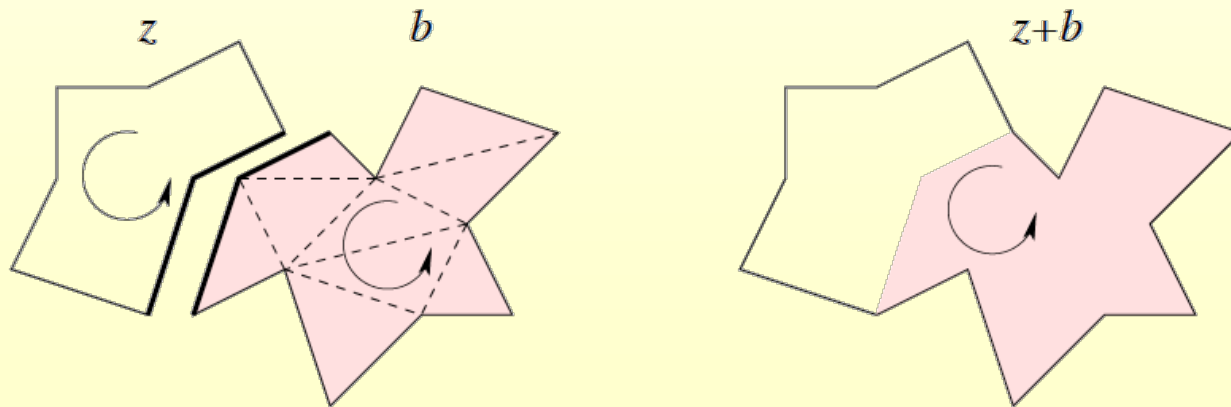


Nesting behavior

# Equivalent Cycles

- $z$  is a  $k$ -cycle
- $b$  is a  $k$ -boundary
- We would like to have  $z + b$  be equivalent to  $z$
- That is, if  $z_1 - z_2 = b$  where  $b$  is a boundary, then  $z_1 \sim z_2$
- Any boundary would do!

Cosets!

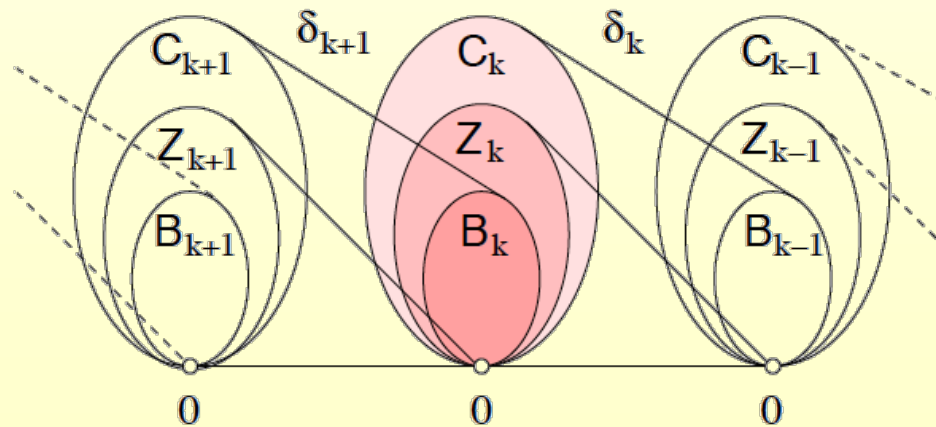


# Simplicial Homology

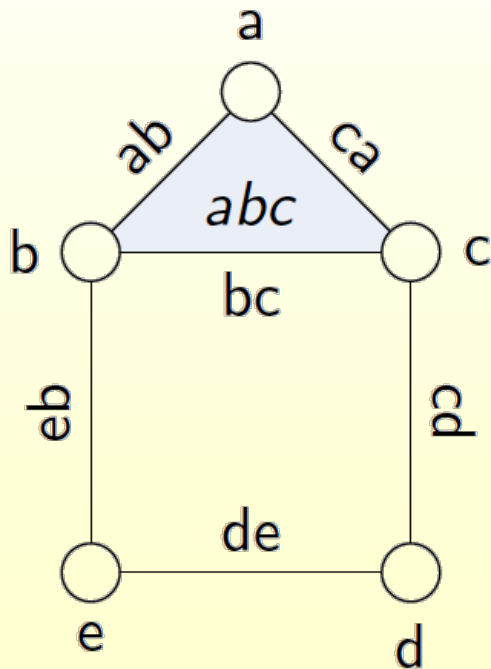
- The  $k$ th homology group is

$$H_k = Z_k / B_k = \ker \partial_k / \text{im } \partial_{k+1}.$$

- If  $z_1 = z_2 + B_k$ ,  $z_1, z_2 \in Z_k$ , we say  $z_1$  and  $z_2$  are **homologous**
- $z_1 \sim z_2$ .



# To Repeat



In other words..

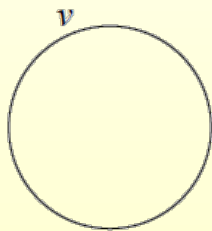
- ▶ The kernel (null space) of  $\partial_k$  is the vector space of cycles in dimension  $k$ .
- ▶ The image of  $\partial_k$  is the subspace of boundary cycles in dimension  $k - 1$ .

Homology of a space  $X$  is the quotient:

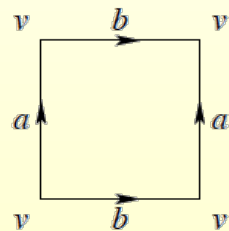
$$H_k(X) = \ker(\partial_k) / \text{im}(\partial_{k+1})$$

# Homology of 2-Manifolds

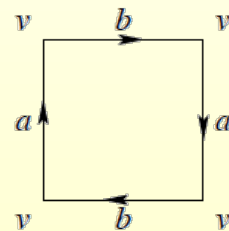
2-manifold	$H_0$	$H_1$	$H_2$
sphere	$\mathbb{Z}$	$\{0\}$	$\mathbb{Z}$
torus	$\mathbb{Z}$	$\mathbb{Z} \times \mathbb{Z}$	$\mathbb{Z}$
projective plane	$\mathbb{Z}$	$\mathbb{Z}_2$	$\{0\}$
Klein bottle	$\mathbb{Z}$	$\mathbb{Z} \times \mathbb{Z}_2$	$\{0\}$



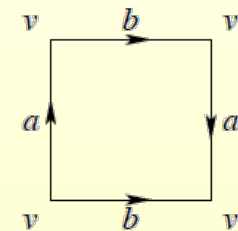
(a) Sphere



(b) Torus



(c) Projective plane



(d) Klein bottle

# Homology Groups

- Homology groups are finitely generated abelian.
- (Theorem) Every finitely generated abelian group is isomorphic to product of cyclic groups of the form

$$\mathbb{Z}_{m_1} \times \mathbb{Z}_{m_2} \times \dots \times \mathbb{Z}_{m_r} \times \mathbb{Z} \times \mathbb{Z} \times \dots \times \mathbb{Z},$$

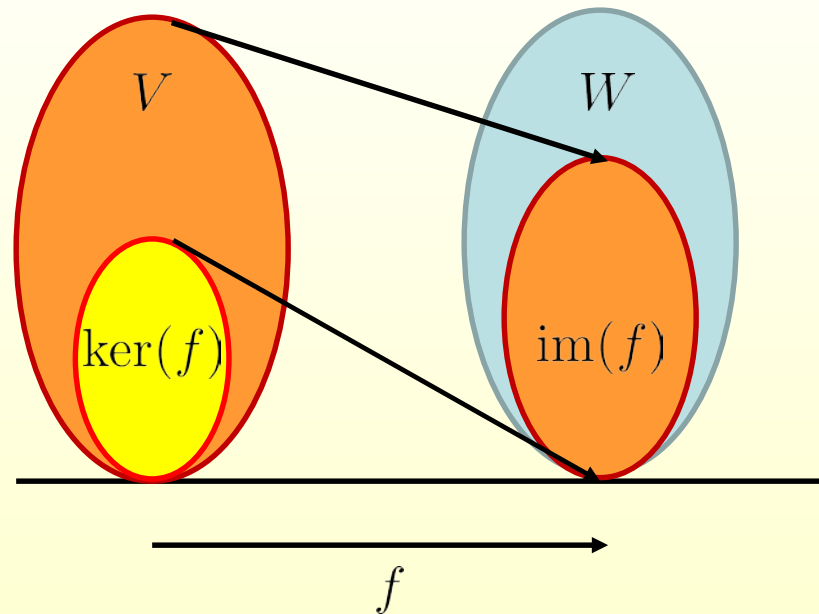
- The  $k$ th Betti number  $\beta_k$  of a simplicial complex  $K$  is  $\beta_k = \beta(\mathbf{H}_k)$ , the rank of the free part of  $\mathbf{H}_k$ .
- Torsion coefficients
  - Alexander Duality:
    - $\beta_0$  measures the number of components of the complex.
    - $\beta_1$  is the rank of a basis for the tunnels.
    - $\beta_2$  counts the number of voids in the complex.

# Invariance of Homology Groups

- (Hauptvermutung) Any two triangulations of a topological space have a common refinement (Poincaré 1904)
  - True for polyhedra of dimension  $\leq 2$  (Papakyriakopoulos 1943)
  - True for 3-manifolds (Moïse 1953)
  - False in dimensions  $\geq 6$  (Milnor 1961)
  - False for manifolds of dimension  $\geq 5$  (Kirby and Siebenmann 1969)
- Singular homology



# In Vector Spaces



$$V \approx \ker(f) \oplus \text{im}(f)$$

# Euler Revisited

- Let  $K$  be a simplicial complex and  $s_i = |\{\sigma \in K \mid \dim \sigma = i\}|$ . The Euler characteristic  $\chi(K)$  is

$$\chi(K) = \sum_{i=0}^{\dim K} (-1)^i s_i = \sum_{\sigma \in K - \{\emptyset\}} (-1)^{\dim \sigma}.$$

- We have new language!
- Let  $\mathbf{C}_*$  be the chain complex on  $K$
- $\text{rank}(\mathbf{C}_i) = |\{\sigma \in K \mid \dim \sigma = i\}|$  ( $= n_i = z_i + b_{i-1}$ )
- $\chi(K) = \chi(\mathbf{C}_*) = \sum_i (-1)^i \text{rank}(\mathbf{C}_i)$ .

$$\sum_i (-1)^i (z_i + b_{i-1}) = \sum_i (-1)^i (z_i - b_i)$$

# Euler - Poincaré

- Homology functors  $H_*$
- $H_*(C_*)$  is a chain complex:

$$\dots \rightarrow H_{k+1} \xrightarrow{\partial_{k+1}} H_k \xrightarrow{\partial_k} H_{k-1} \rightarrow \dots$$

- What is its Euler characteristic?
- (Theorem)  $\chi(K) = \chi(C_*) = \chi(H_*(C_*))$ .

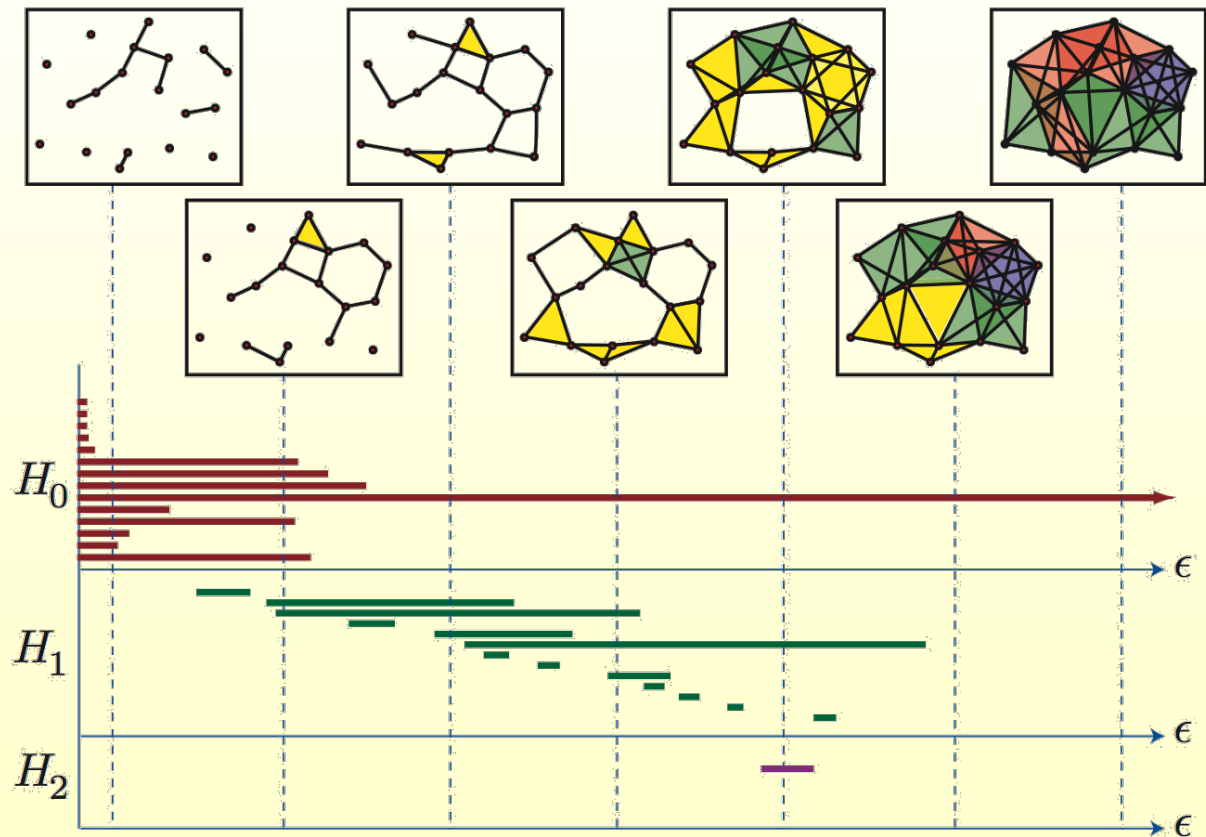
- $\sum_i (-1)^i s_i = \sum_i (-1)^i \text{rank}(H_i) = \sum_i (-1)^i \beta_i$

- Sphere:  $2 = 1 - 0 + 1$

- Torus:  $0 = 1 - 2 + 1$   $\sum_i (-1)^i (z_i + b_{i-1}) = \sum_i (-1)^i (z_i - b_i)$

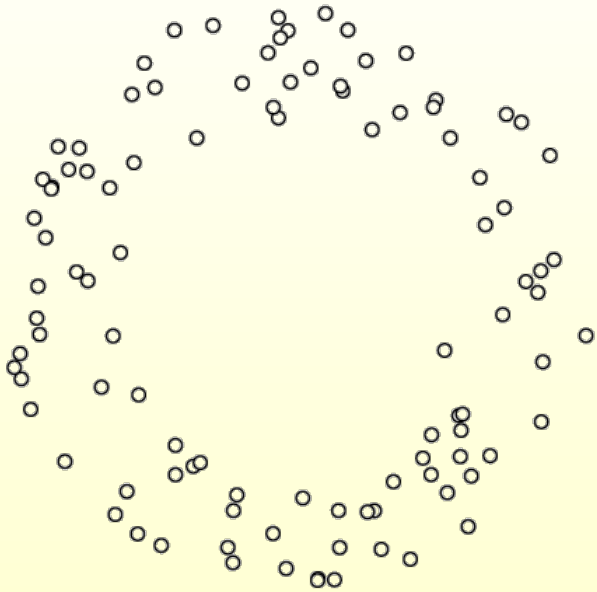
# Persistent Homology

# Persistent Homology



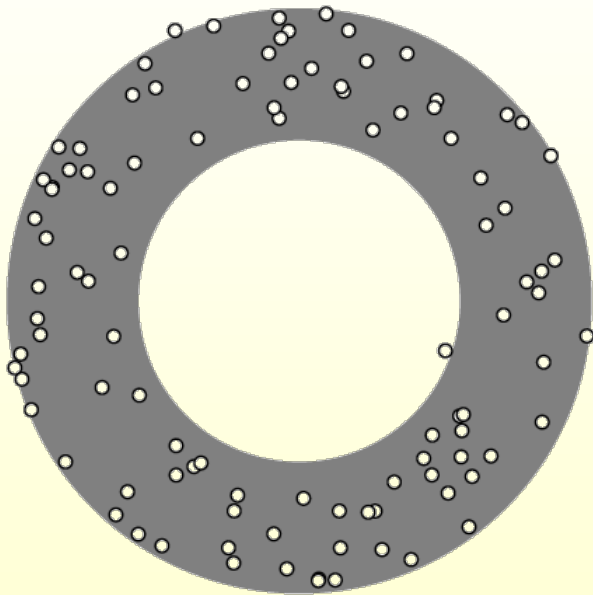
Slides ack: Afra  
Zomorodian, Ryan Lewis,  
Fred Chazal, Robert Ghrist

# Sampled Data Has “Shape”



2-dimensional  
Approximates annulus

# Sampled Data Has “Shape”



2-dimensional

Approximates annulus

Topological features of annulus:

1 component ( $\beta_0 = 1$ )

1 loop ( $\beta_1 = 1$ )

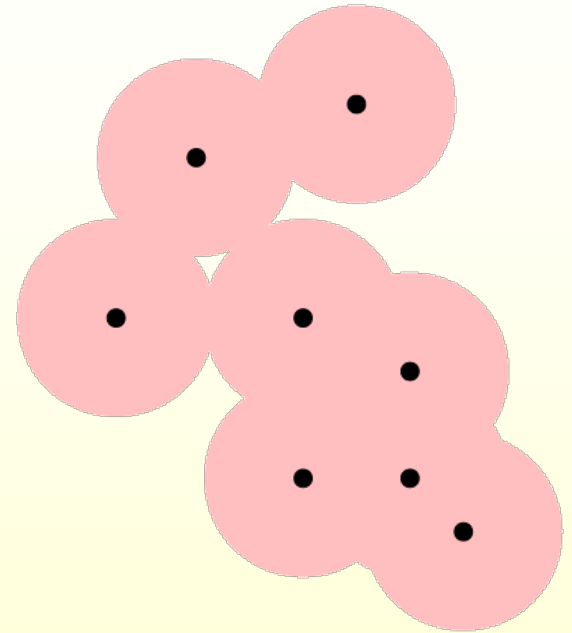
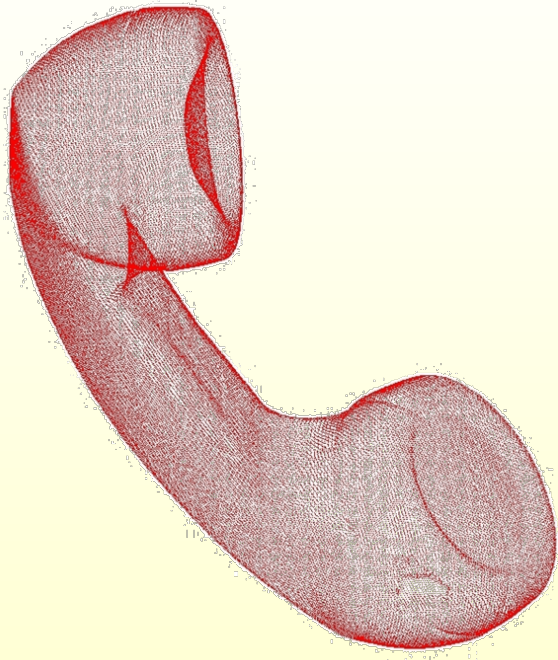
**Goal:** Recover *topology* of annulus from point cloud

We do so by building various complexes on the point cloud

# Complexes on Point Clouds



# $\epsilon$ -Balls



- $\epsilon$ -ball:  $B_\epsilon(x) = \{y \mid d(x, y) < \epsilon\}$ .
- Open sets and topology
- Manifold is  $\tilde{M} = \bigcup_{m_i \in M} B_\epsilon(m_i)$

# A Model Space

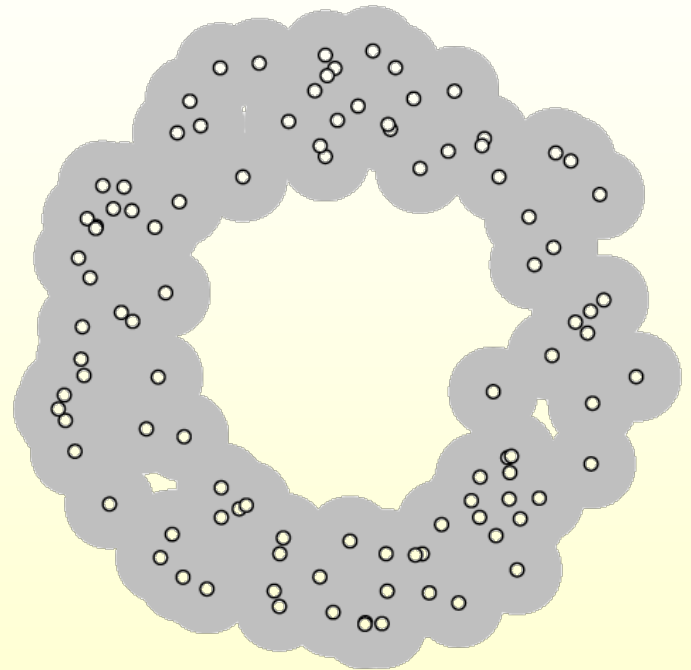
For a dataset  $X$  we study the topology of the *union of balls*

$$M_\epsilon = \bigcup_{x \in X} B_\epsilon(x)$$

**Two Issues:**

**Scale:** No natural choice of  $\epsilon$ !

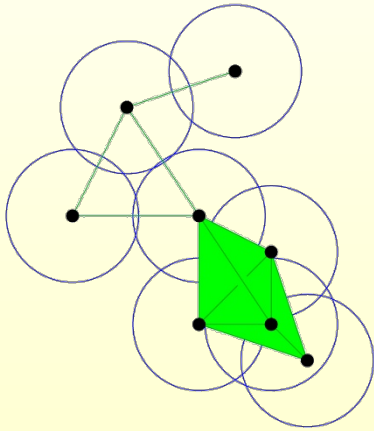
**Conception:** How to encode  $M_\epsilon$  on computer?



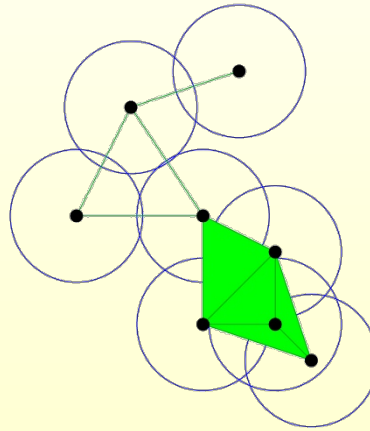
# Complex Zoo

Must choose which simplices to introduce

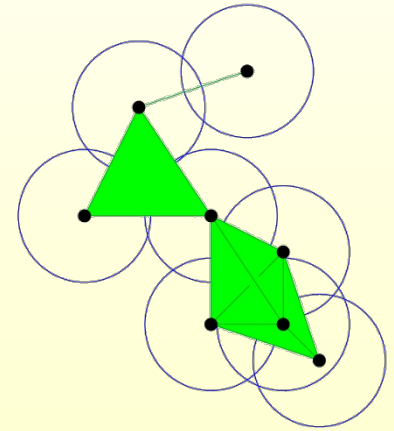
Čech



Alpha

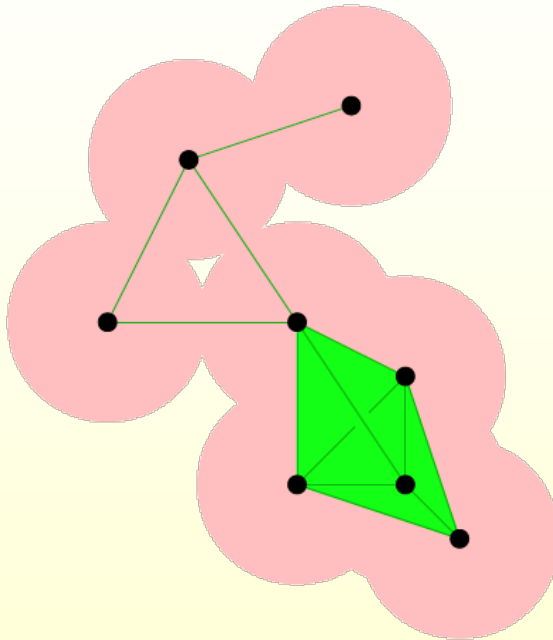


Rips



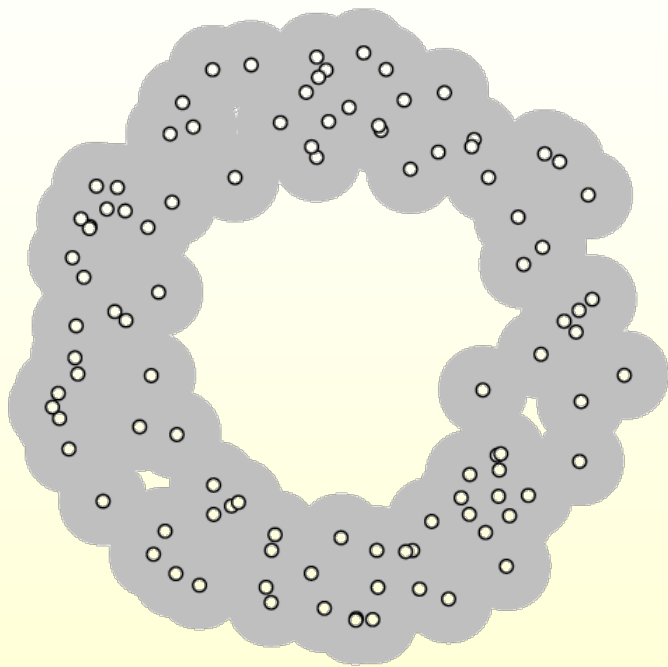
Combinatorial complexes provide discrete representations of the underlying space

# Čech Complex



- $C_\epsilon(M) = \{\text{conv } T \mid T \subseteq M, \bigcap_{m_i \in T} B_\epsilon(m_i) \neq \emptyset\}$ .
- $\sum_{k=0}^m \binom{m}{k} = 2^{m+1} - 1$
- $C_\epsilon(M) \simeq \tilde{M}$

# Čech Complex

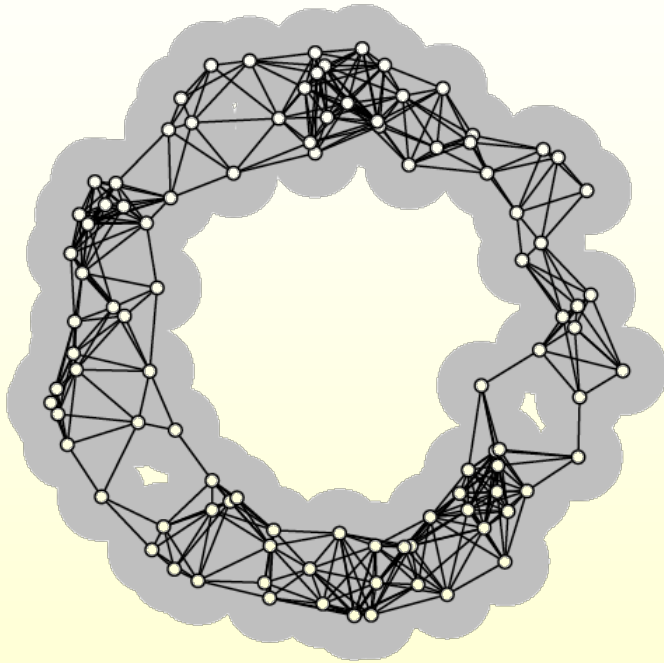


$\check{C}.6$

The Čech complex  $\check{C}_\epsilon$  encodes the intersection pattern of  $M_\epsilon$ : Encode:

Points as *vertices*  
(0-cells)

# Čech Complex



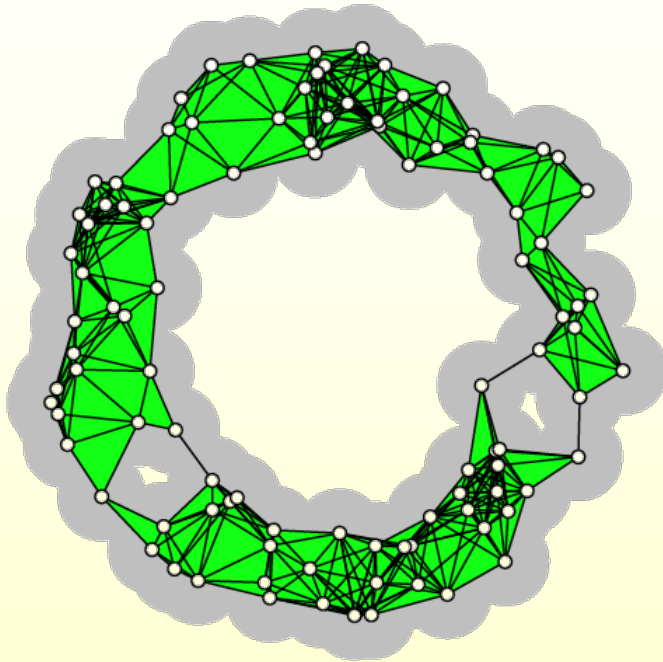
$\check{C}_{.6}$

The Čech complex  $\check{C}_\epsilon$  encodes the intersection pattern of  $M_\epsilon$ : Encode:

Points as *vertices*  
(0-cells)

Pairwise intersections  
as *edges* (1-cells)

# Čech Complex



$\check{C}_\epsilon$

The Čech complex  $\check{C}_\epsilon$  encodes the intersection pattern of  $M_\epsilon$ : Encode:

Points as *vertices*  
(0-cells)

Pairwise intersections  
as *edges* (1-cells)

Threeway intersections  
as *triangles* (2-cells)

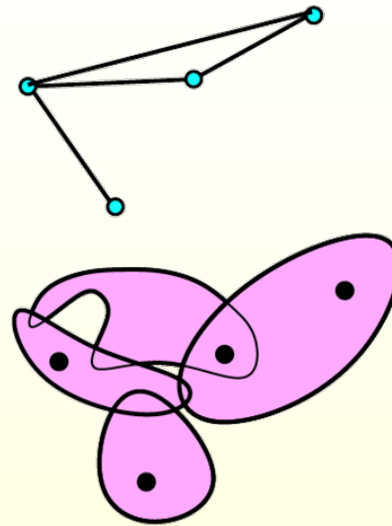
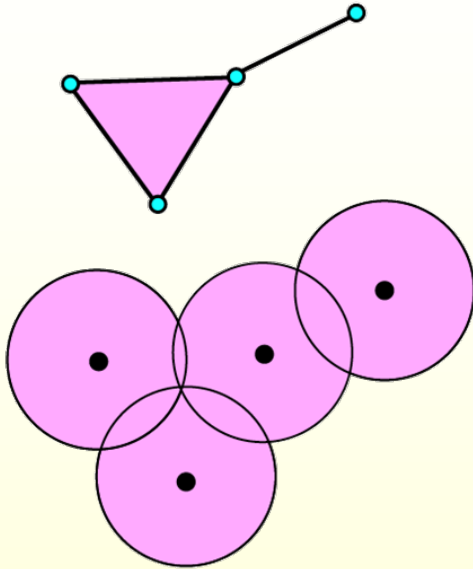
$k$ -way intersections as  
( $k+1$ )-cells

Lemma (Nerve Lemma, Leray '45)

$\check{C}_\epsilon$  is topologically equivalent to  $M_\epsilon$ .

Can be hard to compute ...

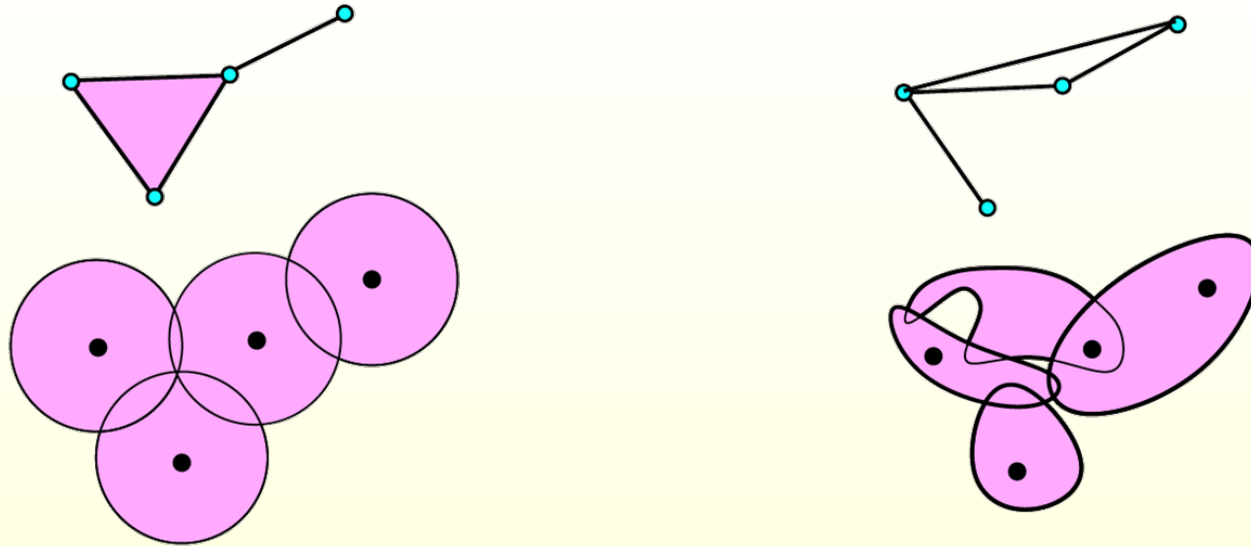
# General Čech Complex



- Let  $\mathcal{U} = (U_i)_{i \in I}$  be a covering of a topological space  $X$  by open sets:  
 $X = \cup_{i \in I} U_i$ .
- The Čech complex  $C(\mathcal{U})$  associated to the covering  $\mathcal{U}$  is the simplicial complex defined by:
  - the vertex set of  $C(\mathcal{U})$  is the set of the open sets  $U_i$
  - $[U_{i_0}, \dots, U_{i_k}]$  is a  $k$ -simplex in  $C(\mathcal{U})$  iff  $\cap_{j=0}^k U_{i_j} \neq \emptyset$ .



# General Čech Complex



**Nerve theorem (Leray):** If all the intersections between opens in  $\mathcal{U}$  are either empty or contractible then  $C(\mathcal{U})$  and  $X = \cup_{i \in I} U_i$  are homotopy equivalent.

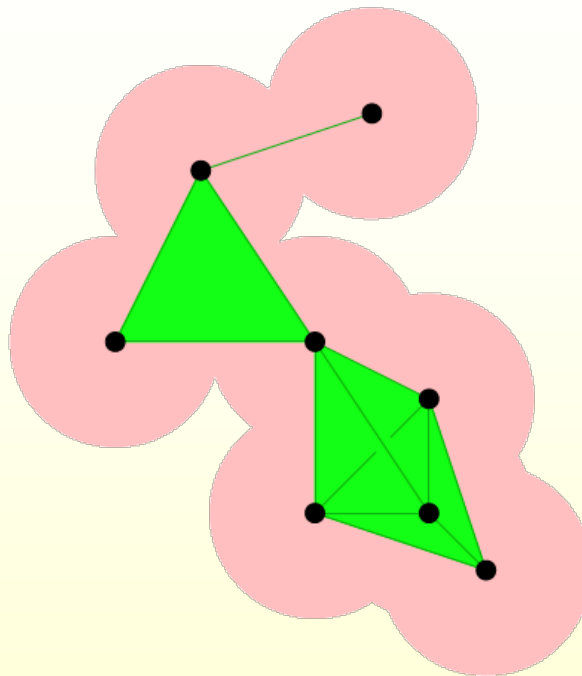
$\Rightarrow$  The combinatorics of the covering (a simplicial complex) carries the topology of the space.

**Warning:** even when the open sets are euclidean balls, the computation of the Čech complex is a very difficult task!

# Rips-Vietoris Complex

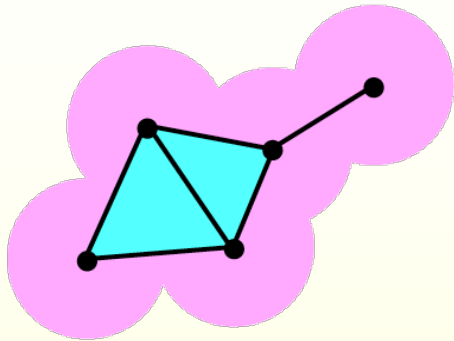
The “poor man’s” alternative  
to the Čech

This is a common complex  
For computations

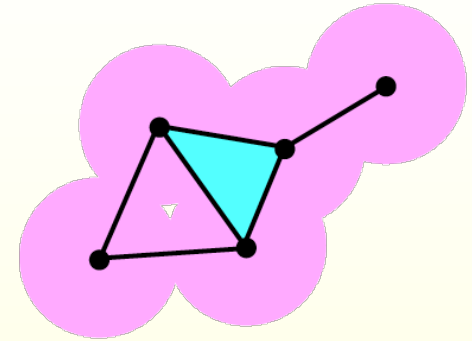


- $R_\epsilon(M) = \{\text{conv } T \mid T \subseteq M, d(m_i, m_j) < \epsilon, m_i, m_j \in T\}.$
- Still  $O\left(\binom{m}{k}\right)$  for the  $k$ th skeleton
- Need  $(k + 1)$ st skeleton for computing  $H_k$

# Rips vs. Čech



Rips vs Čech



Let  $L = \{p_0, \dots, p_n\}$  be a (finite) point cloud (in a metric space).

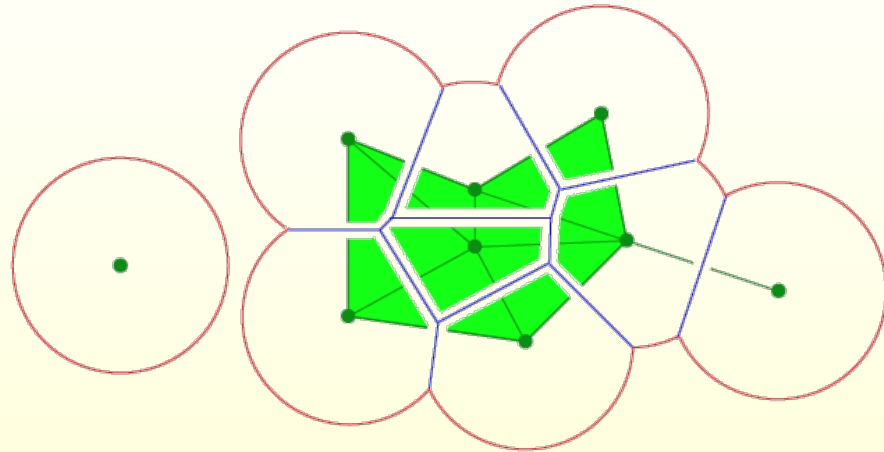
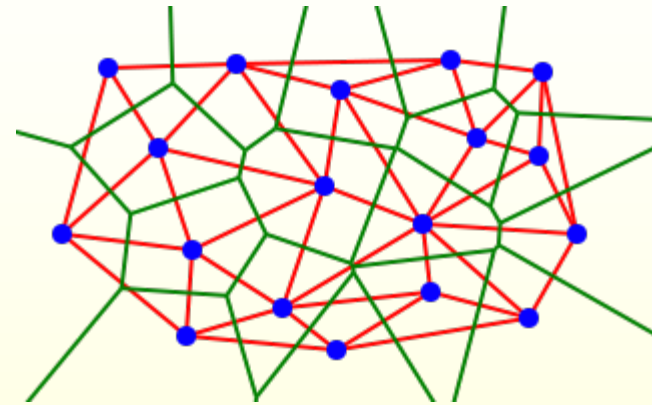
The **Rips complex**  $\mathcal{R}^\alpha(L)$ : for  $p_0, \dots, p_k \in L$ ,

$$\sigma = [p_0 p_1 \dots p_k] \in \mathcal{R}^\alpha(L) \text{ iff } \forall i, j \in \{0, \dots, k\}, d(p_i, p_j) \leq \alpha$$

- Easy to compute and fully determined by its 1-skeleton
- Rips-Čech interleaving: for any  $\alpha > 0$ ,

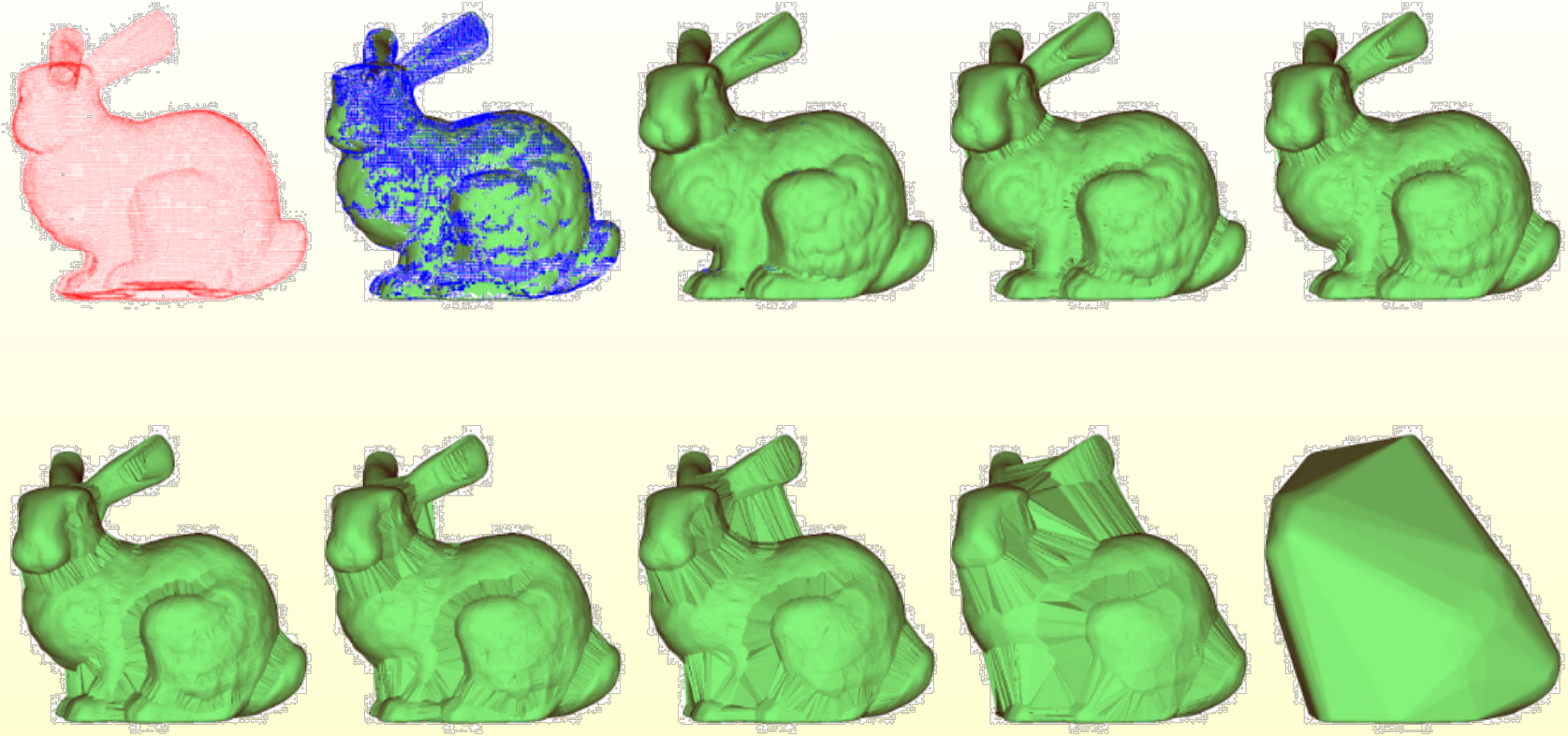
$$\mathcal{C}^{\frac{\alpha}{2}}(L) \subseteq \mathcal{R}^\alpha(L) \subseteq \mathcal{C}^\alpha(L) \subseteq \mathcal{R}^{2\alpha}(L) \subseteq \dots$$

# Alpha Complex



- $V(m_i) = \{x \in \mathbb{R}^3 \mid d(x, m_i) \leq d(x, m_j) \forall m_j \in M\}$
- $\hat{V}(m_i) = B_\epsilon(m_i) \cap V(m_i)$
- $A_\epsilon = \left\{ \text{conv} T \mid T \subseteq M, \bigcap_{m_i \in T} \hat{V}(m_i) \neq \emptyset \right\}$
- $A_\epsilon(M) \simeq \tilde{M}$ ,  $A_\epsilon \subseteq D$ , the **Delaunay complex**
- $O(n \log n + n^{\lceil d/2 \rceil})$

# Alpha Complexes on the Stanford Bunny



- 34,834 points, 1,026,111 complexes

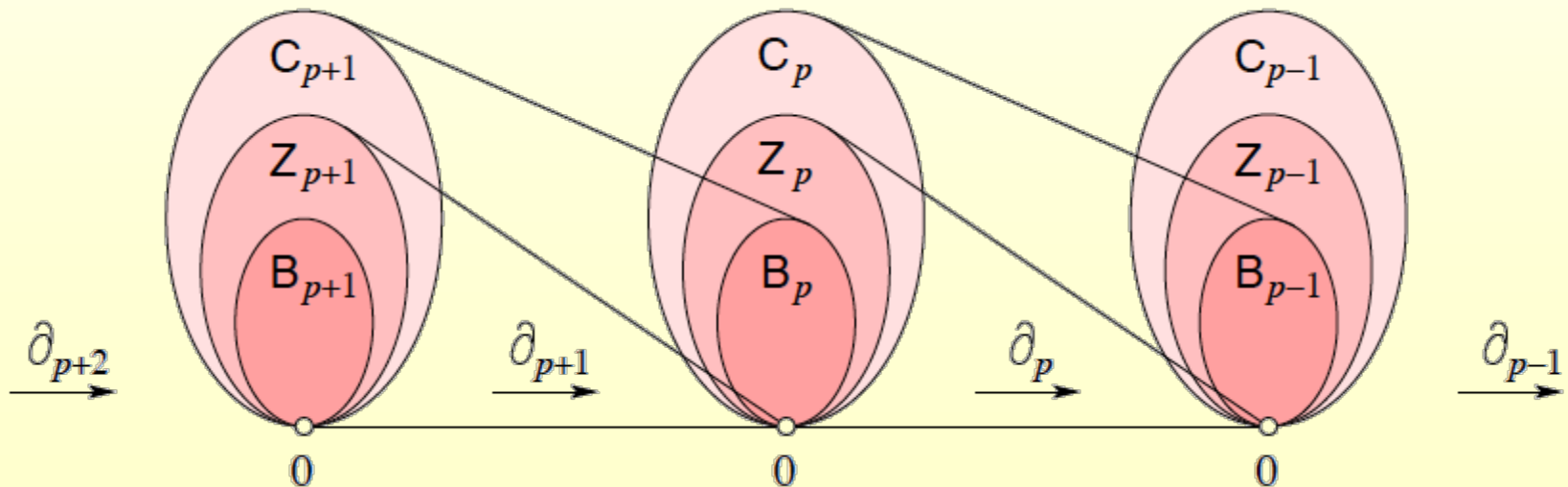
# Computing Homology via Bases

# Homology

- The  $k$ th homology group is

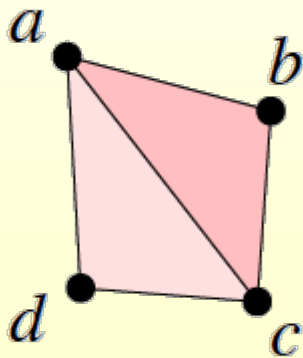
$$H_k = Z_k / B_k = \ker \partial_k / \text{im } \partial_{k+1}.$$

- Compute a basis for  $\ker \partial_k$
- Compute a basis for  $\text{im } \partial_{k+1}$



# Matrix Representation of $\partial$

- Boundary homomorphism is linear, so it has a matrix
- $\partial_k : \mathbf{C}_k \rightarrow \mathbf{C}_{k-1}$
- Use oriented simplices as bases for domain and codomain!
- $M_k$  is the **standard matrix representation** for  $\partial_k$



$$M_1 = \begin{bmatrix} & ab & bc & cd & ad & ac \\ a & -1 & 0 & 0 & -1 & -1 \\ b & 1 & -1 & 0 & 0 & 0 \\ c & 0 & 1 & -1 & 0 & 1 \\ d & 0 & 0 & 1 & 1 & 0 \end{bmatrix}$$

[Two glued triangles, not the tetrahedron ...]



# Elementary Matrix Operations

- The **elementary row operations** on  $M_k$  are
  1. exchange row  $i$  and row  $j$ ,
  2. multiply row  $i$  by  $-1$ ,
  3. replace row  $i$  by  $(\text{row } i) + q(\text{row } j)$ , where  $q$  is an integer and  $j \neq i$ .
- Similar **elementary column operations** on columns
- Effect: change of bases

# Questions

- How do we find cycles?
- How do we find boundaries?
- What does a free generator correspond to?
- How does a torsional element appear?



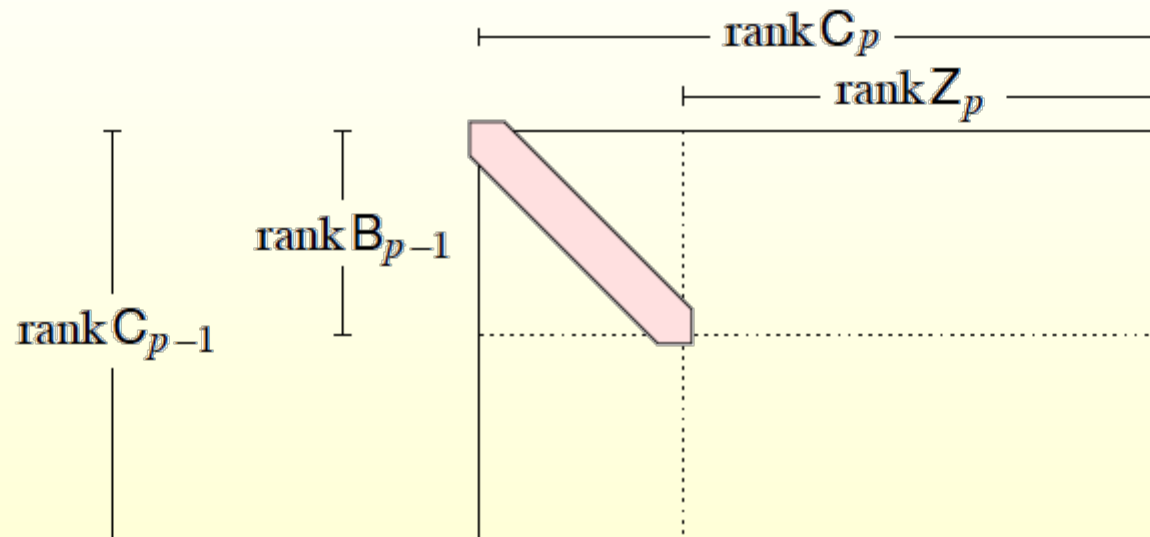
# Reduction Algorithm

- Like Gaussian elimination, we keep changing the basis to get to the **(Smith) normal form**:

$$\tilde{M}_k = \left[ \begin{array}{cc|c} b_1 & 0 & 0 \\ & \ddots & 0 \\ 0 & b_{l_k} & 0 \\ \hline & 0 & 0 \end{array} \right]$$

- $l_k = \text{rank } M_k = \text{rank } \tilde{M}_k, b^i \geq 1$
- $b_i | b_{i+1}$  for all  $1 \leq i < l_k$   $b_i = 1 \quad \forall i, \text{ if no torsion}$

# Smith Normal Form

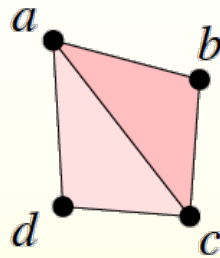


Introduce columns from left to right

Keep doing Gaussian elimination steps ...

For a complex with  $m$  simplices, this can take  $O(m^3)$  operations

# Reduction Example



$$\tilde{M}_1 = \left[ \begin{array}{c|ccc|cc} & cd & bc & ab & z_1 & z_2 \\ \hline d-c & 1 & 0 & 0 & 0 & 0 \\ c-b & 0 & 1 & 0 & 0 & 0 \\ b-a & 0 & 0 & 1 & 0 & 0 \\ \hline a & 0 & 0 & 0 & 0 & 0 \end{array} \right]$$

- $z_1 = ad - bc - cd - ab$  and  $z_2 = ac - bc - ab$  form a basis for  $\mathbf{Z}_1$
- $\{d - c, c - b, b - a\}$  is a basis for  $\mathbf{B}_0$

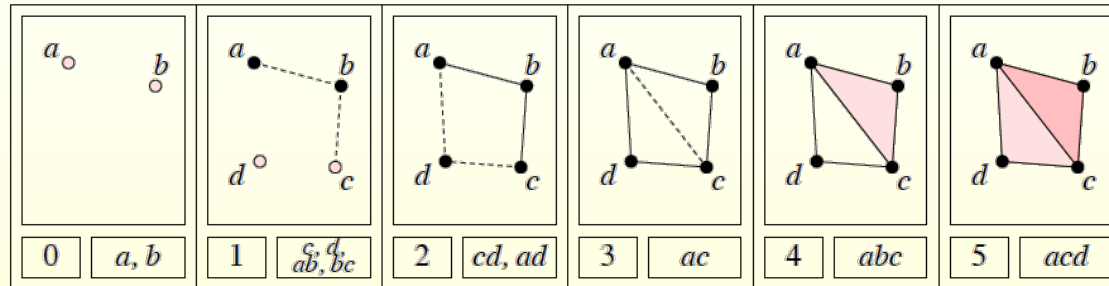
# Reduction Example

$$M_2 = \left[ \begin{array}{c|cc} & abc & acd \\ \hline ac & -1 & 1 \\ ad & 0 & -1 \\ cd & 0 & 1 \\ bc & 1 & 0 \\ ab & 1 & 0 \end{array} \right]$$

$$\tilde{M}_2 = \left[ \begin{array}{c|cc} & -abc & -acd + abc \\ \hline ac - bc - ab & 1 & 0 \\ ad - cd - bc - ab & 0 & 1 \\ cd & 0 & 0 \\ bc & 0 & 0 \\ ab & 0 & 0 \end{array} \right]$$

# Can Simplify for Complexes in $R^3 / S^3$

- Use a **filtration**



- A **filtration** of a complex  $K$  is  $\emptyset = K^0 \subseteq K^1 \subseteq \dots \subseteq K^m = K$ .
- A filtration is a partial ordering
- Sort according to dimension
- Break other ties arbitrarily
- Algorithm for  $K = S^3$

# Alexander Duality, Complements

- **Alexander Duality:**
  - $\beta_0$  measures the number of components of the complex.
  - $\beta_1$  is the rank of a basis for the **tunnels**.
  - $\beta_2$  counts the number of **voids** in the complex.
- An incremental approach:
  - add each simplex in turn
  - check to see if we form a new cycle class or destroy one.

$$\beta_k = \text{rank (of the free part) of } H_k$$

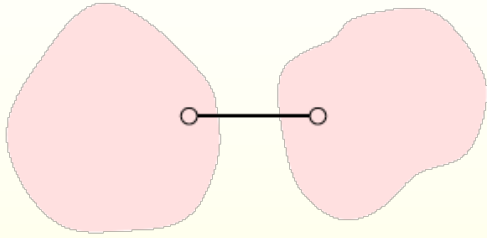


# Vertices

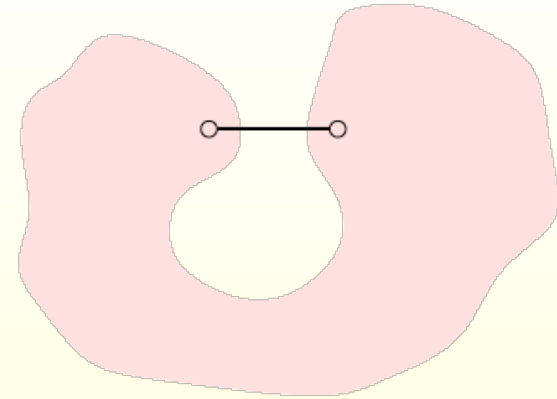
- Vertices always add a new component, so  $\beta_0^{++}$ .
- Union-find data-structure:
  - MAKESET: initializes a set with an item
  - FIND: finds the set an element belongs to
  - UNION: forms the union of two sets
- Very simple to implement
- $O(n)$  space
- Amortized  $\alpha(m)$  FIND, UNION
- MAKESET for each vertex

$\beta_0$  requires maintaining connected components

# Edges



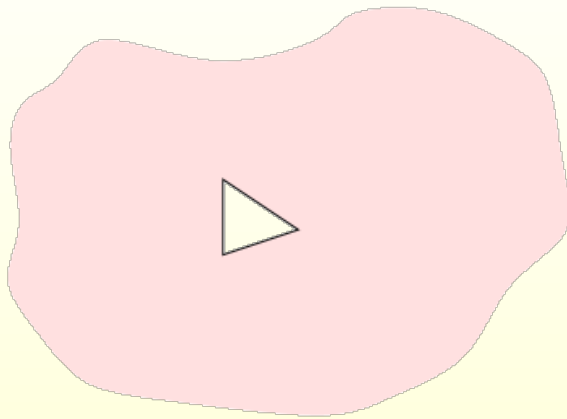
(a)  $\beta_0^{--}$



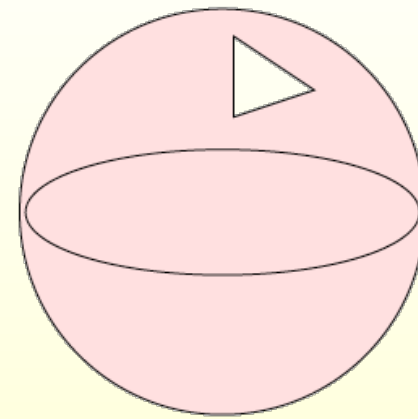
(b)  $\beta_1^{++}$

- (a) Two FINDs, one UNION
- (b) Two FINDs

# Triangles and Tetrahedra

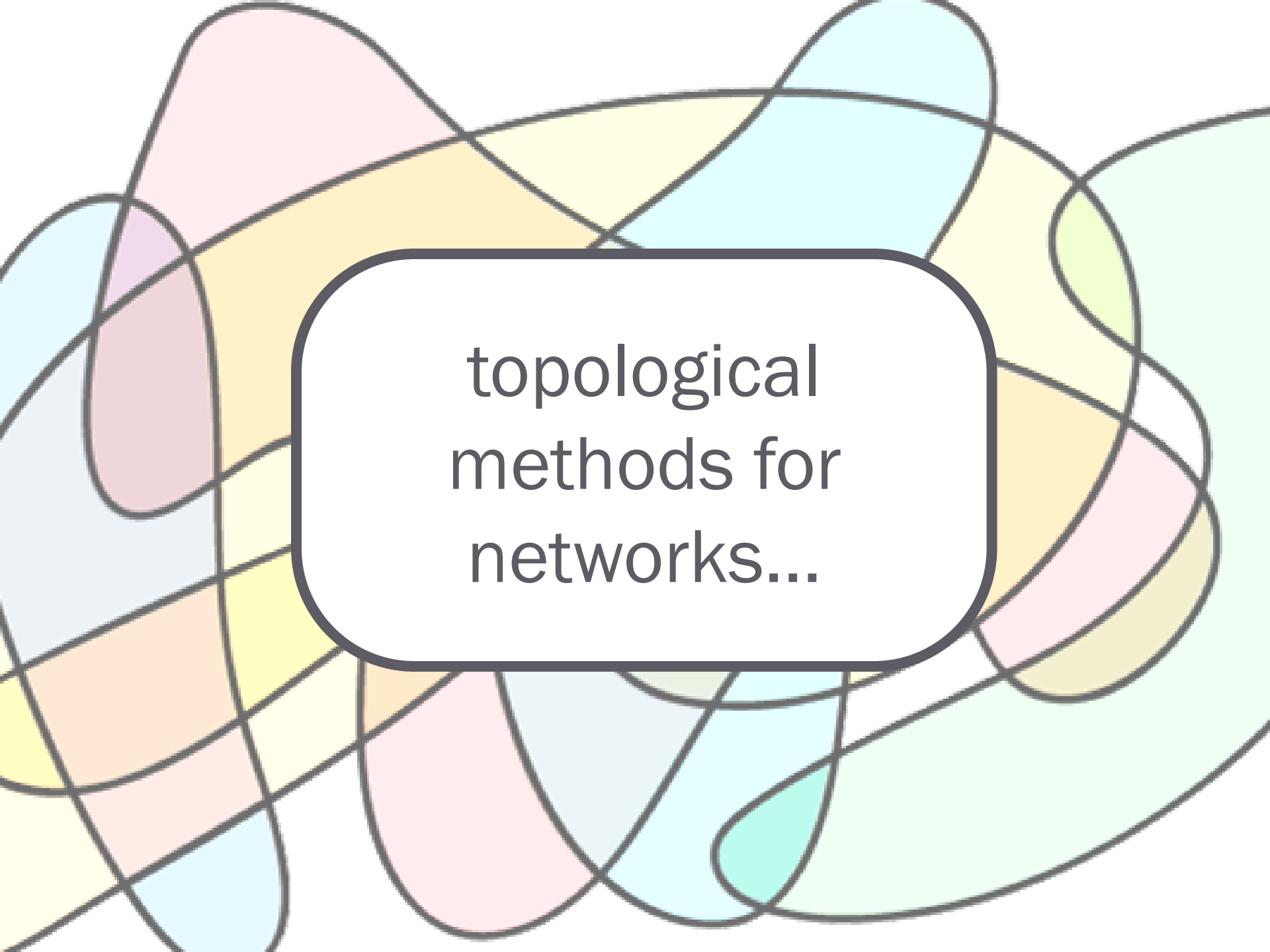


(a)  $\beta_{1--}$



(b)  $\beta_{2++}$

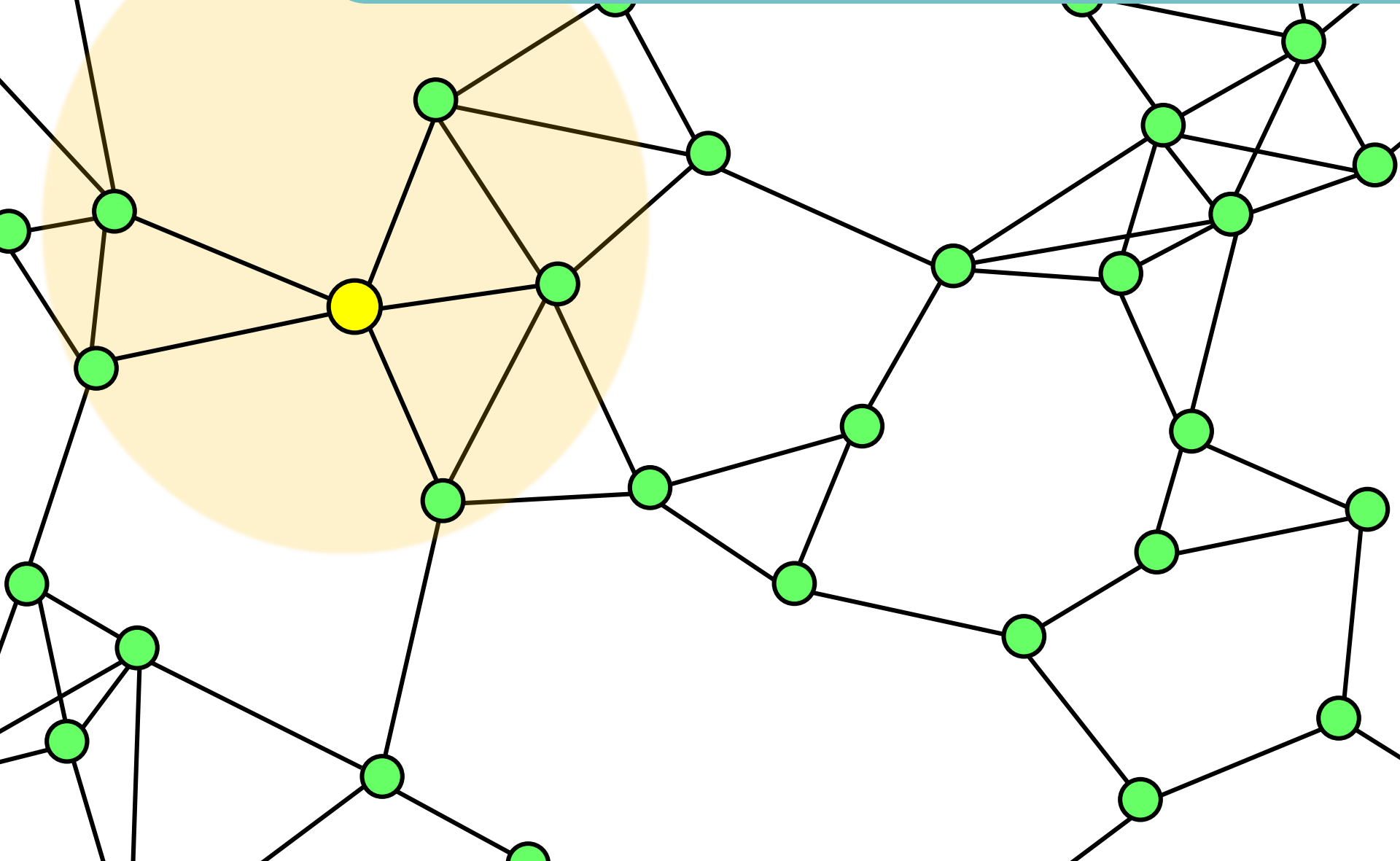
- Tetrahedra always fill voids, so  $\beta_{2--}$

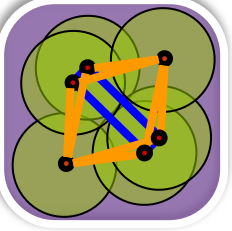
The background consists of several overlapping, rounded, organic shapes in various colors including light blue, yellow, pink, light green, and light purple. These shapes are outlined with thin black lines and overlap each other in a complex, non-linear fashion. In the center of the image, there is a white rounded rectangle with a thick black border. Inside this rectangle, the text "topological methods for networks..." is written in a dark grey, sans-serif font, centered both horizontally and vertically.

topological  
methods for  
networks...

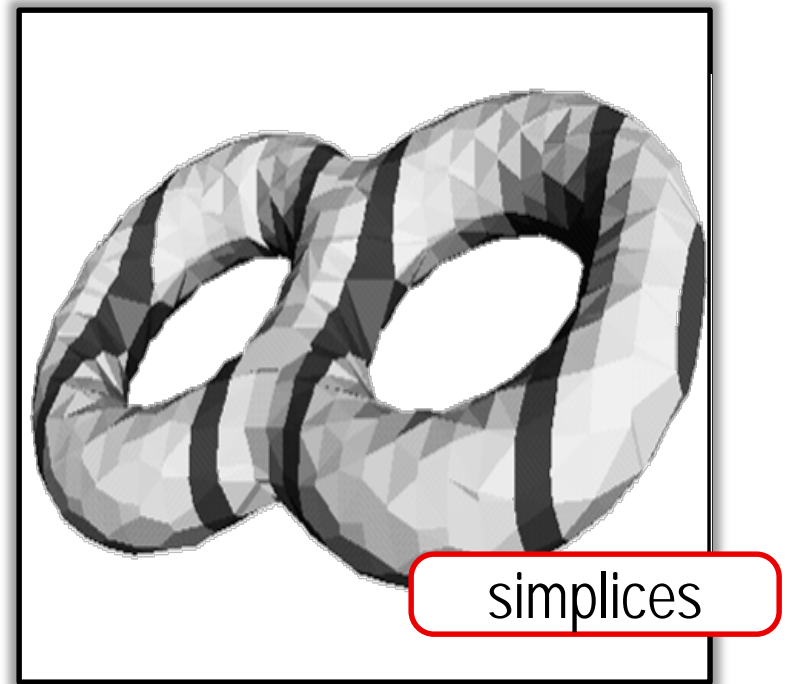
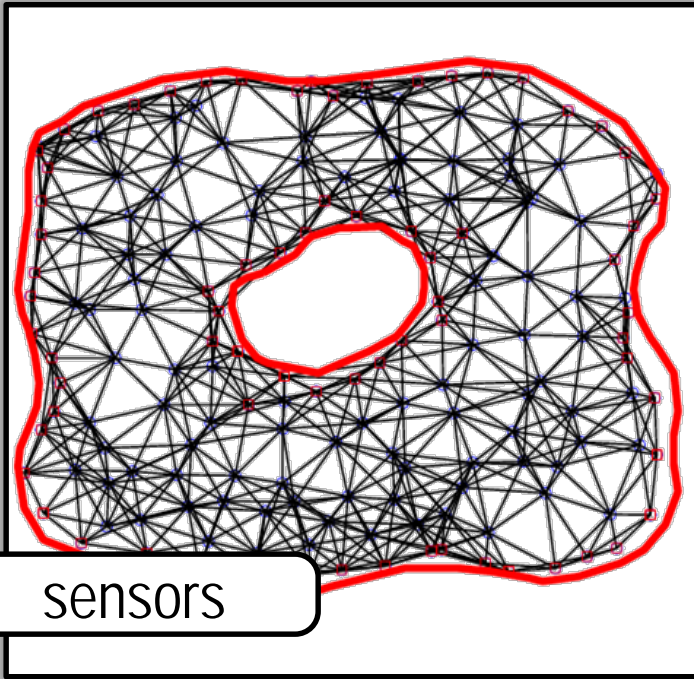


# homological coverage

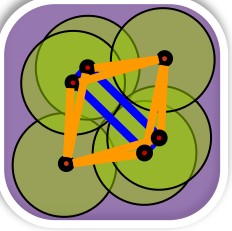




# homological coverage



sensors and simplices each have knowledge only of their identities and of their local connectivity...



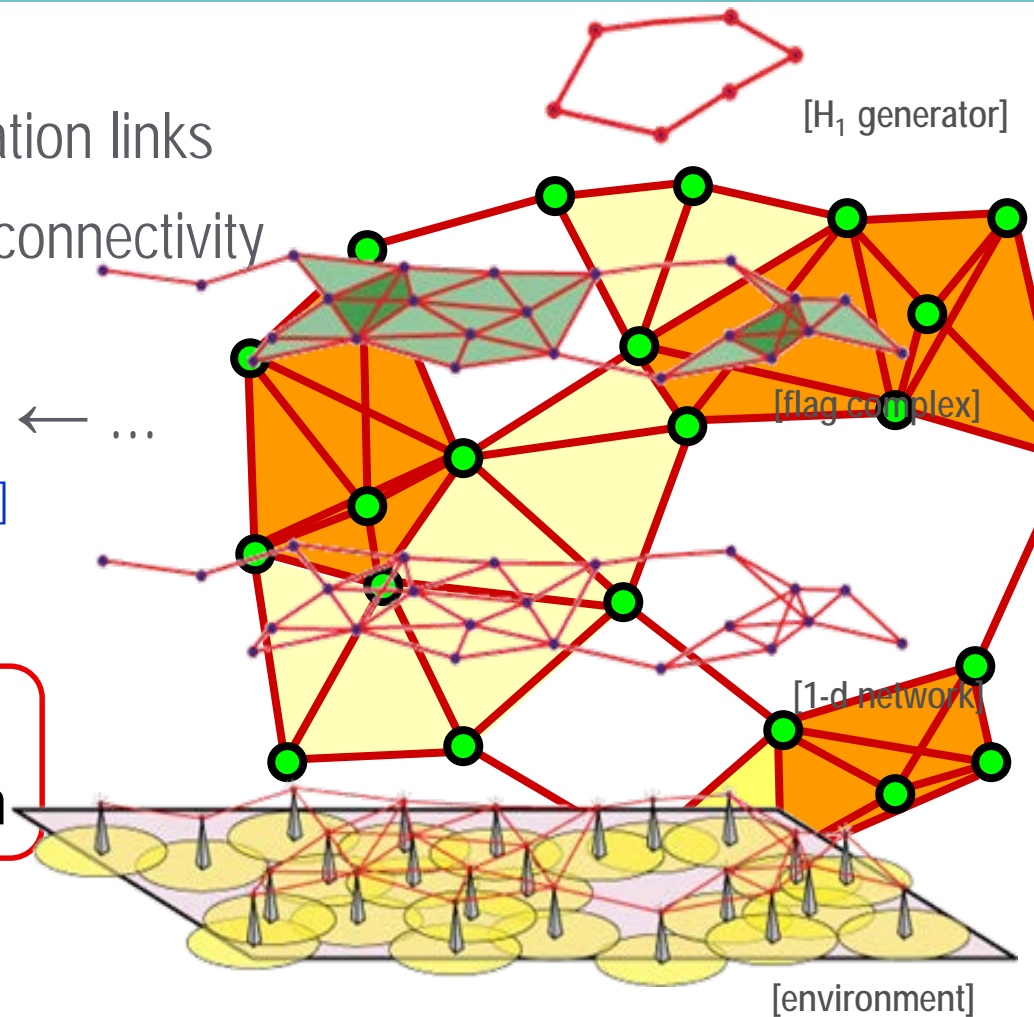
# networks & complexes

given node id's, local communication links  
count nodes & cancel via signal connectivity

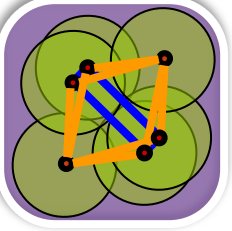
$$C_0 \leftarrow C_1 \leftarrow C_2 \leftarrow C_3 \leftarrow \dots$$

[nodes]    [pairs]    [triples]    [quads]

the **Rips complex** of a network  
is the maximal simplicial completion

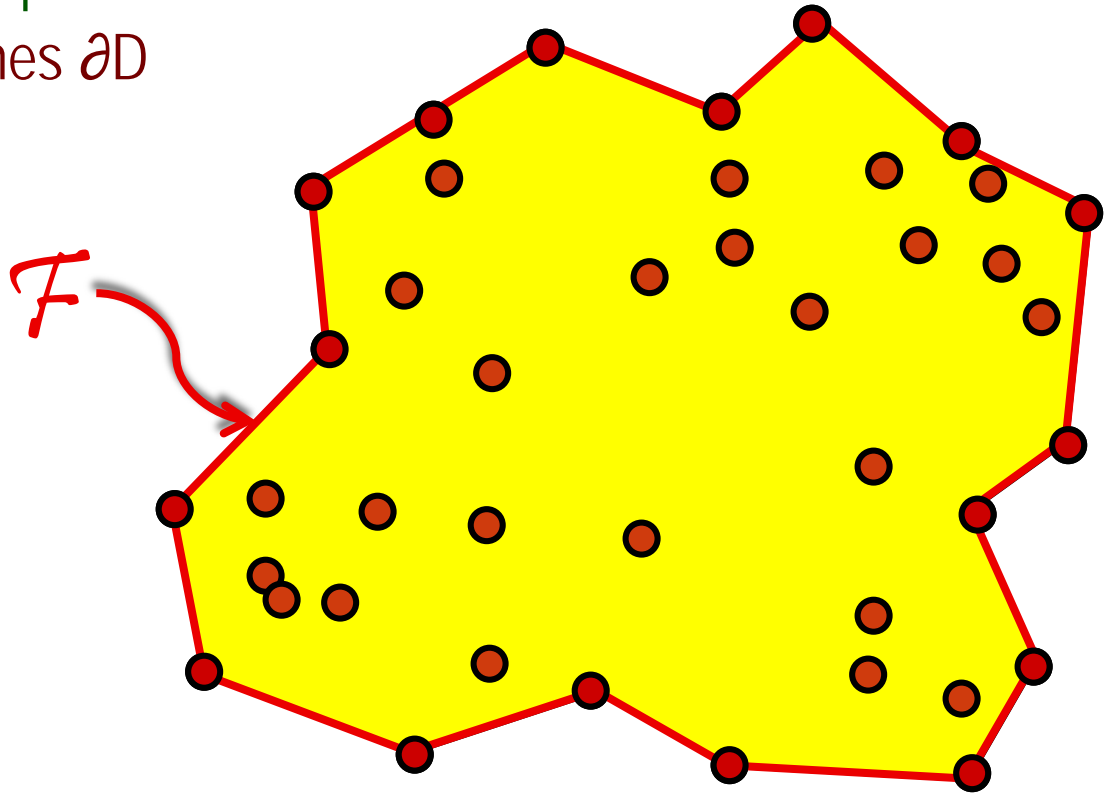
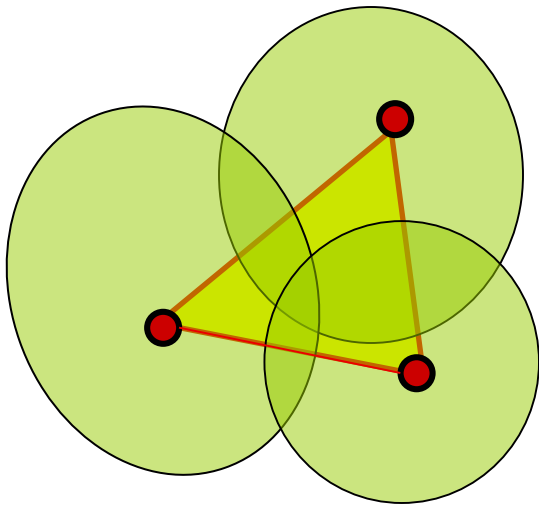


homology converts higher-order network connectivity into global structure...  
...without coordinates; density assumptions; uniform distributions, etc.

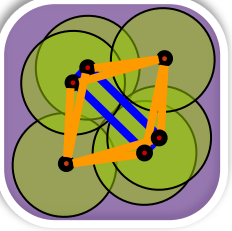


# coverage assumptions

1. compact polygonal domain  $D$  in  $\mathbb{R}^2$
2. nodes broadcast unique id's to neighbors
3. coverage regions of a 2-simplex of connected nodes contain the convex hull
4. dedicated fence cycle defines  $\partial D$







# coverage criterion

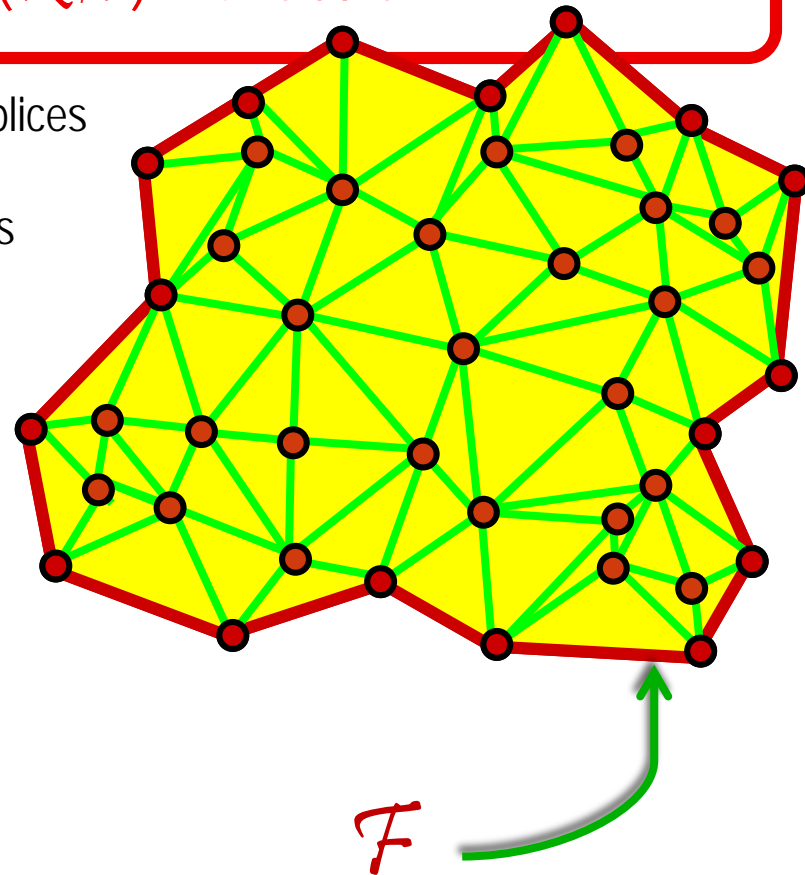
**Theorem [DG]:** under above assumptions, the sensor network covers the domain without gaps if there exists  $[\alpha]$  in  $H_2(\mathcal{R}, \mathcal{F})$  with  $\partial\alpha \neq 0$

intuition:  $[\alpha]$  "triangulates" the domain with covered simplices

proof: build a commutative diagram of homology groups  
 map  $\sigma: (\mathcal{R}, \mathcal{F}) \rightarrow (\mathbb{R}^2, \partial D)$  convex hulls of simplices

$$\begin{array}{ccc}
 H_2(\mathcal{R}, \mathcal{F}) & \xrightarrow{\partial_*} & H_1(\mathcal{F}) \\
 \swarrow & & \downarrow \sigma_* \approx \\
 H_2(\mathbb{R}^2 - p, \partial D) = 0 & & H_1(\partial D) \\
 \searrow & & \uparrow \partial_* \\
 H_2(\mathbb{R}^2, \partial D) & \xrightarrow{\partial_*} & H_1(\partial D)
 \end{array}$$

if  $p$  lies in  $D - \sigma(\mathcal{R})$ , then the left passes through zero  
 commutativity of diagram yields a contradiction



**The End**