



Resource Conservation in Sensor Networks

- Byers, J., and Nasser, G. "Utility-Based Decision-Making in Wireless Sensor Networks". Proceedings of IEEE MobiHOC 2000, Boston, MA, August 2000.
- Zhao, F., Shin, J., and Reich, J. "Information-Driven Dynamic Sensor Collaboration for Tracking Applications". Proceedings of IEEE Signal Processing Magazine, March 2002.



Overview

- Motivation
- Byers 2000
 - Cost-Utility Model
 - Algorithm
 - Experimental Results
- Zhao 2002
 - IDSQ Model
 - Algorithm
 - Experimental Results



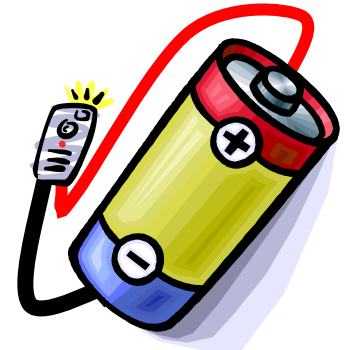
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Power & Sensor Conservation

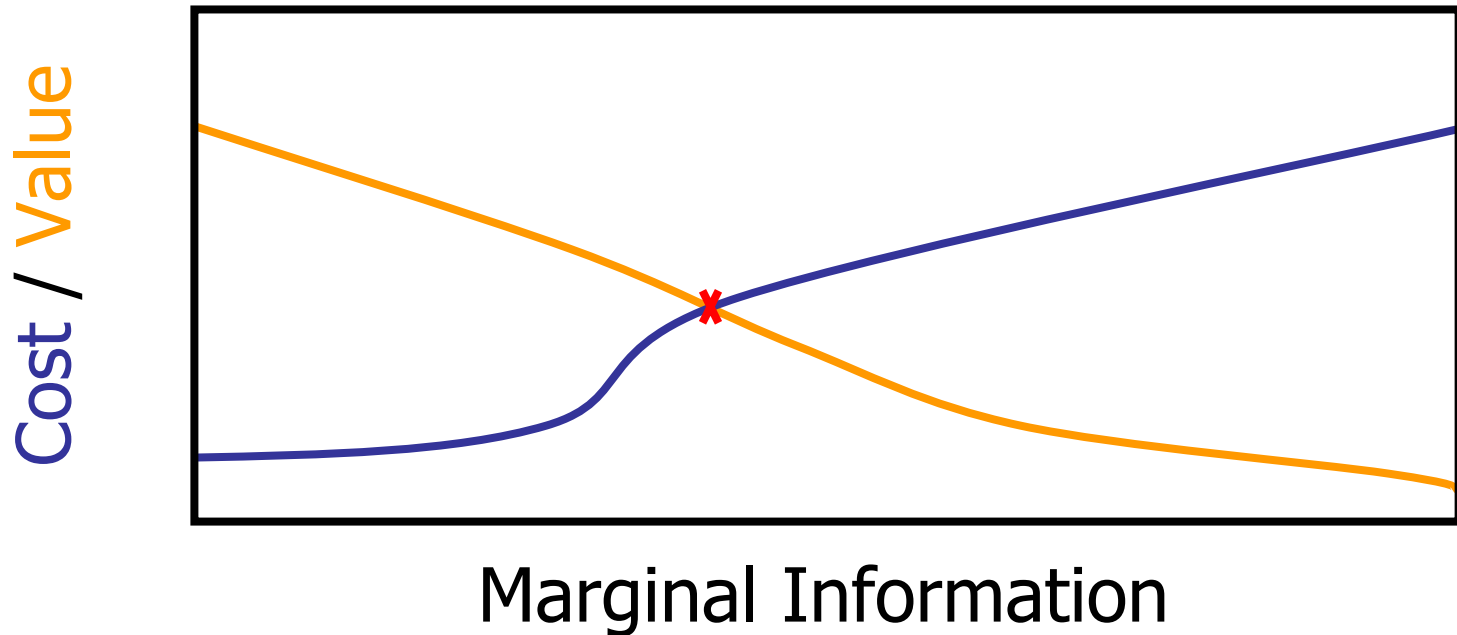
- Sensors have limited power
 - Reliant on non-renewable batteries
 - Battery technology not improving
- Conserving power
 - Increases sensor life
 - Makes sensor smaller
 - Decreases per sensor cost
- Also, minimize sensor count





Information Utility

- Information has diminishing marginal returns

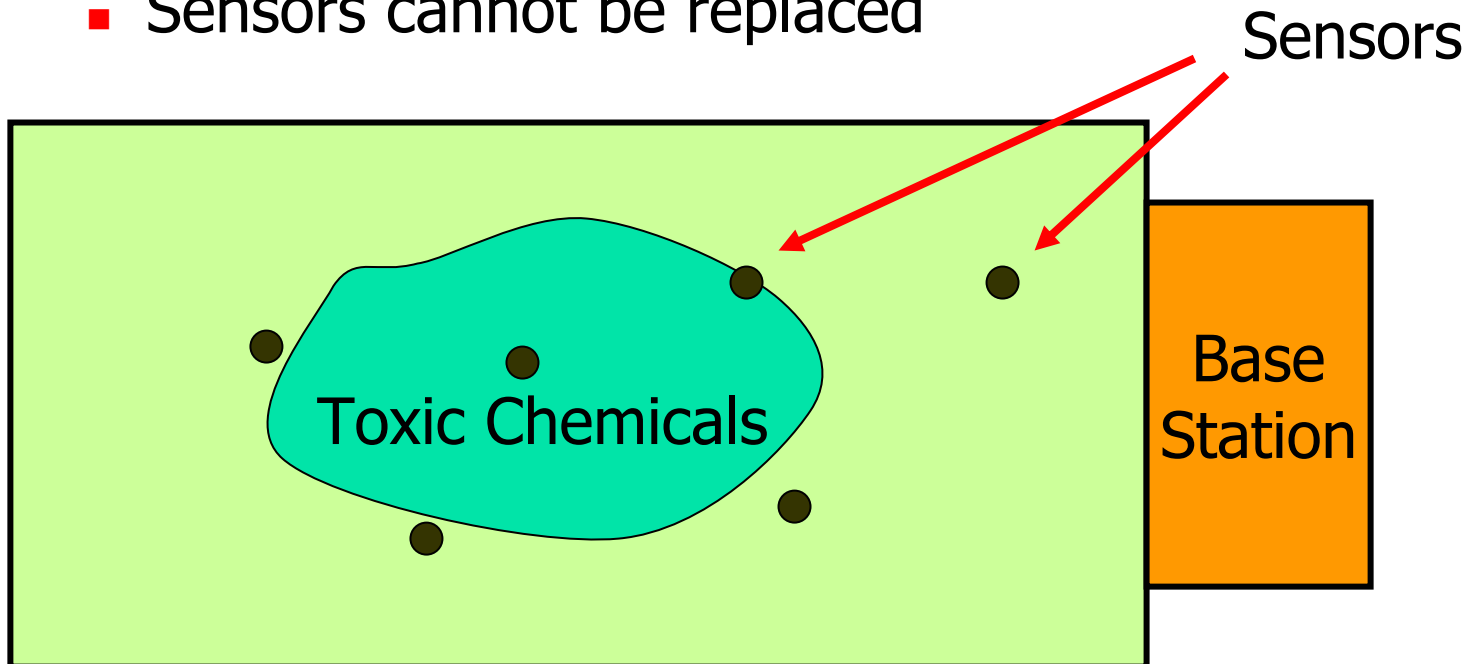




Example: Monitoring Toxicity

- Assumptions:

- Sensors connected to base station
- Sensors cannot be replaced





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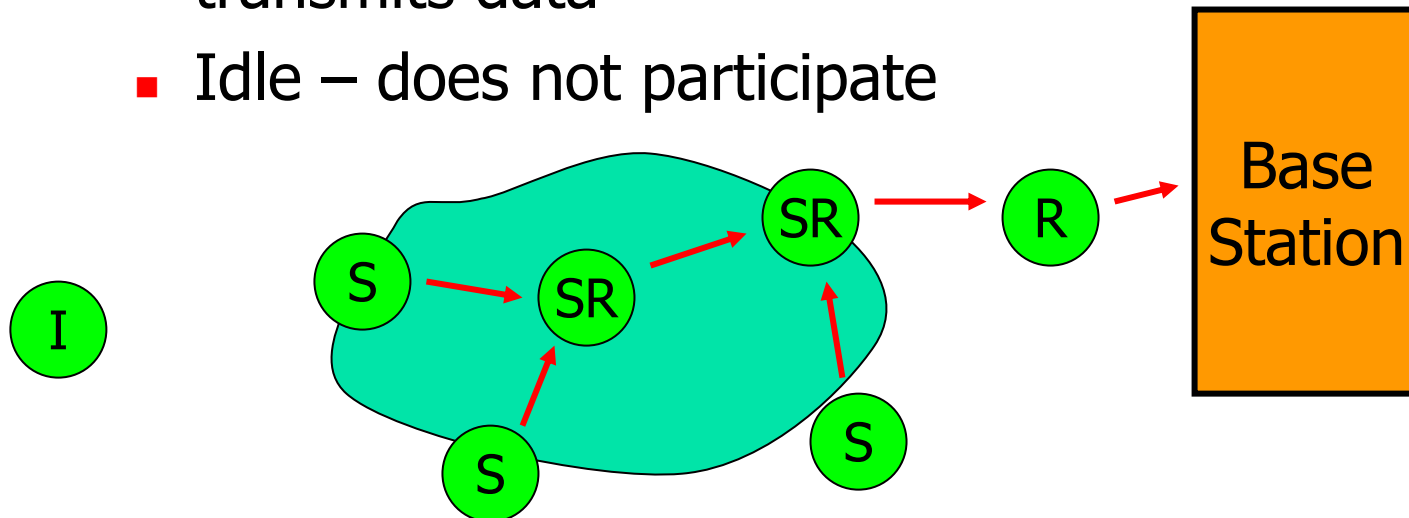
Network Model

- Assumptions:
 - Isotropic sensor transmission with fixed range R
 - No communication medium conflict
- Let $G = (V, E)$
 - V = Set of all operational sensors
 - E = Communication links between sensors
- An edge exists when the physical distance between a pair of nodes is below R



Sensing Model

- Nodes can operate in one of four modes:
 - Only Routing – receives and transmits data
 - Only Sensing – senses and transmits data
 - Both Routing and Sensing – receives, senses, and transmits data
 - Idle – does not participate





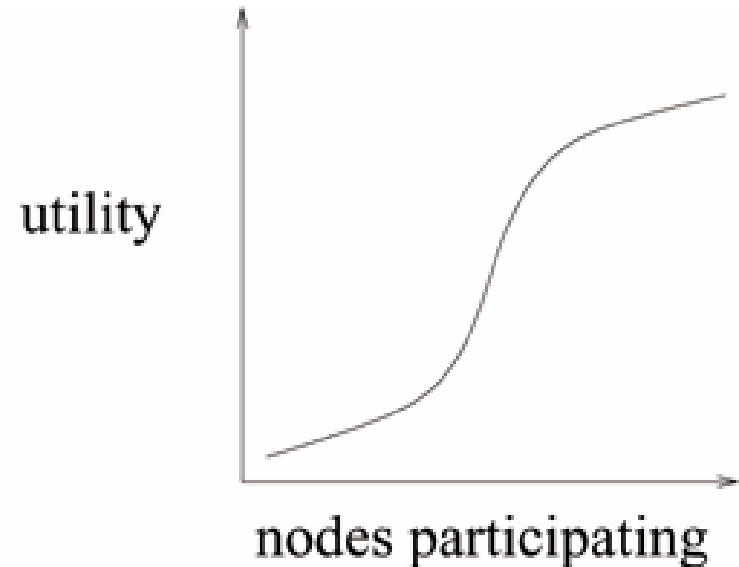
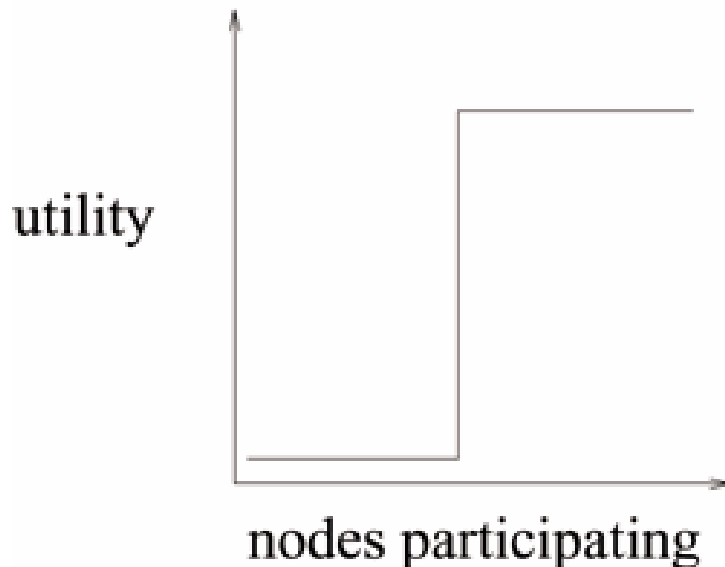
Data Aggregation

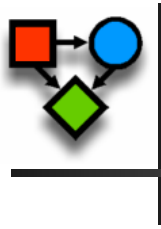
- Want to perform lossy compression of data to save communication costs
- Most useful when:
 - Information is of little value (e.g. little change)
 - Sensors provide redundant information
 - Large tree depth



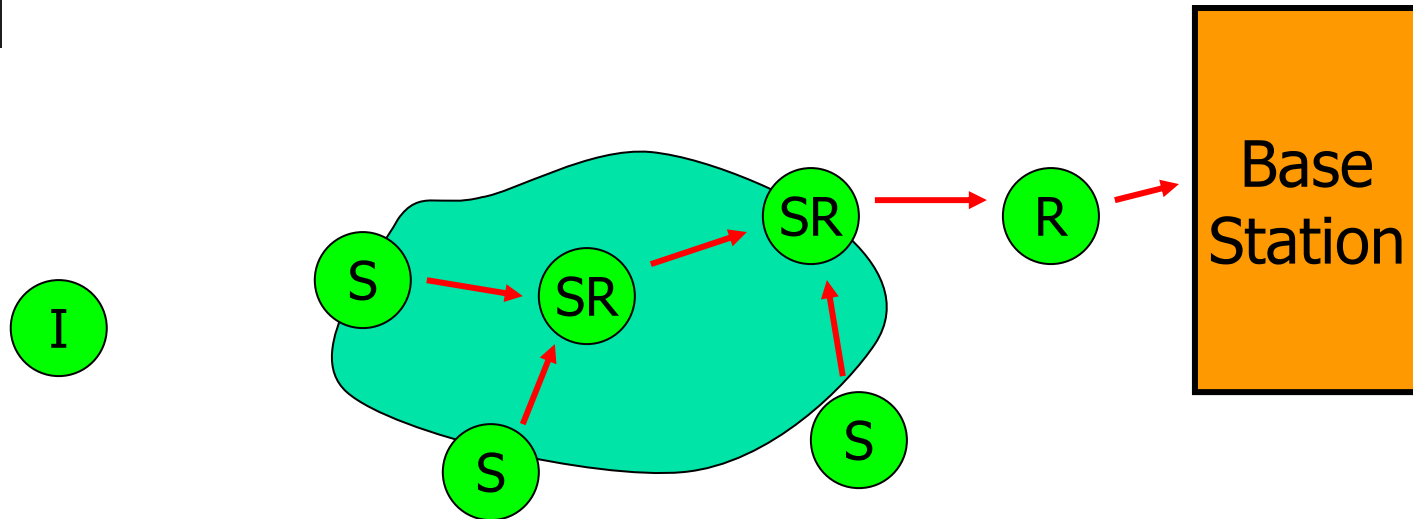
Information Utility

- Let S be the sensors that are chosen to sense
- Model utility based on S such that there exists a mapping $U: S^* \rightarrow [0,1]$
- Possible utilities:

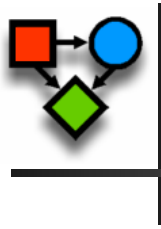




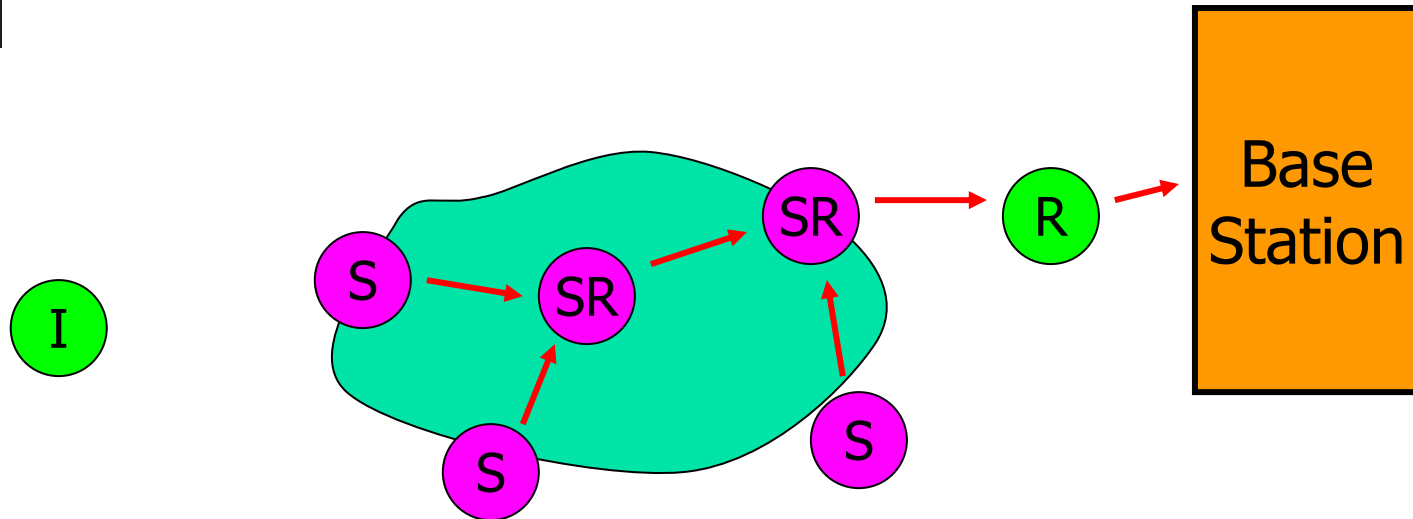
Costs



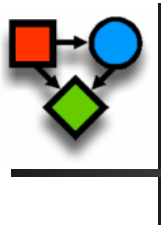
- S_t is the set of sensors sensing at time t
- $U(S_t)$ = utility from the sensors active at time t
- R_t is the set of sensors routing at time t
- c_s = sensing cost, c_t = transmitting cost, c_r = receiving cost, and c_a is aggregation cost



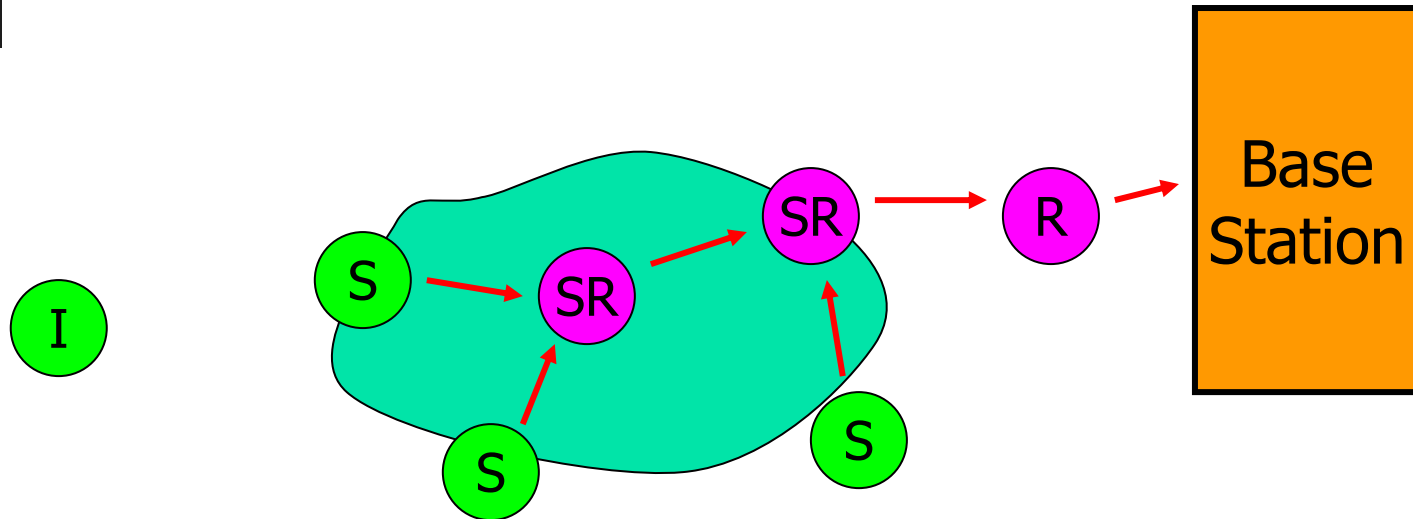
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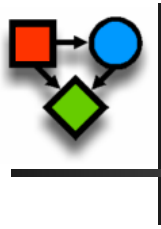
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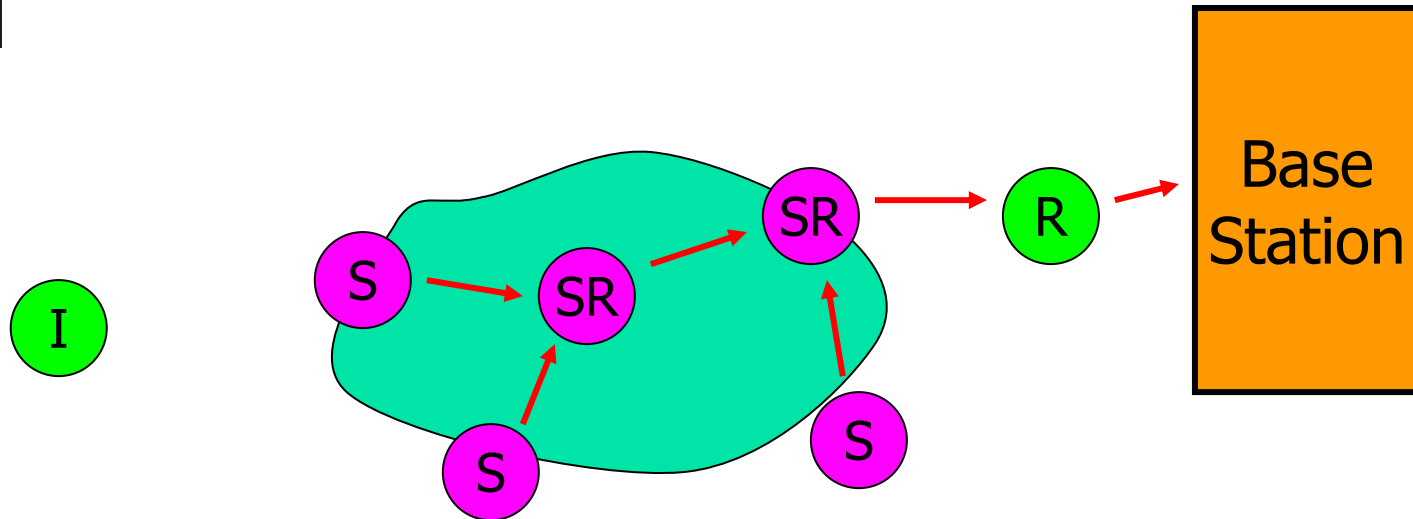
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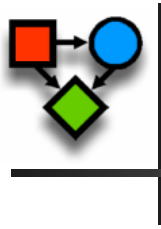
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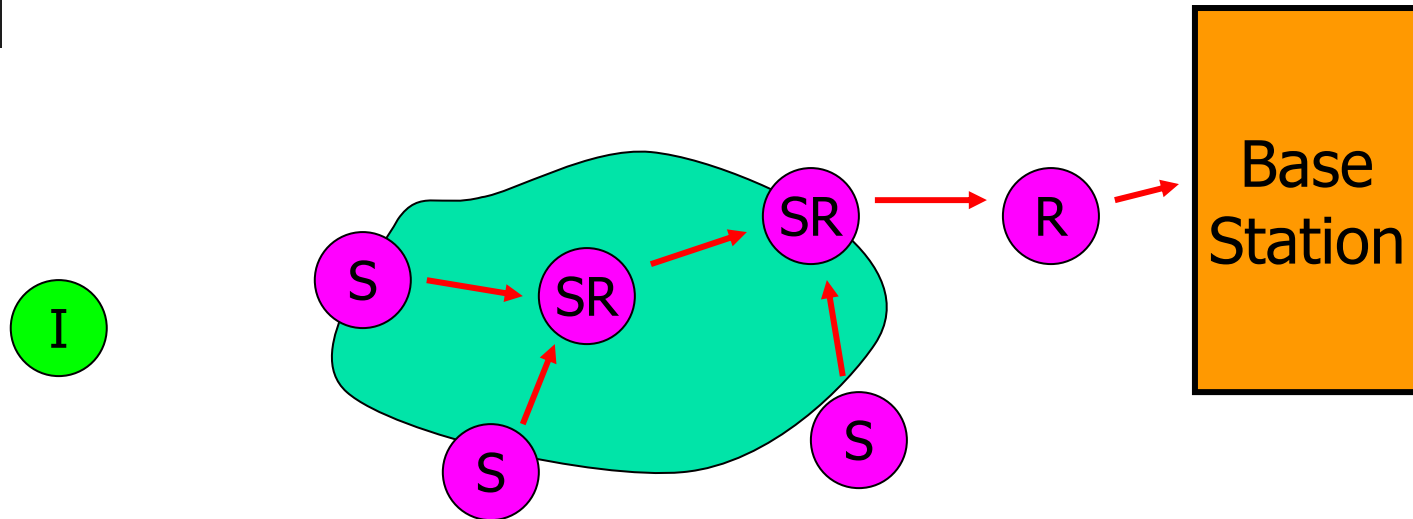
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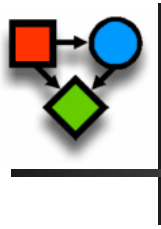
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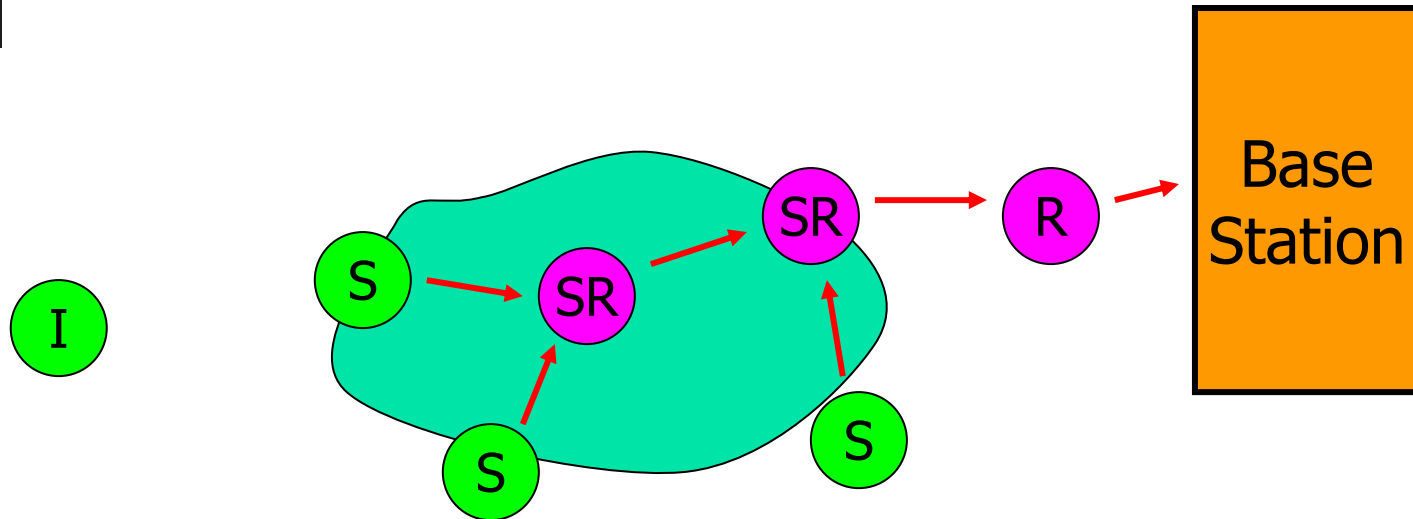
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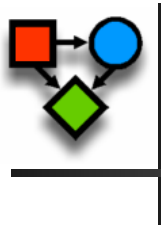
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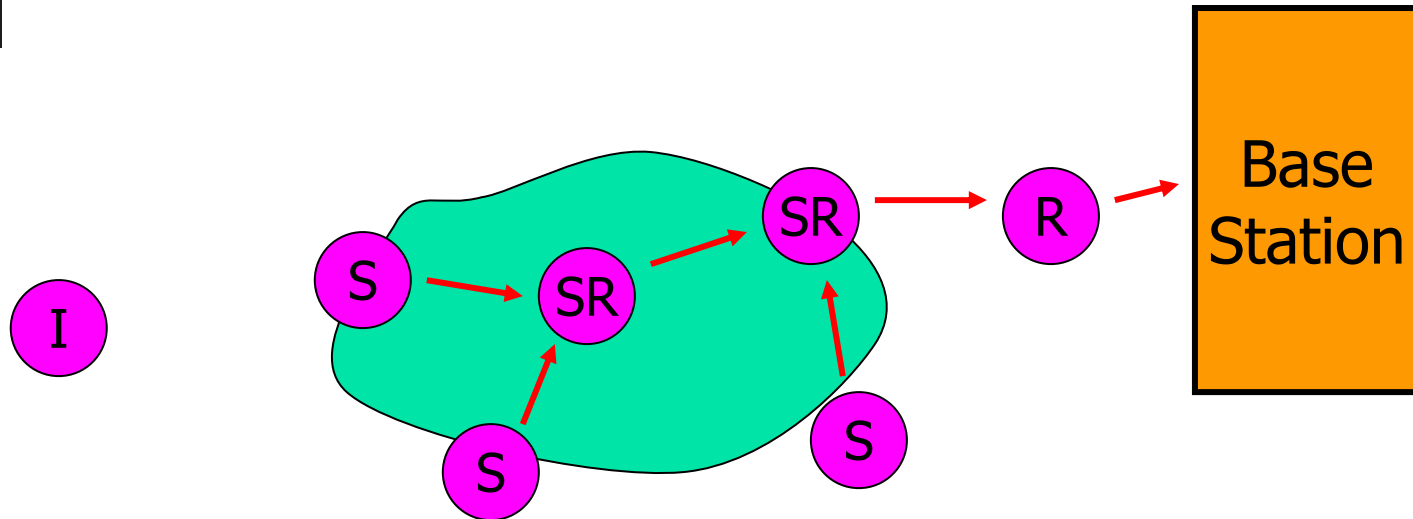
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Objective Function

- Objective: maximize utility across all time
- Subject to:
for each node across all time,
all costs incurred by node \leq power at node
- c_s = sensing cost
- c_t = transmitting cost
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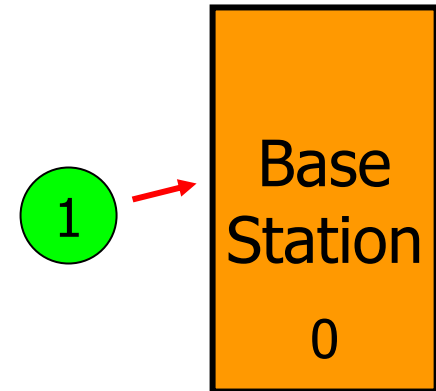


Network Initialization



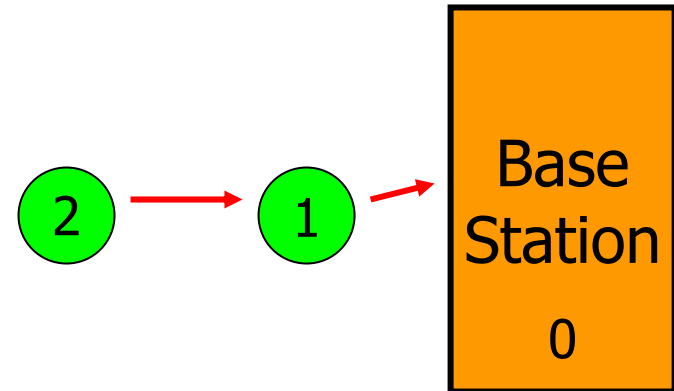


Network Initialization



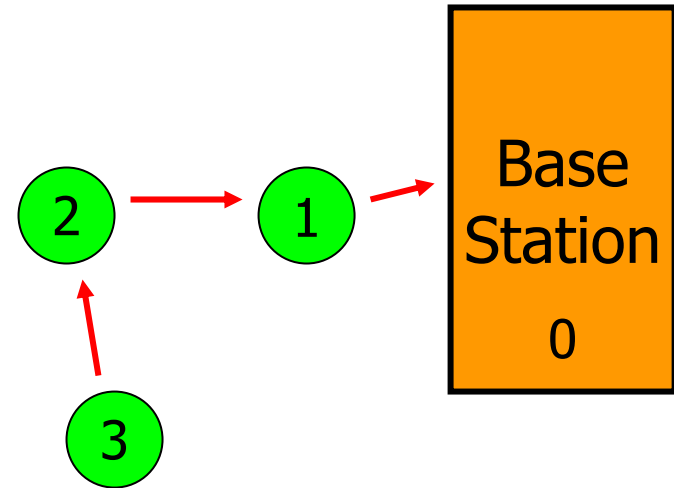


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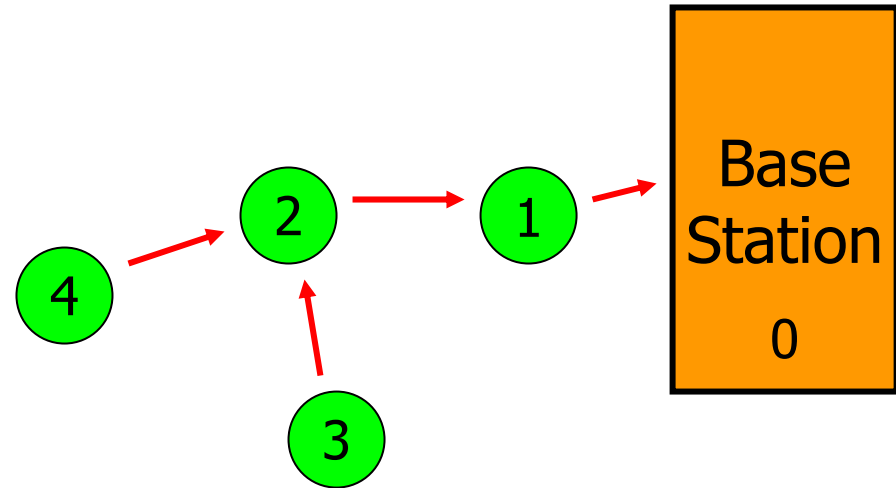


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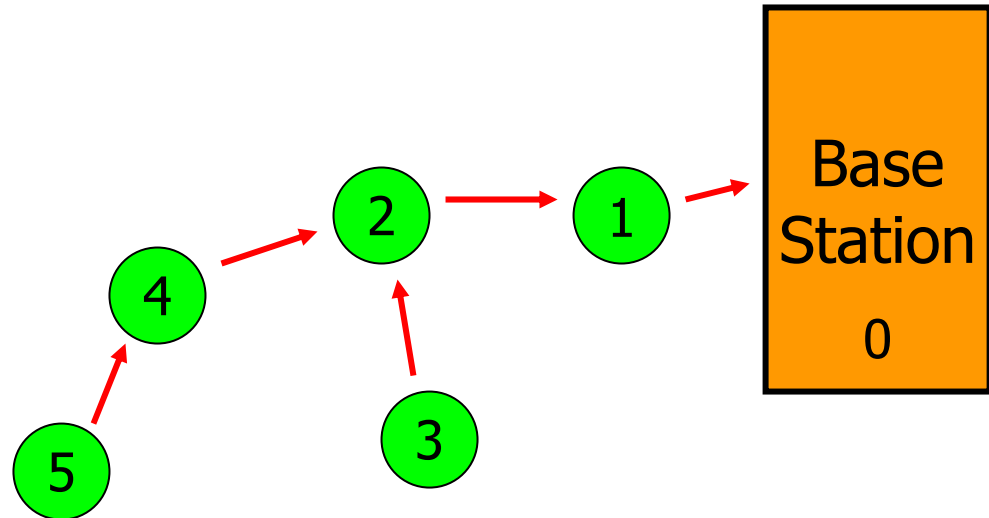


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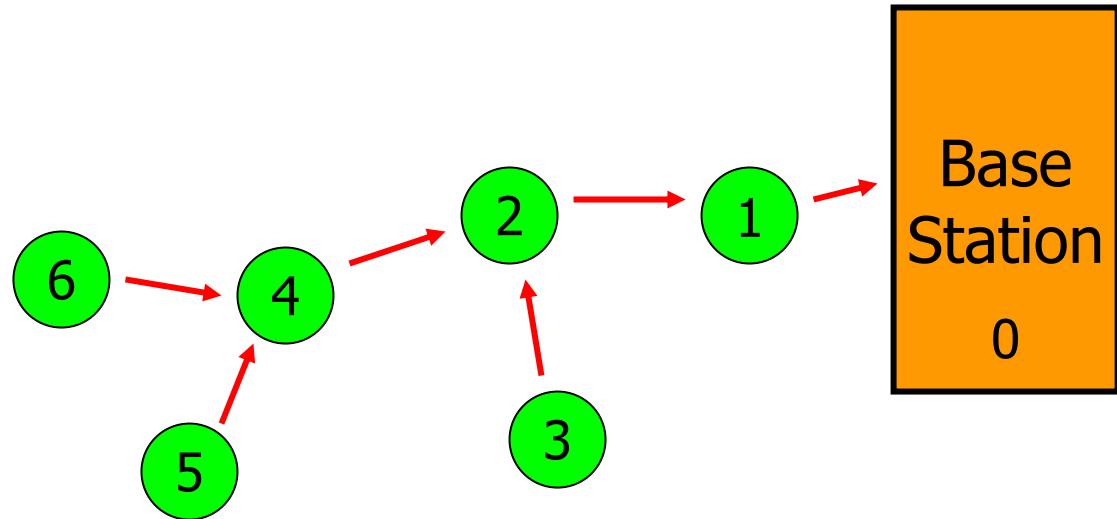


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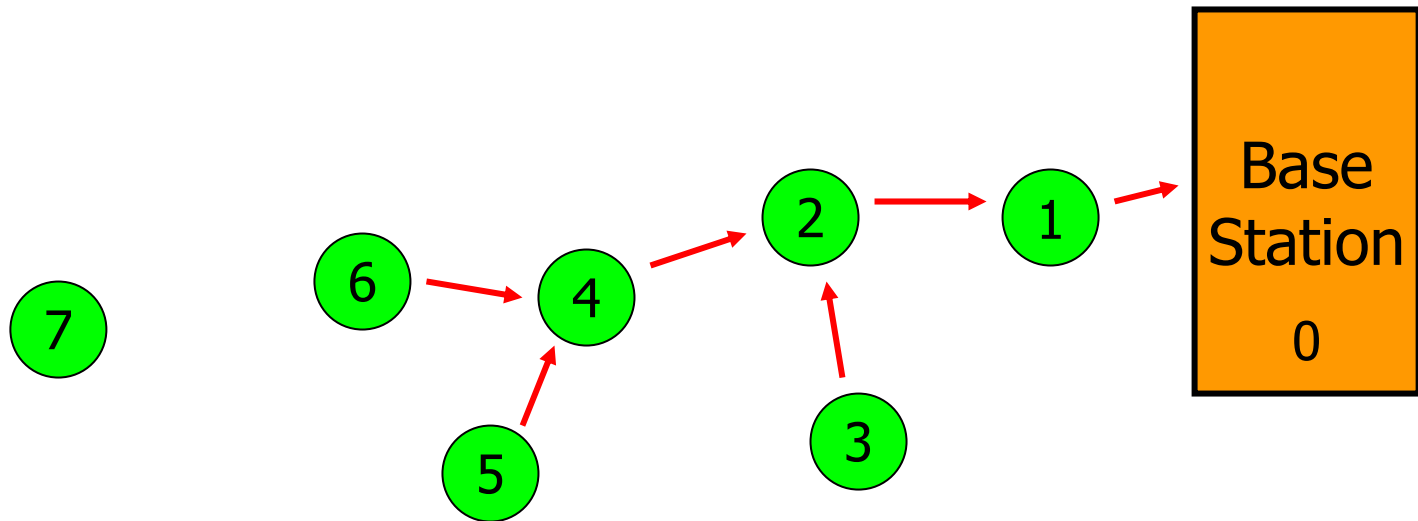


Network Initialization





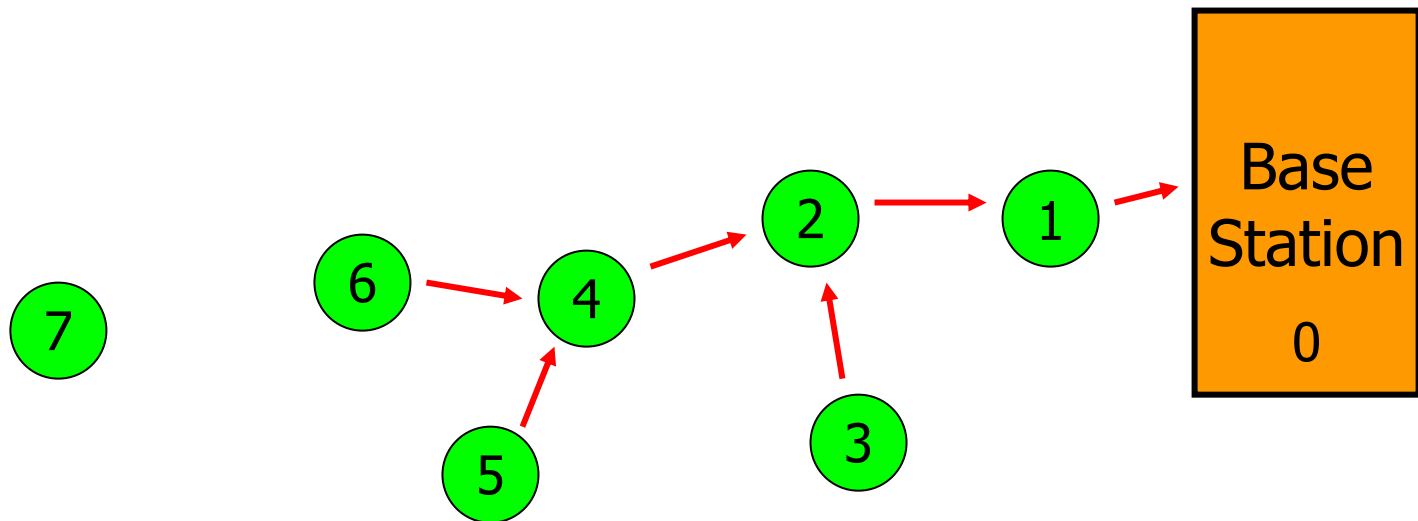
Network Initialization





Adaptive Routing Algorithm

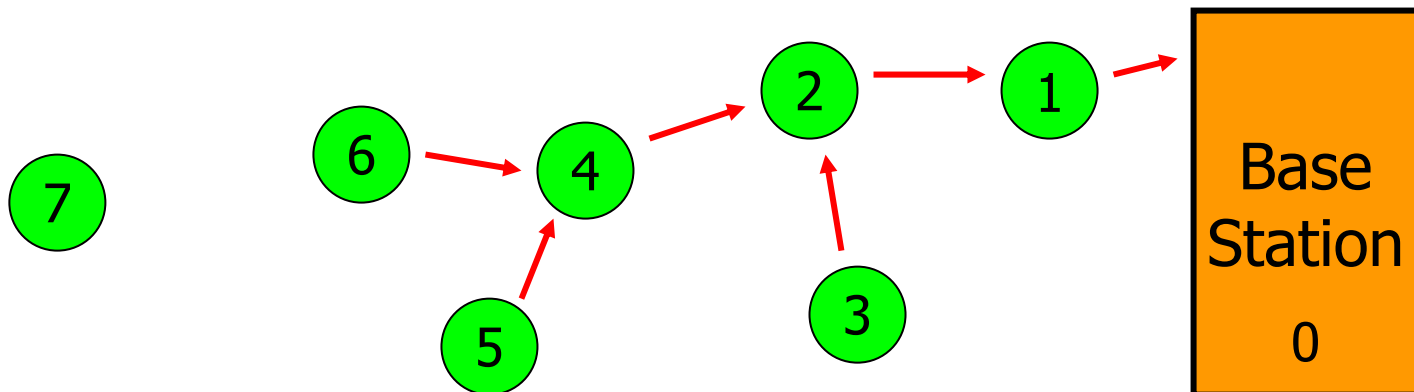
- Base station decides that N nodes must sense
- Variant of N Participants (Information Theory)
- Invariant: sense only if all children sensing





Inheriting Orphaned Subtrees

- Leaf failure is rectified via N Participants
- If routing node fails:
 - Each node independently emits search ping
 - Connected nodes within distance R become prospective parents
 - Orphans choose parent at minimum depth in tree



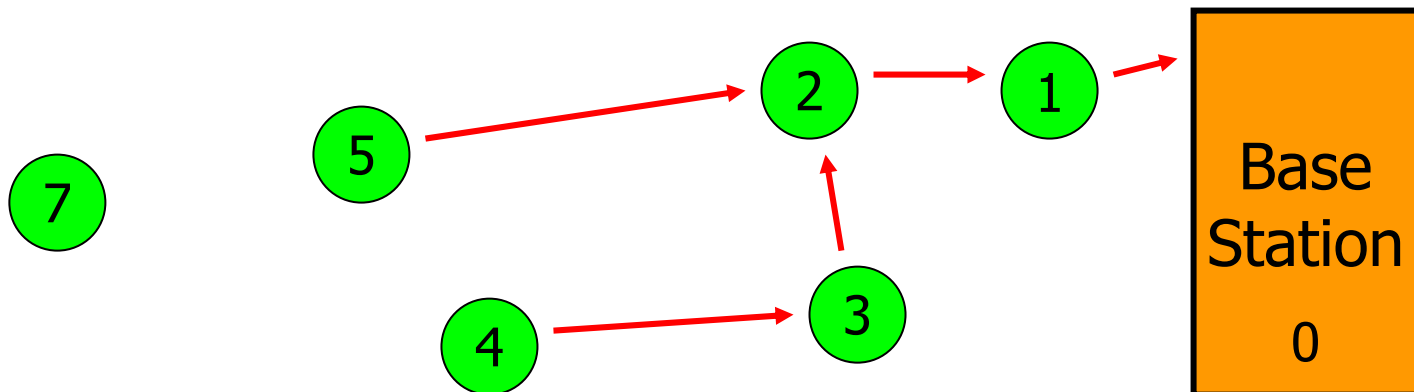


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- The diagram illustrates a network topology for a shortest path algorithm. It features a Base Station (0) represented by an orange rectangle on the right. A red arrow points from the Base Station to node 1 (a green circle). Node 1 is connected to node 2 (a green circle) by a red arrow. Node 2 is connected to node 3 (a green circle) by a red arrow. To the left of node 2, there are two isolated green circles labeled 7 and ?. The entire diagram is set against a light blue background.



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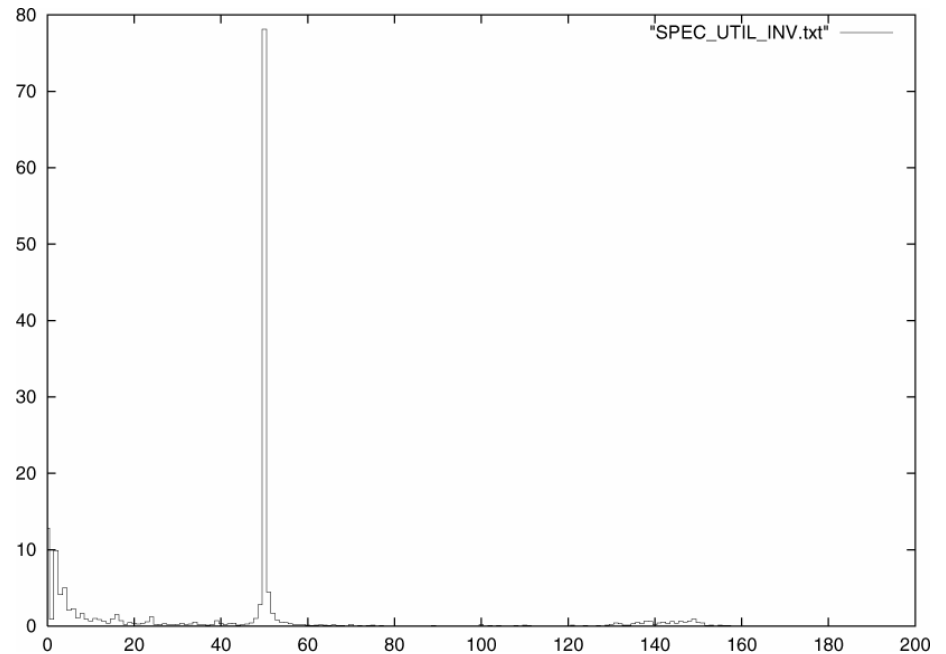


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Simulated Results



Heuristic	Lifetime	Total Utility	Total Energy Consumed	Utility/Energy
NAIVE	92.98	62.3	22039.3	.002827
SIMPLE UTIL	349.633	64.855	15418.01	.004206
SPEC_INV	216	104.87	14505.54	.00723

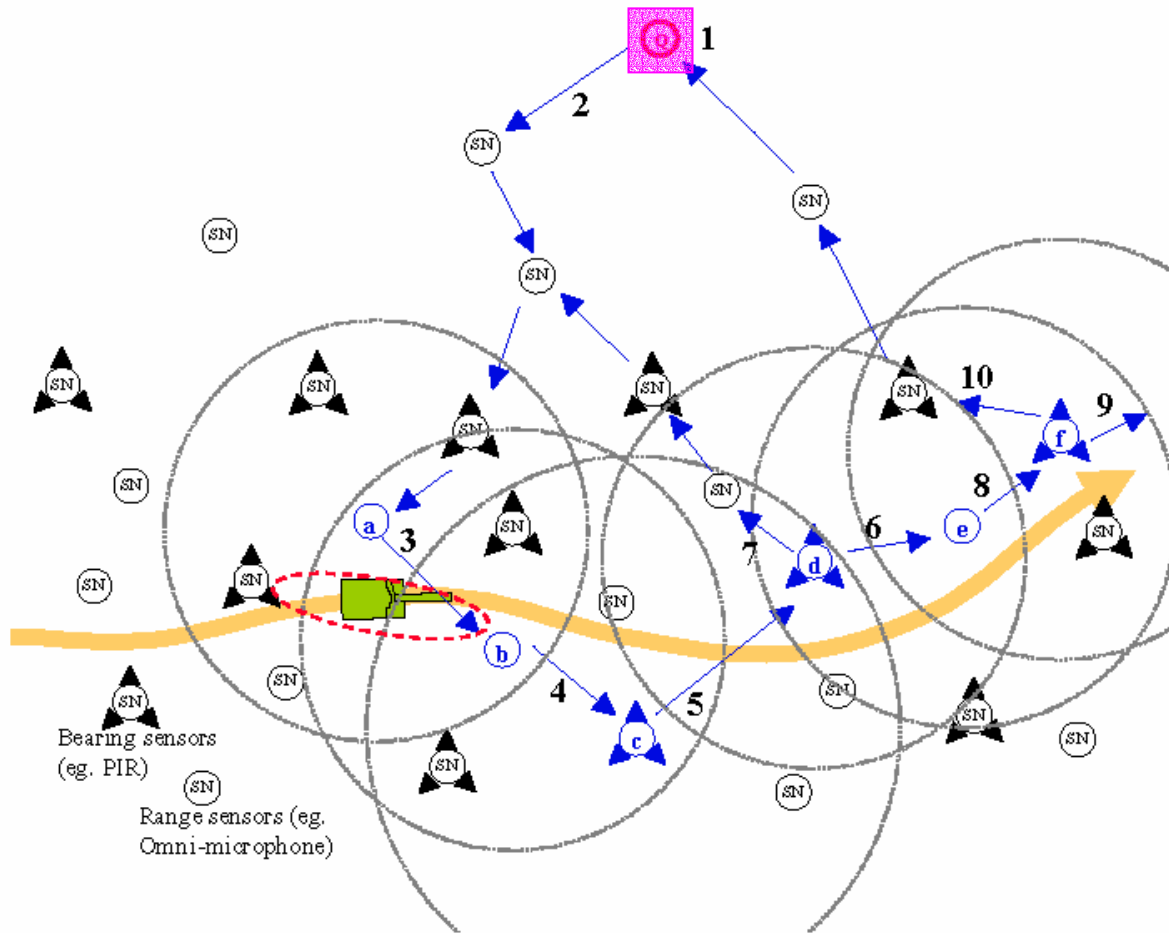


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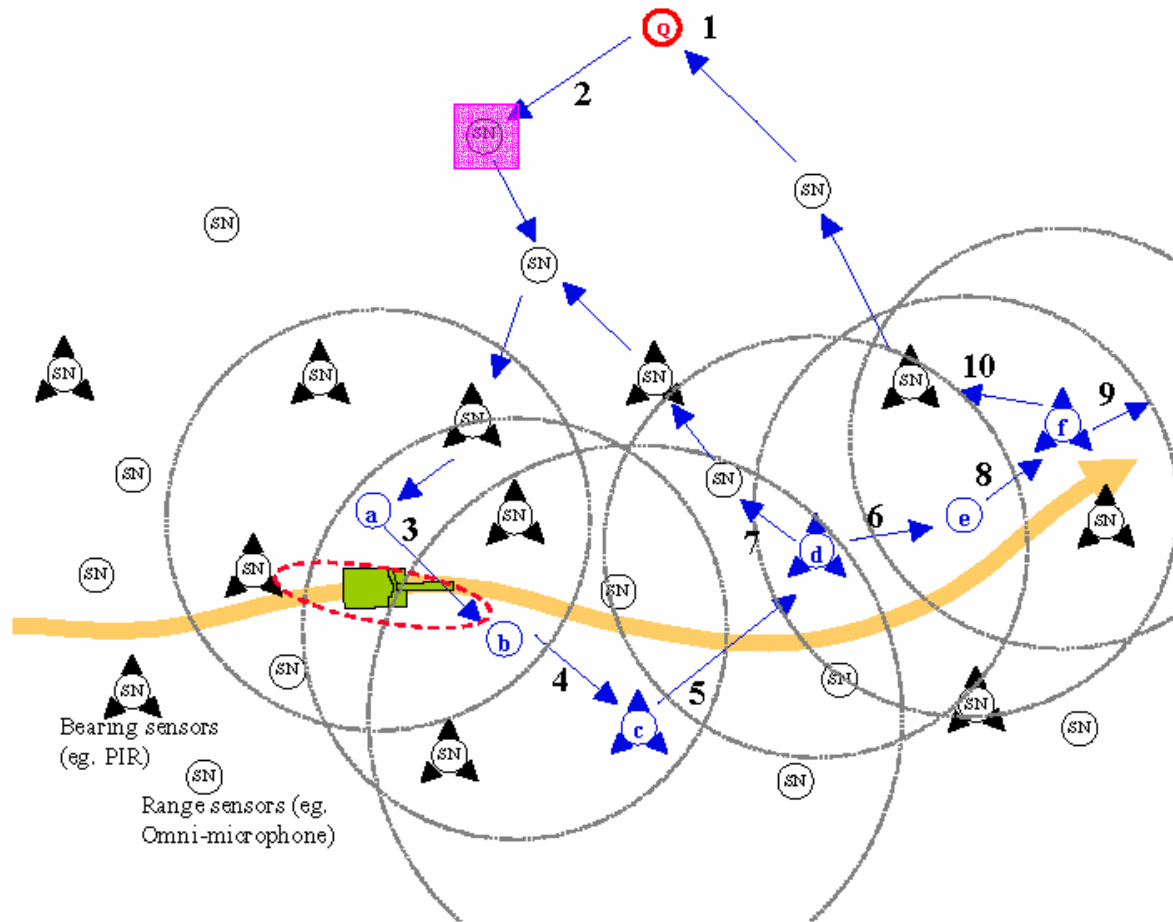
Tracking Scenario



- Query must be routed to the node best able to answer it

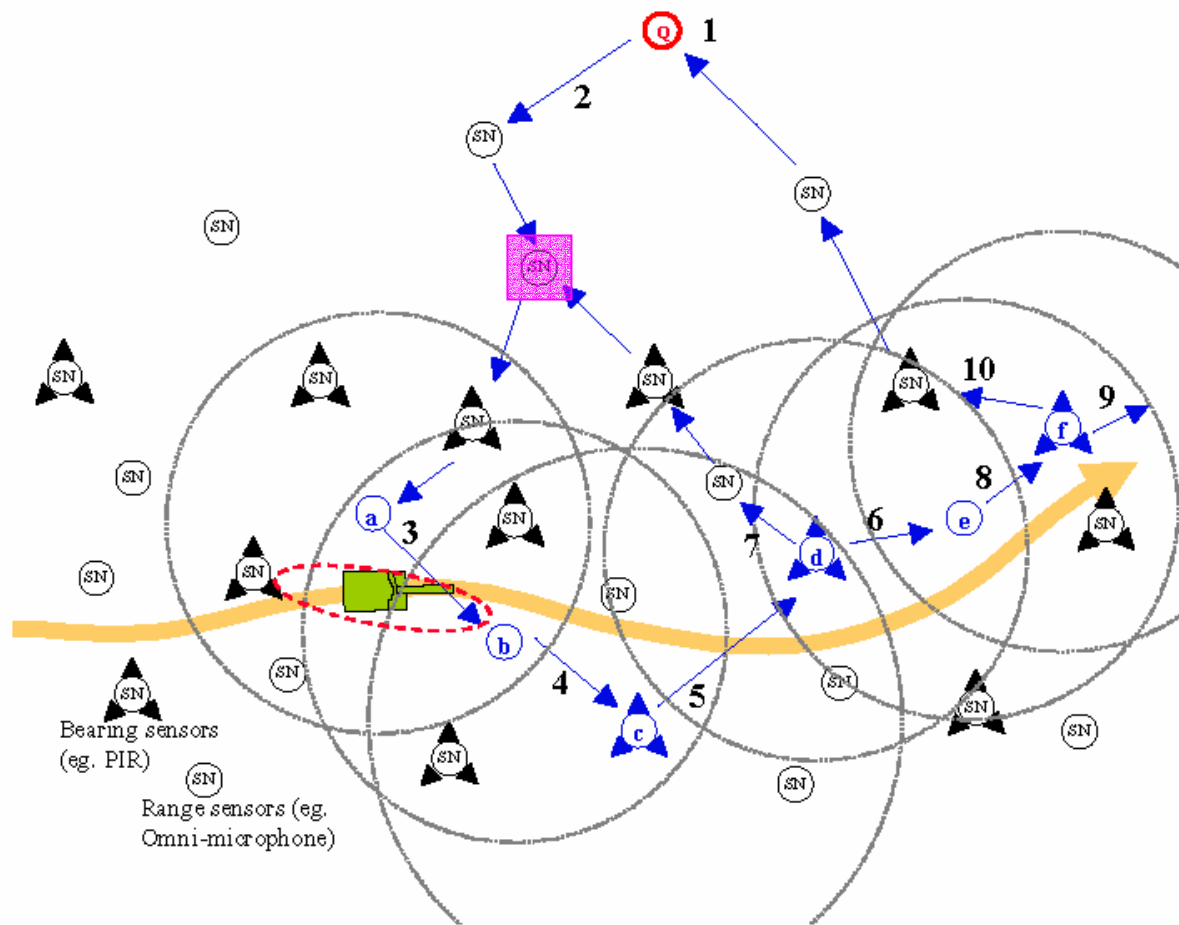


Tracking Scenario



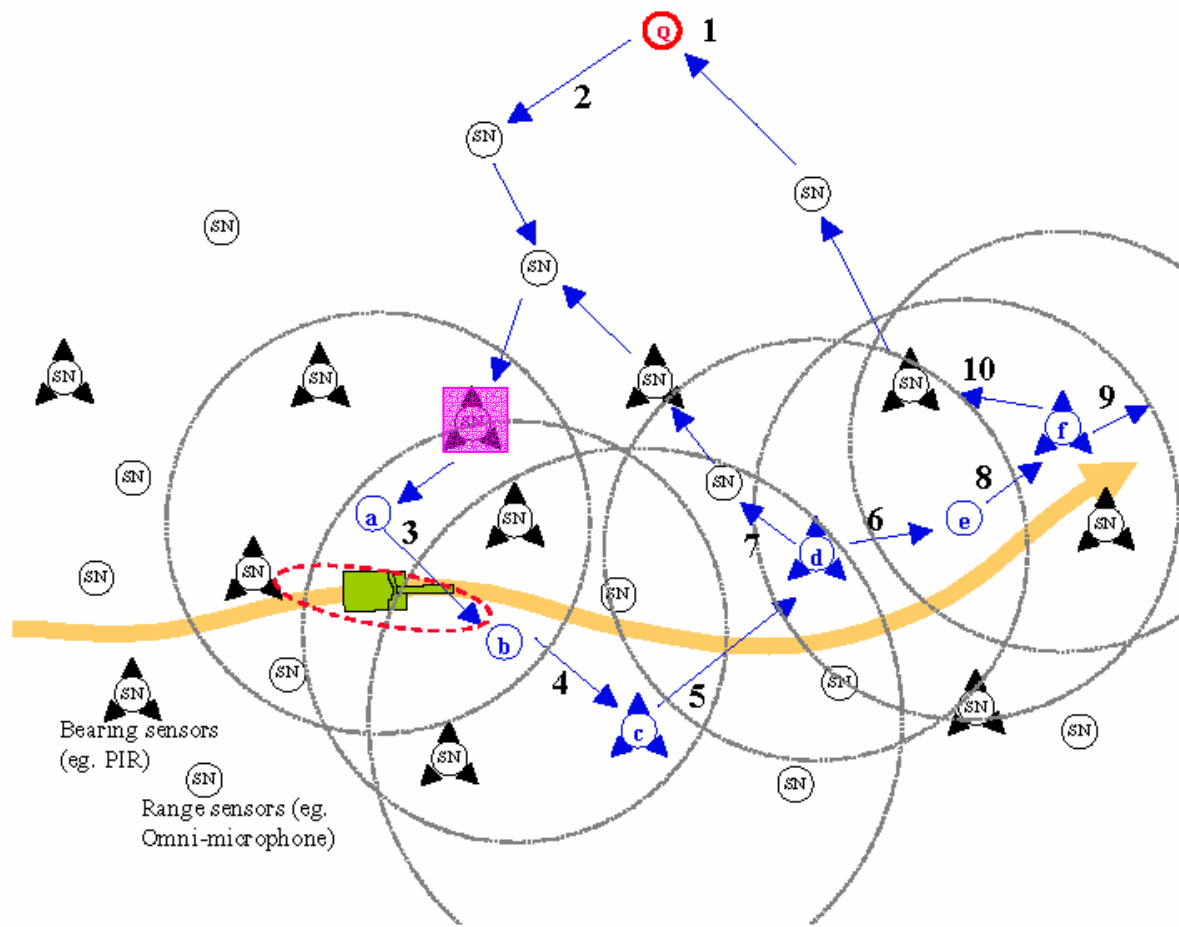


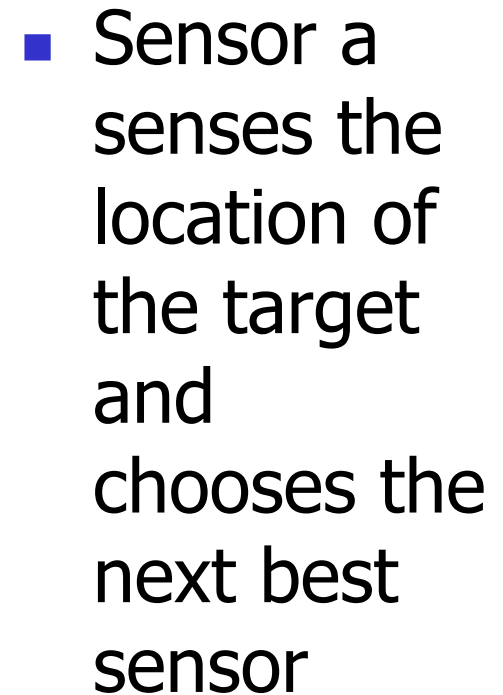
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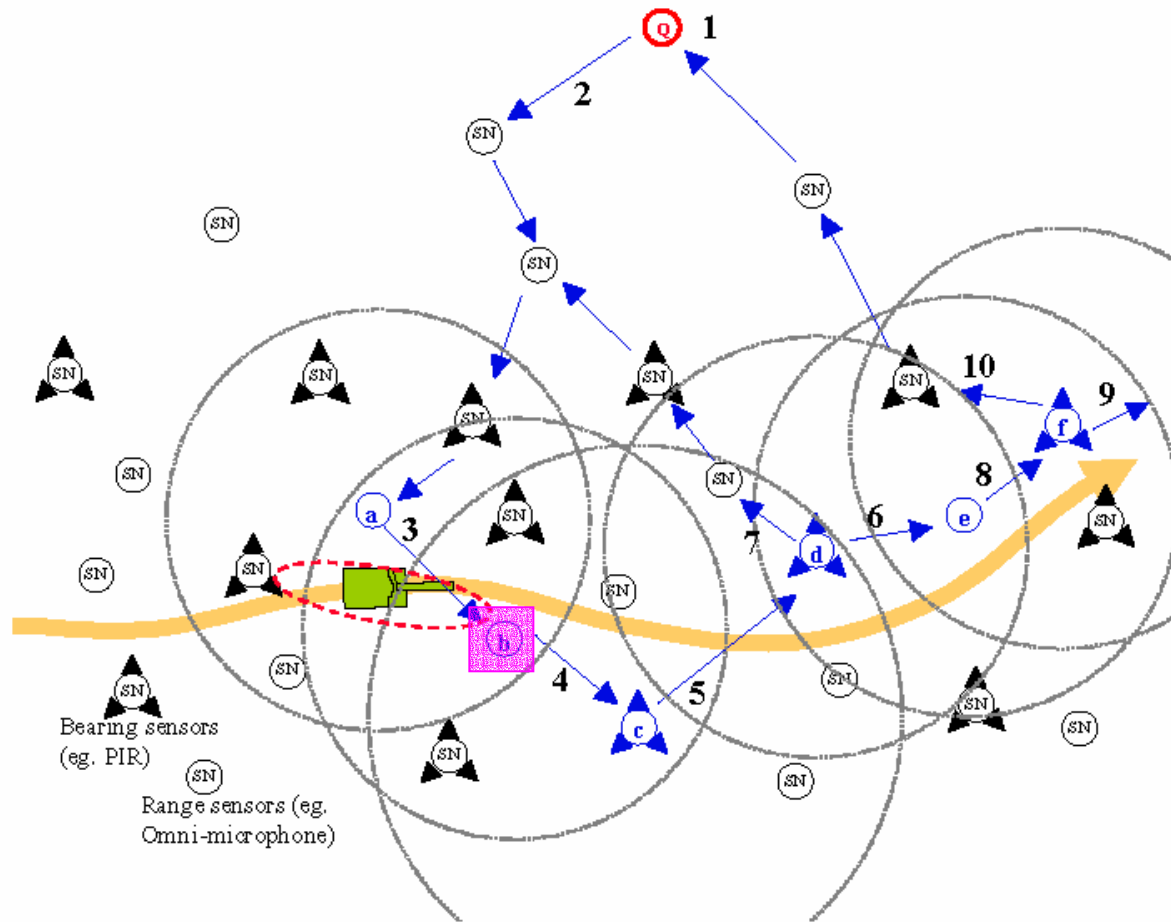
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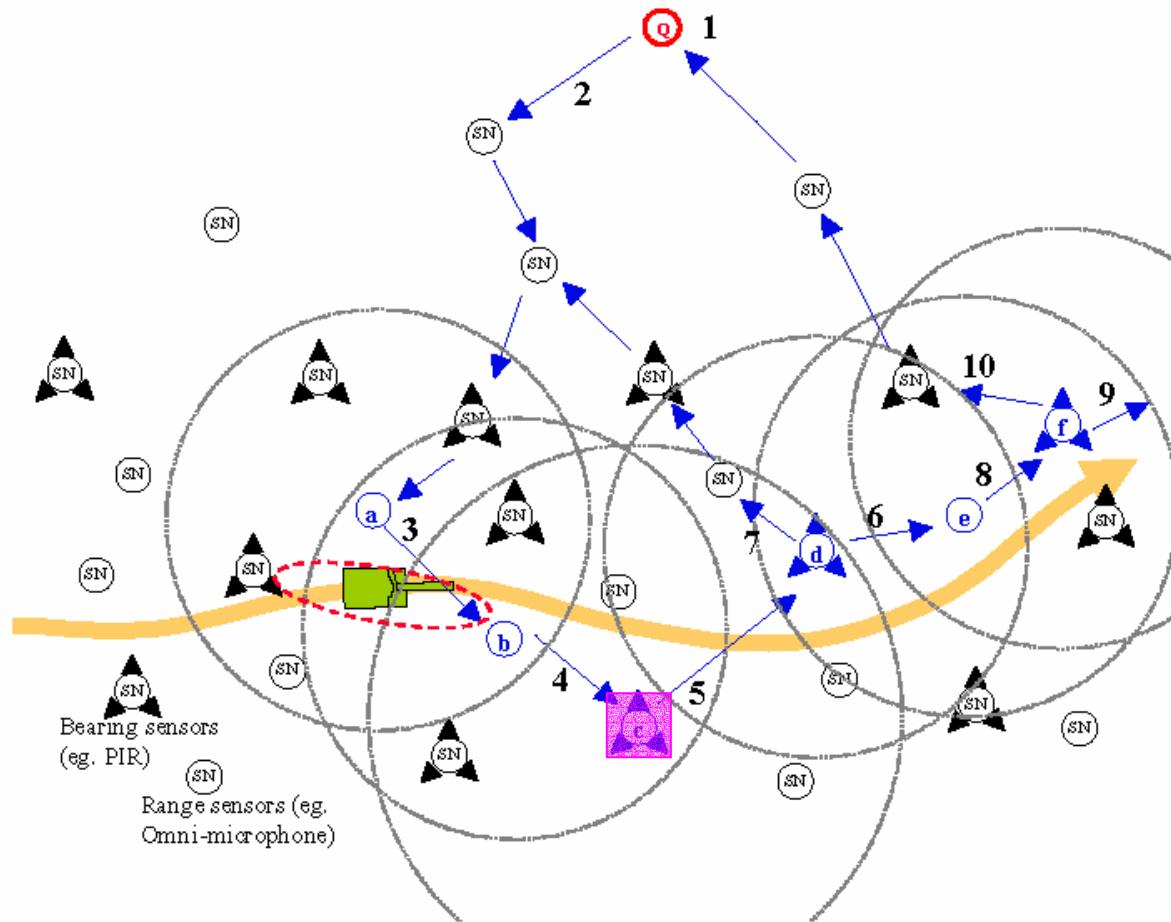
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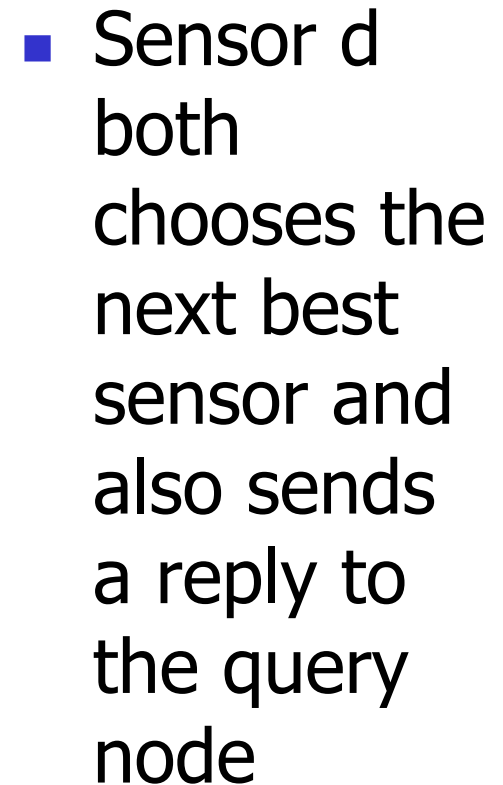


- Sensor b does the same



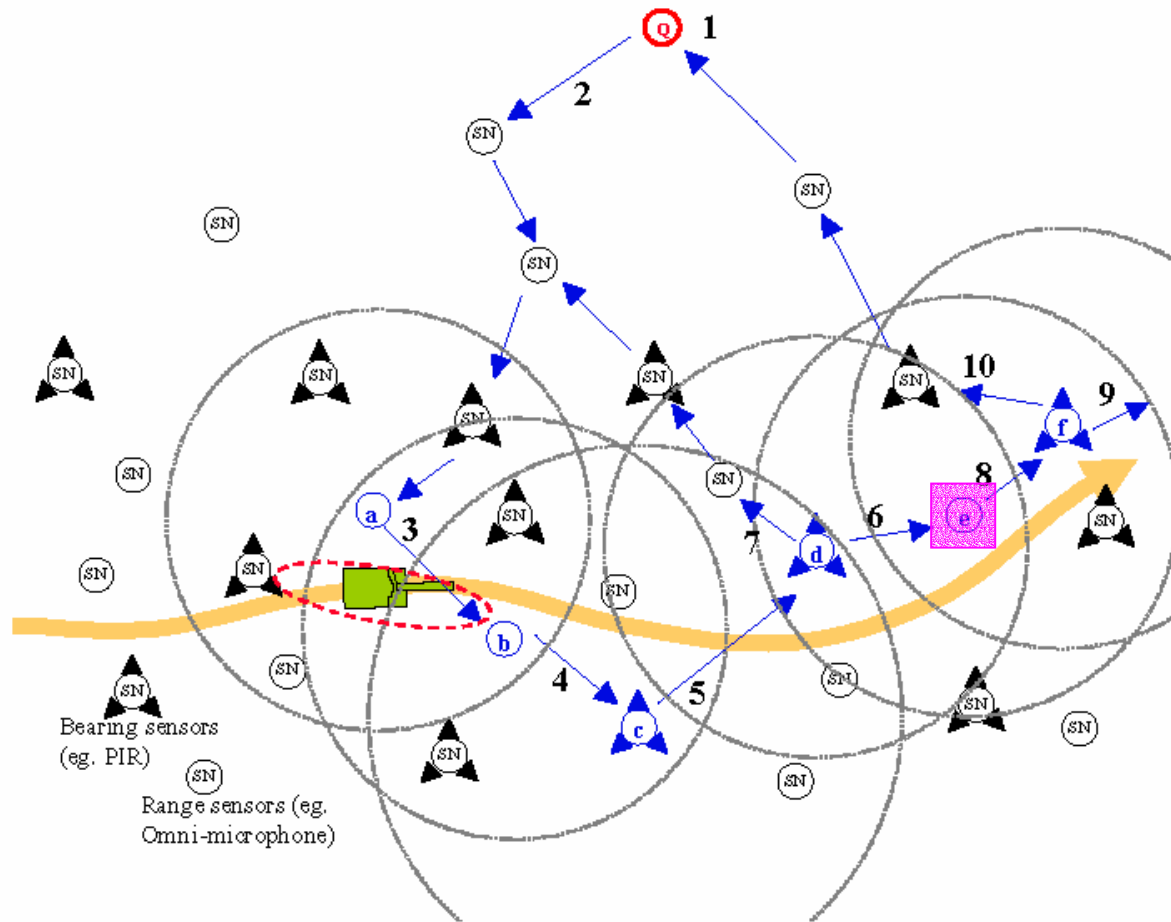
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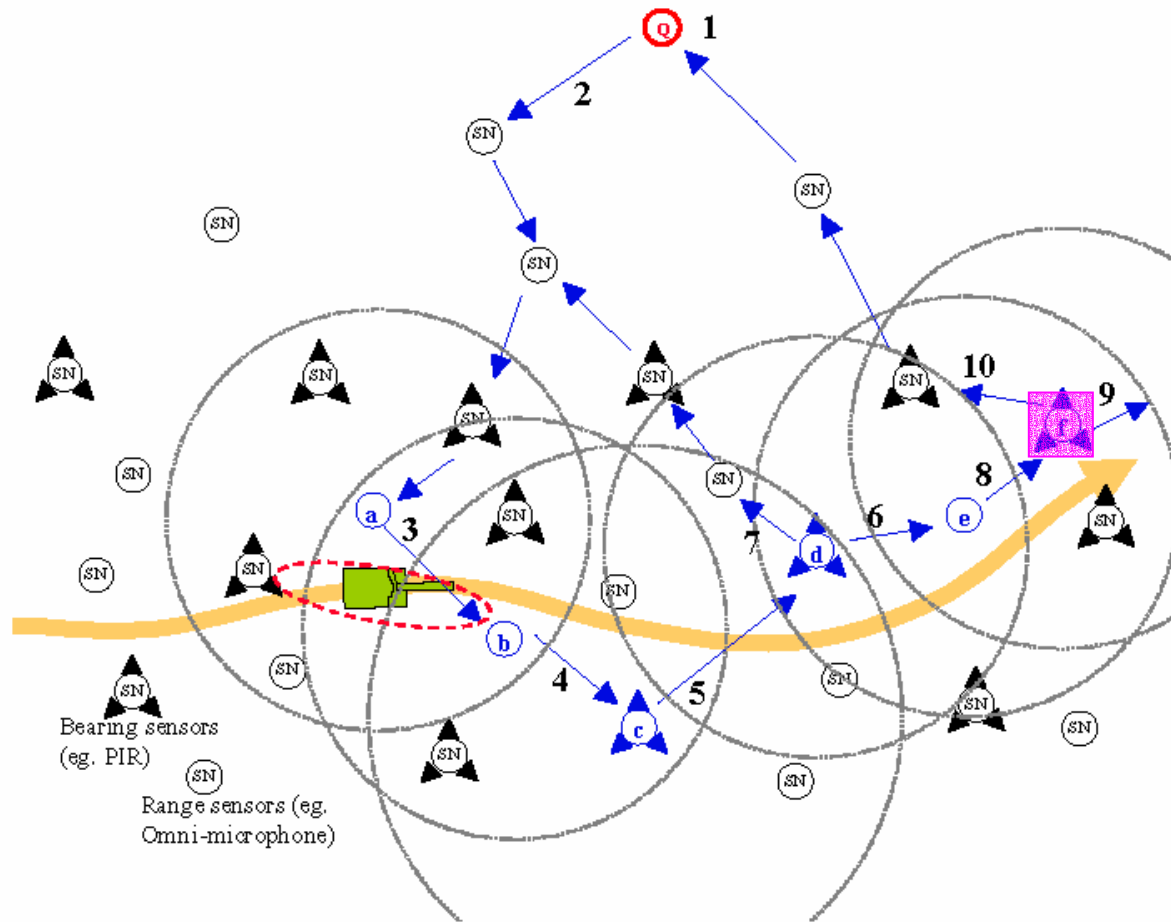


Tracking Scenario





Tracking Scenario



- Sensor f loses the target and sends the final response back to the query node



IDSQ: Information Driven Sensor Querying and Data Routing

- Algorithm to repeatedly choose next best sensor
- Objective:

Cost of Acquiring z_j

$$M(p(\mathbf{x} | \mathbf{z}_1, \dots, \mathbf{z}_j)) = \alpha \cdot \varphi_{Utility}(p(\mathbf{x} | \mathbf{z}_1, \dots, \mathbf{z}_j)) - (1 - \alpha) \cdot \varphi_{Cost}(\mathbf{z}_j)$$

Net Value of Acquiring z_j

Value of Acquiring z_j

Relative weighting of utility and cost

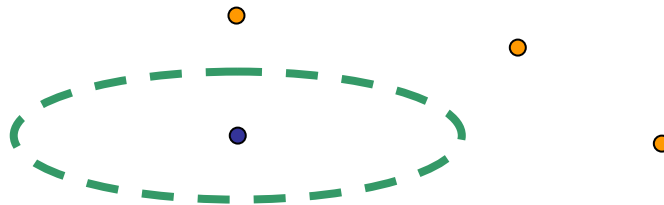


Aside: Mahalanobis Distance

- The quantity r in

$$r^2 = (\mathbf{x} - \mathbf{m}_x)' \mathbf{C}_x^{-1} (\mathbf{x} - \mathbf{m}_x)$$

is called the Mahalanobis distance from the feature vector \mathbf{x} to the mean vector \mathbf{m}_x , where \mathbf{C}_x is the covariance matrix for \mathbf{x}



(all orange points are Mahalanobis equidistant)



Information Utility

- Target location estimated from observations z_1 through z_{j-1}
- Choose next observation providing maximal information utility when incorporated into belief:

$$\hat{j} = \operatorname{argmax}_{j \in V} \varphi_{Utility}(p(\mathbf{x} | \{\mathbf{z}_i\}_{i \in U} \cup \{\mathbf{z}_j\}))$$

Value of Information

New Information

Current Information



Candidate Utility 1: Relative Entropy

- Information utility:

$$\varphi(p(\mathbf{x} \mid \{\mathbf{z}_i\}_{i \in U} \cup \{\mathbf{z}_j\})) \triangleq -H_p(x)$$

- Quick Definitions
 - Entropy is uncertainty
 - Relative Entropy is difference in entropy
- Best metric but impossible to implement



Candidate Utility 2: Mahalanobis Distance

- Information utility:

$$\varphi(\mathbf{x}_j, \hat{\mathbf{x}}, \hat{\Sigma}) = -(\mathbf{x}_j - \hat{\mathbf{x}})^T \hat{\Sigma}^{-1} (\mathbf{x}_j - \hat{\mathbf{x}})$$

Sensor Position

Mean, Covariance of Belief

- Works well for Gaussian distributions
- Does not generalize well to heterogeneous sensors



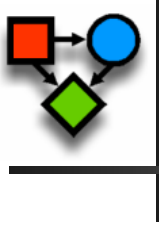
Candidate Utility 3: Expected Relative Entropy

- Observation at sensors unknown
- Need observations for relative entropy
- Posterior gives expected observations
- Useful for finding expected relative entropy

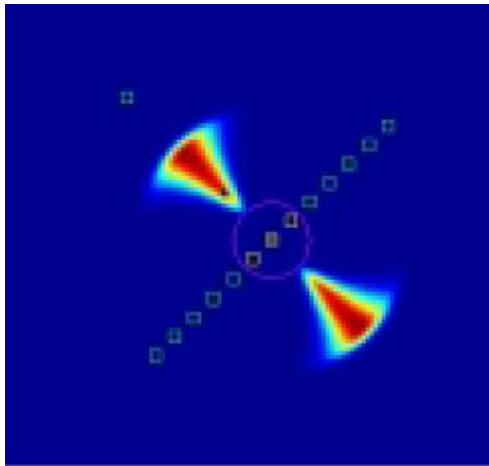


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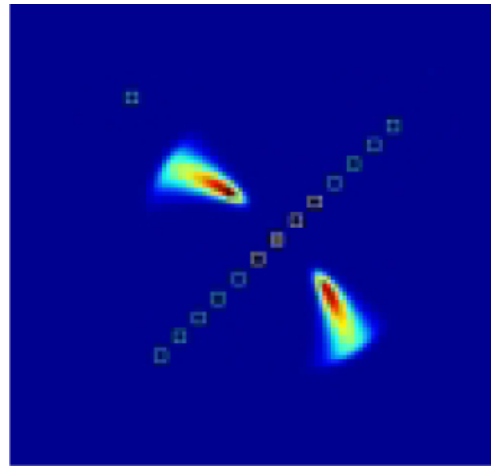
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Experimental Results: Sensor Choice



Original Belief



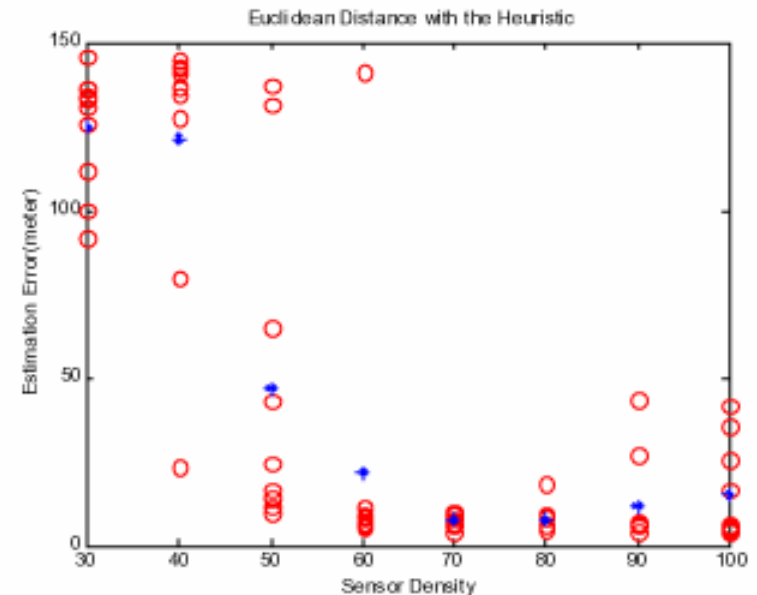
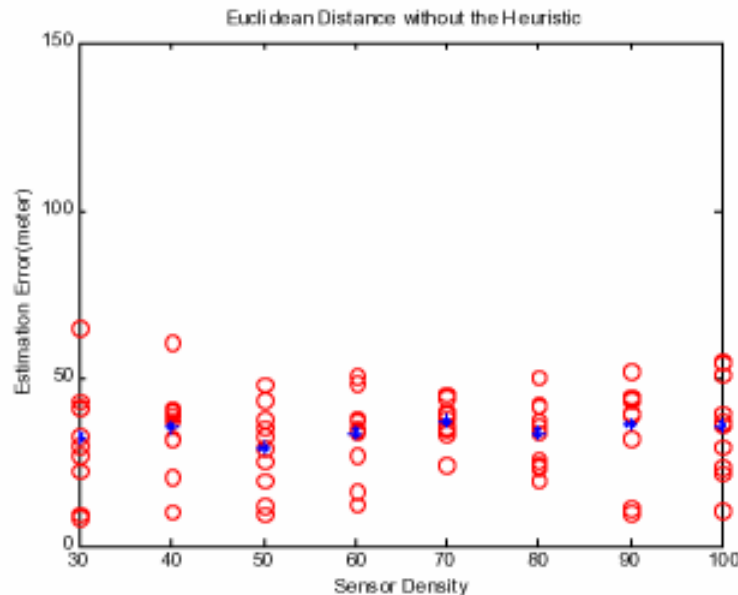
Using Adjacent
Sensor



Using Mahalanobis
Information Utility



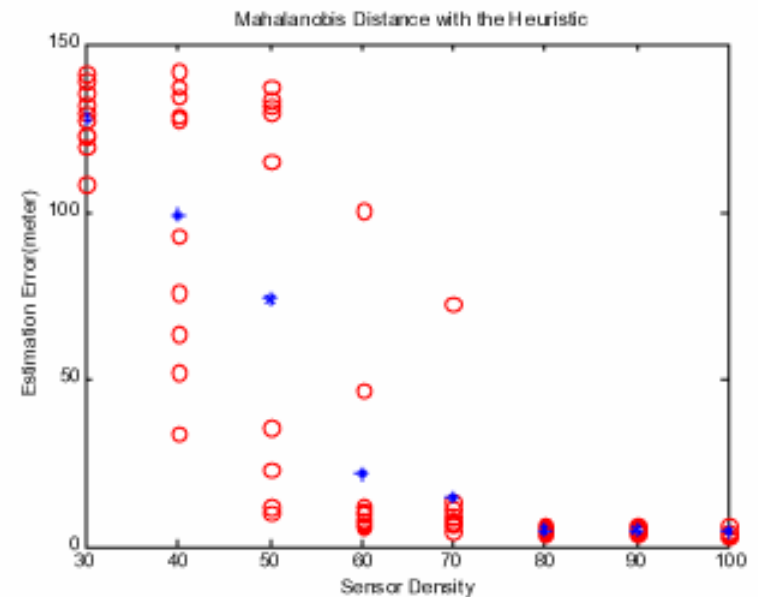
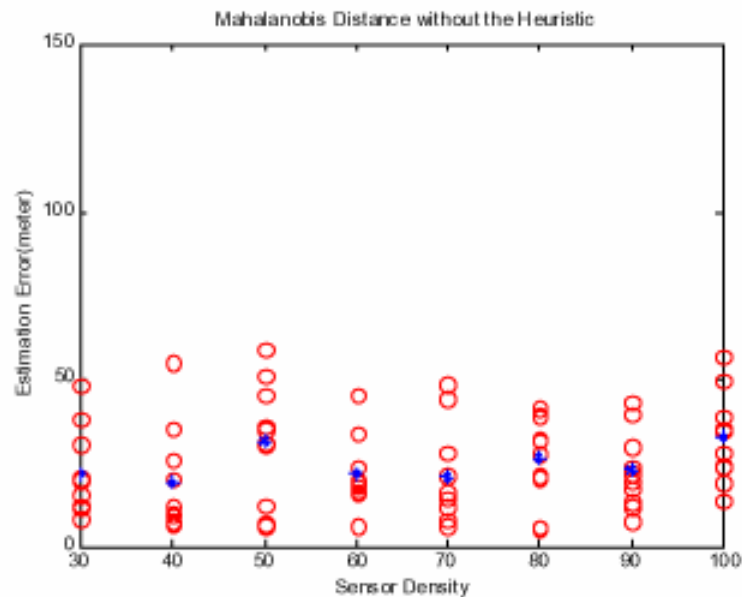
Experimental Results: Euclidean Distance to Mean



(heuristic prevents node from being selected N times)

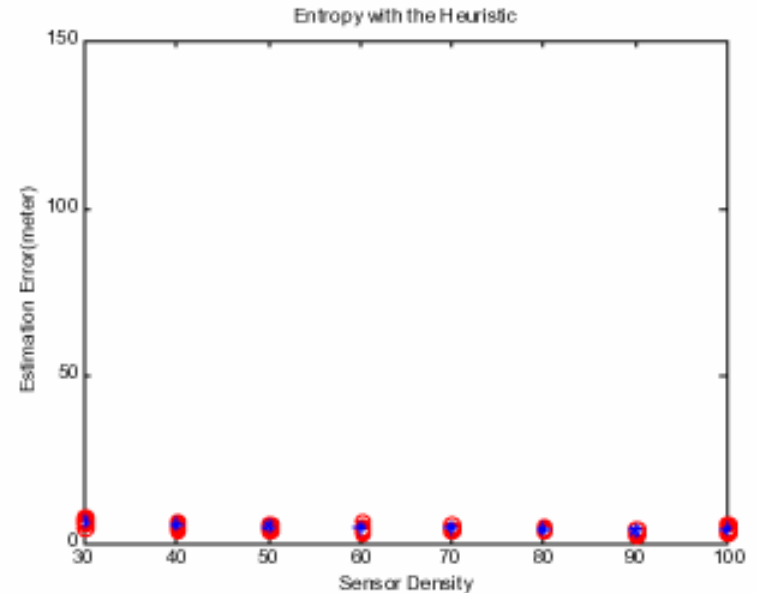
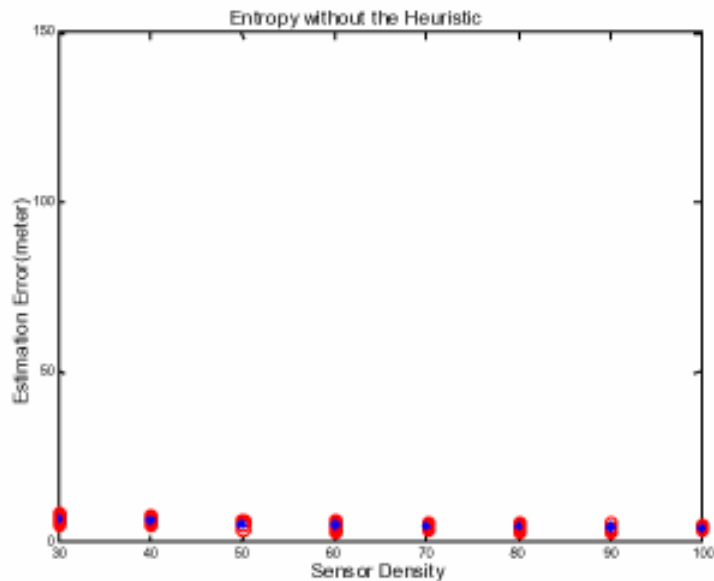


Experimental Results: Mahalanobis Distance to Mean



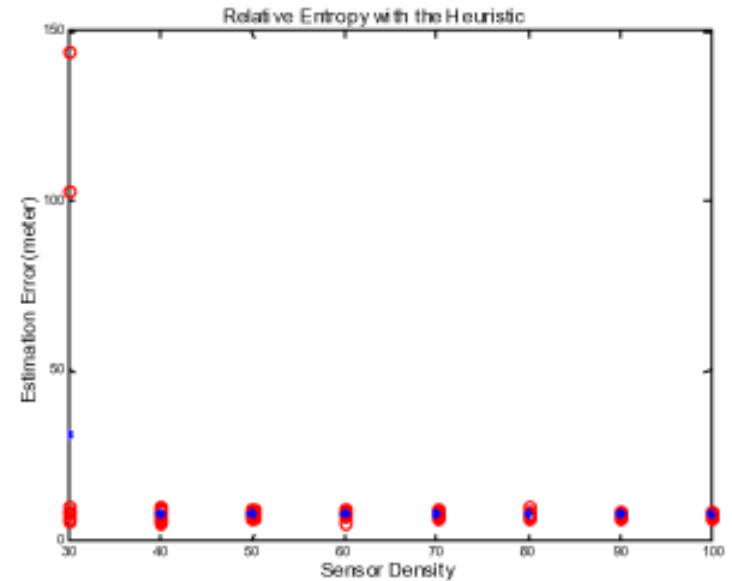
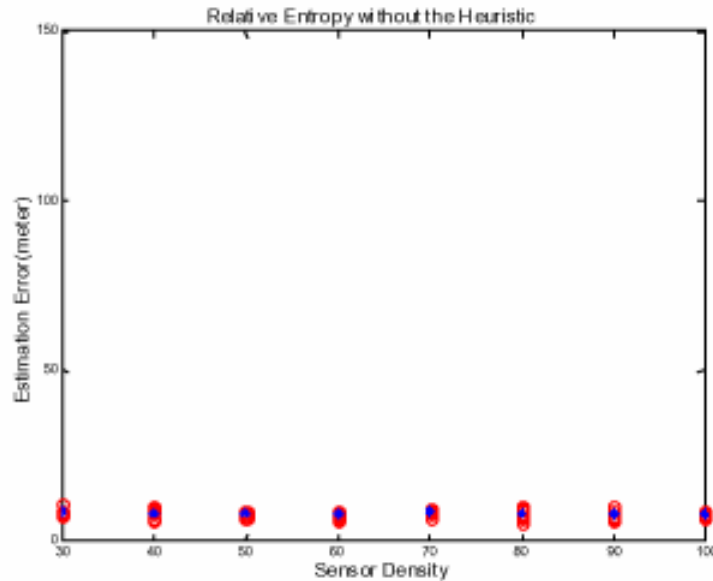


Experimental Results: Entropy (unattainable)





Experimental Results: Relative Entropy





Experimental Results: Tracker Performance

	# Lost tracks/ total runs	Mean error
(a) Nearest neighbor	75/80	34.39
(b) Nearest neighbor with heuristic	37/80	44.79
(c) Mahalanobis distance	70/80	24.86
(d) Mahalanobis distance with heuristic	32/80	44.20
(e) Entropy	0/80	5.13
(f) Entropy with heuristic	0/80	5.05
(g) Relative entropy	0/80	8.09
(h) Relative entropy with heuristic	2/80	10.79



Conclusion

- Power-aware sensor management:
 - Increases sensor lifespan
 - Decreases number of sensors needed
- Information utility:
 - Directs sensing to find more valuable information
 - Balances power consumption and information acquisition