

# Time Synchronization in Sensor Networks

- *Time Synchronization in Ad Hoc Networks* – Kay Roemer

- *Fine Grained Network Time Synchronization using Reference Broadcasts* – J. Elson, L. Girod, D. Estrin

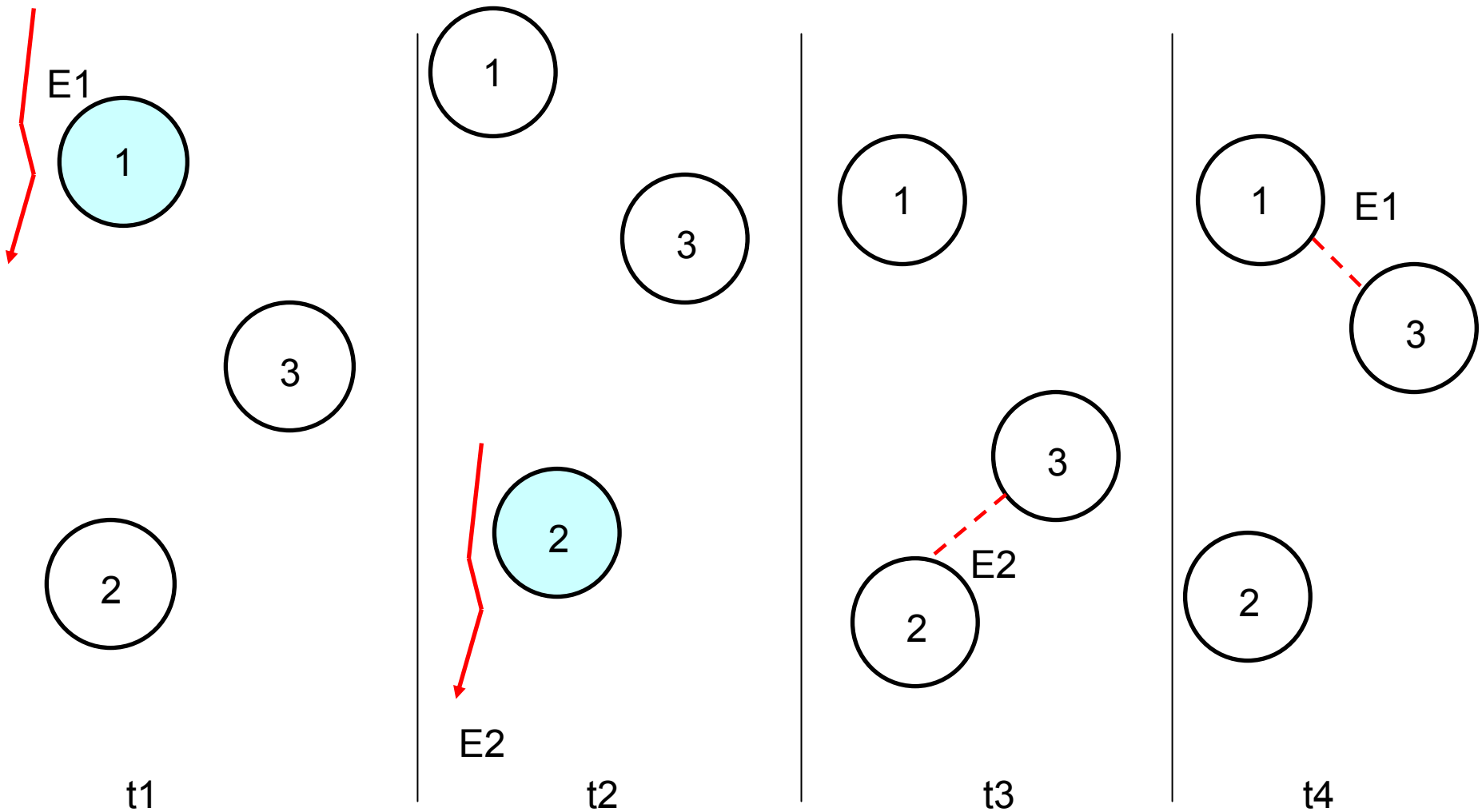
# Need For Time Synchronization

- detect direction of movement of a phenomenon
- form a TDMA radio schedule
- configure a beamforming array
- estimate velocity of a moving phenomenon
- suppress redundant transmissions by recognizing duplicate detection of same event

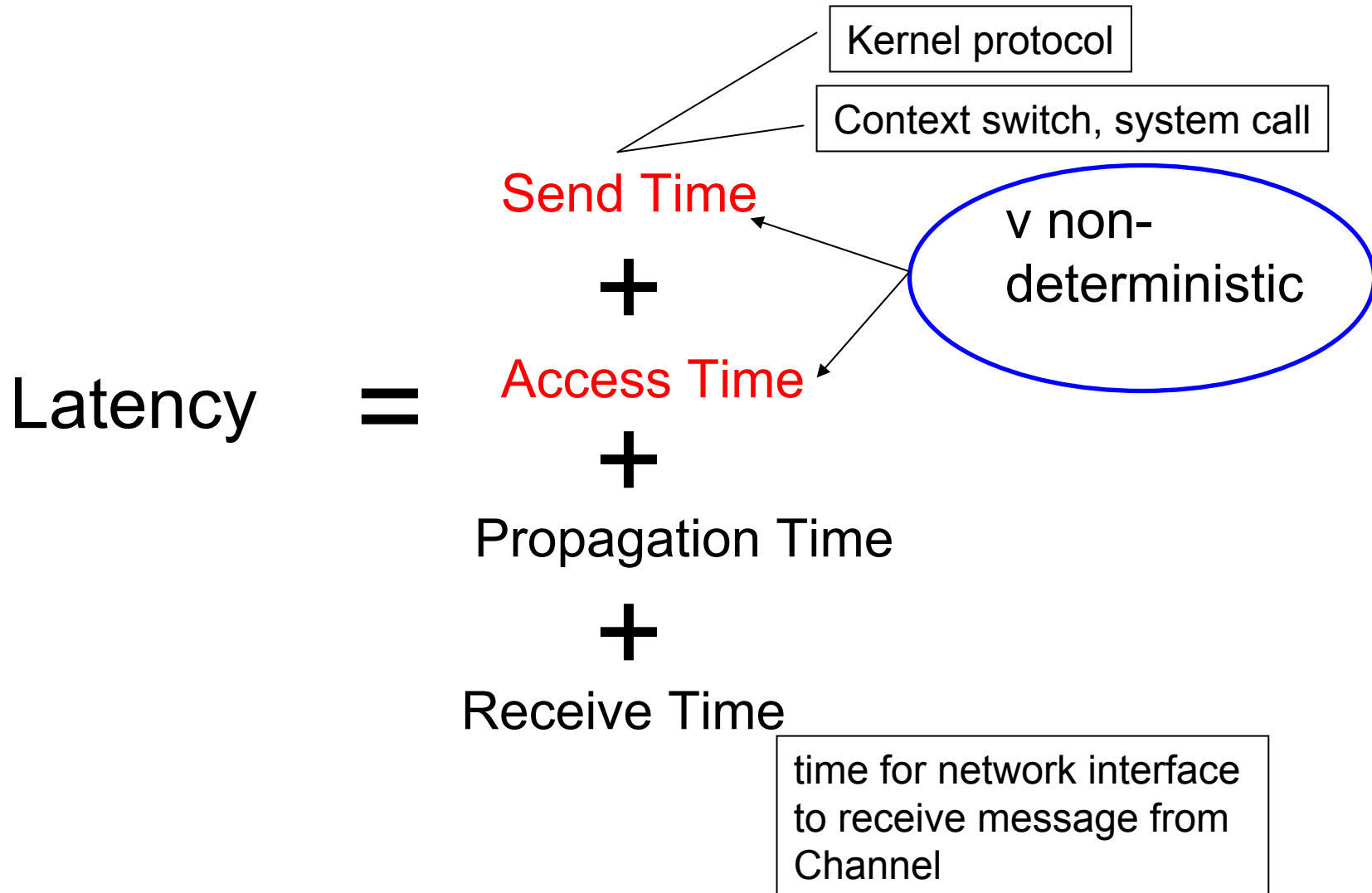
# Traditional Synchronization Methods

- Exchange messages between the nodes
- Estimate one way latency

# Difficulty – Topology Change



# Difficulty – Non Determinism



- *Time Synchronization in Ad Hoc Networks* – Kay Roemer

- attacks **problem of topology change**
- determines upper and lower bounds for real time

- *Fine Grained Network Time Synchronization using Reference Broadcasts* – J. Elson, L. Girod, D. Estrin

- attacks problem of **non determinism in latency**
- uses physical broadcast to synchronize receivers

# K. Roemer – Time Transformation

- Time stamps generated by unsynchronized clocks
- Transform local time of sender to local time of receiver
- Determines lower and upper bounds for such transformations
- Temporal ordering
  - $t(E_1) < t(E_2)$
  - $t(E_1) > t(E_2)$
  - Unsure

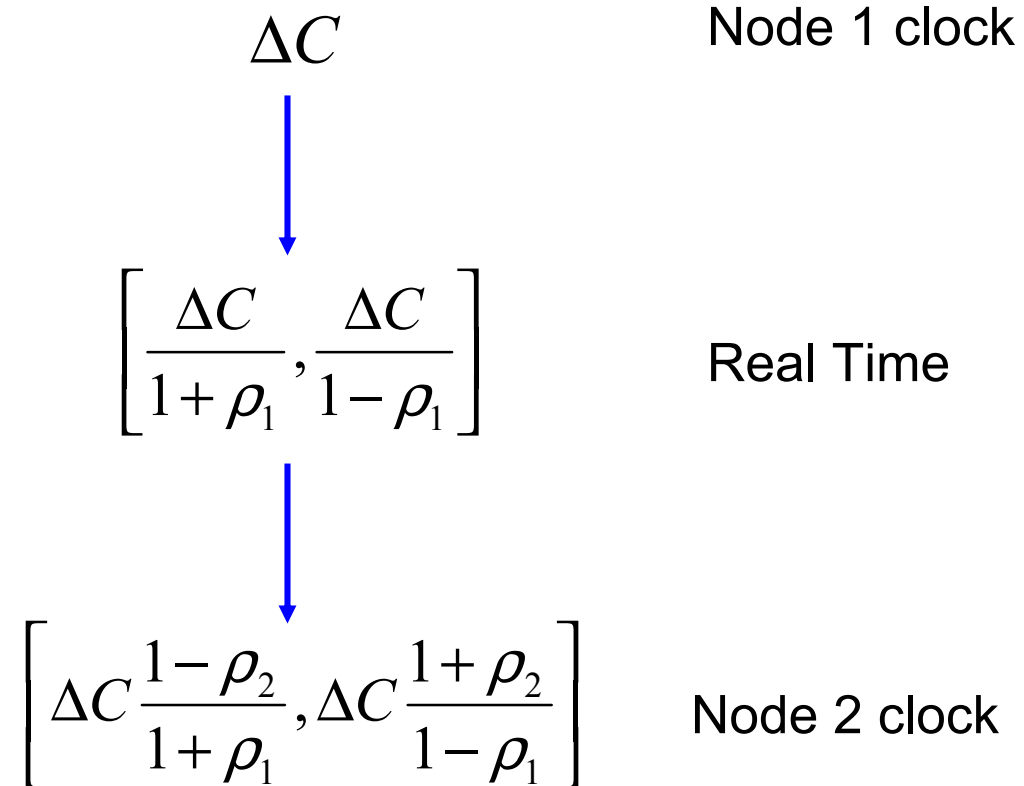
# Clock Drift

$$1 + \hat{u} \approx \frac{dC}{dt} \approx 1 + \hat{u}$$

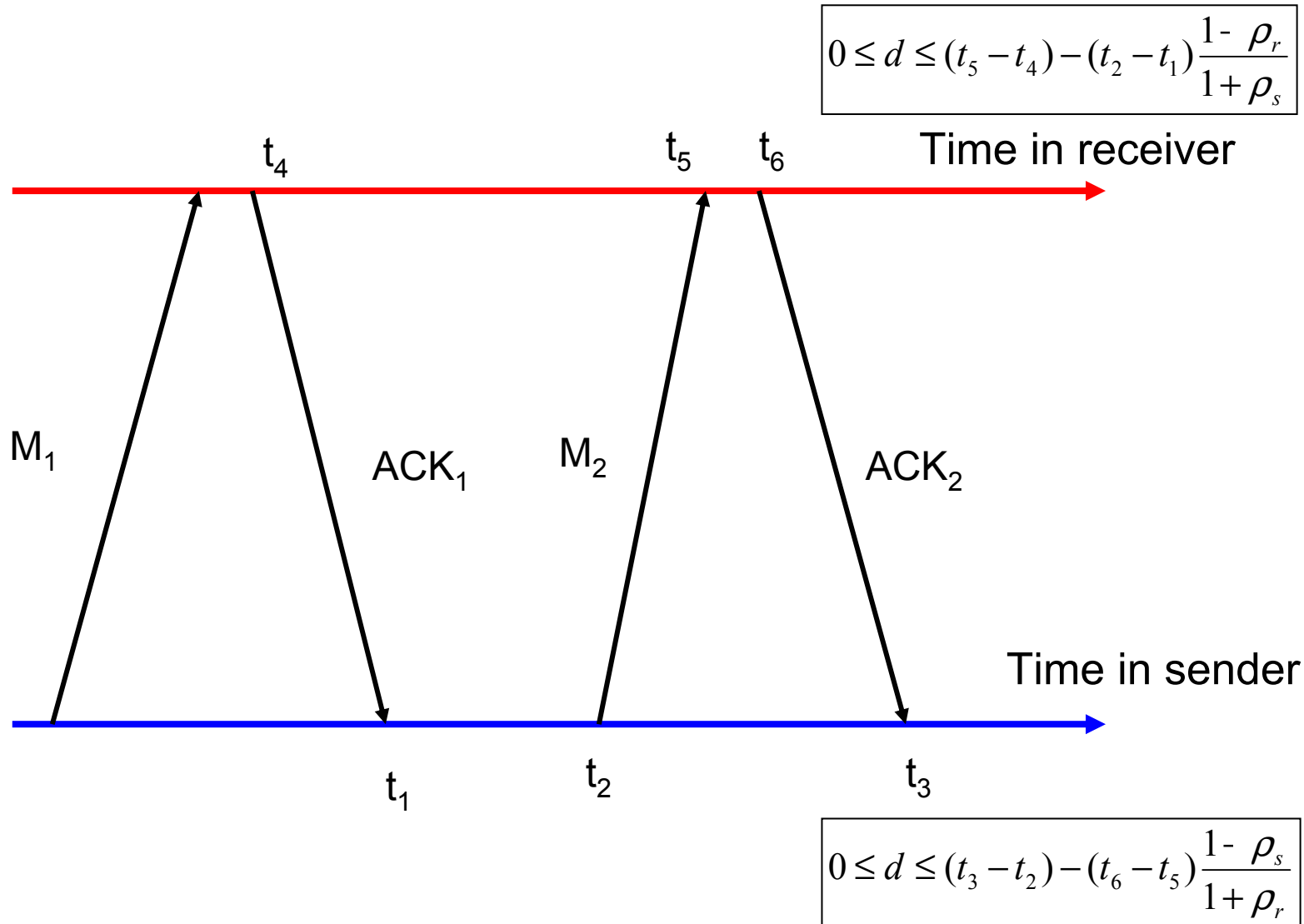
- $\hat{u}$  is of the order of  $10^{-6}$
- different computer clocks have different drifts
- **Approximate** as differences to get
  - $(1 - \rho)\Delta t \leq \Delta C \leq (1 + \rho)\Delta t$
  - $\frac{\Delta C}{1 + \rho} \leq \Delta t \leq \frac{\Delta C}{1 - \rho}$



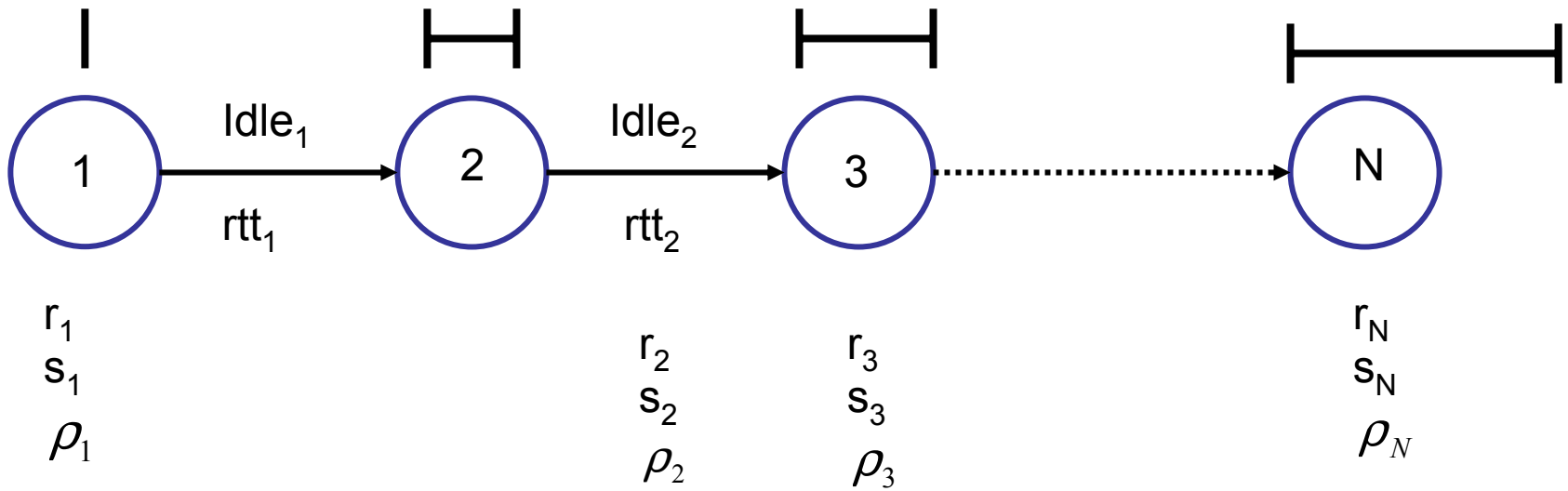
# Time Transformation



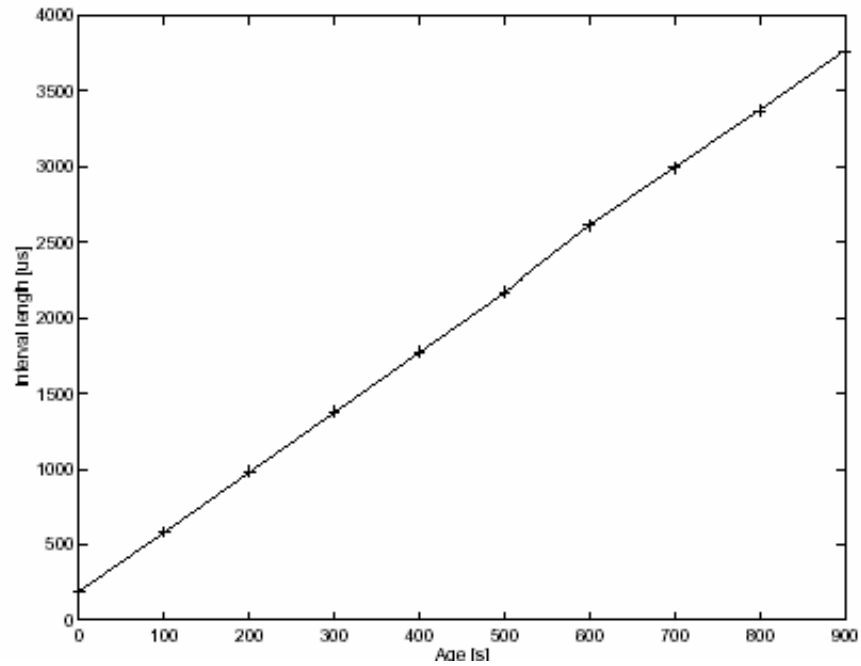
# Message Delay Estimation



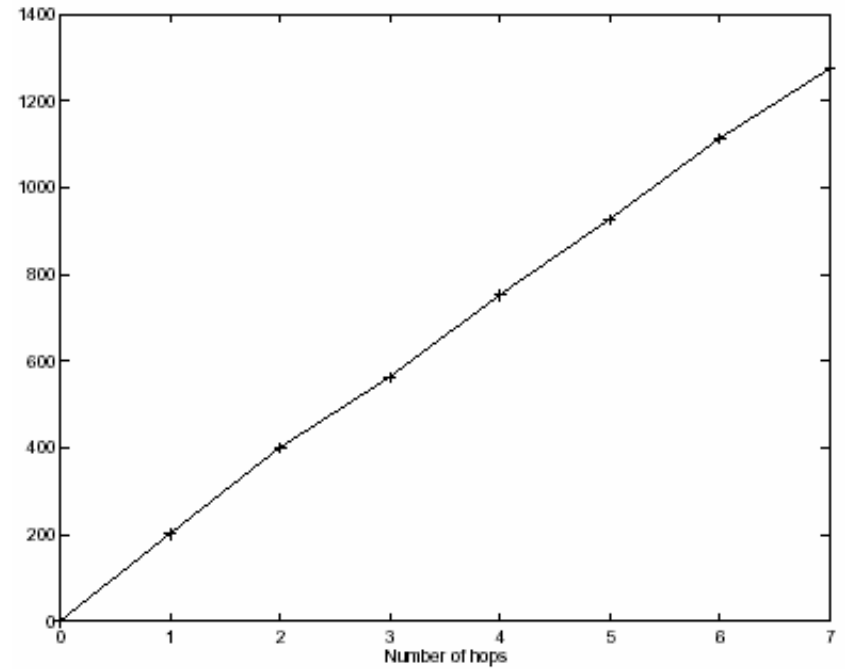
# Message Flow



# Accuracy

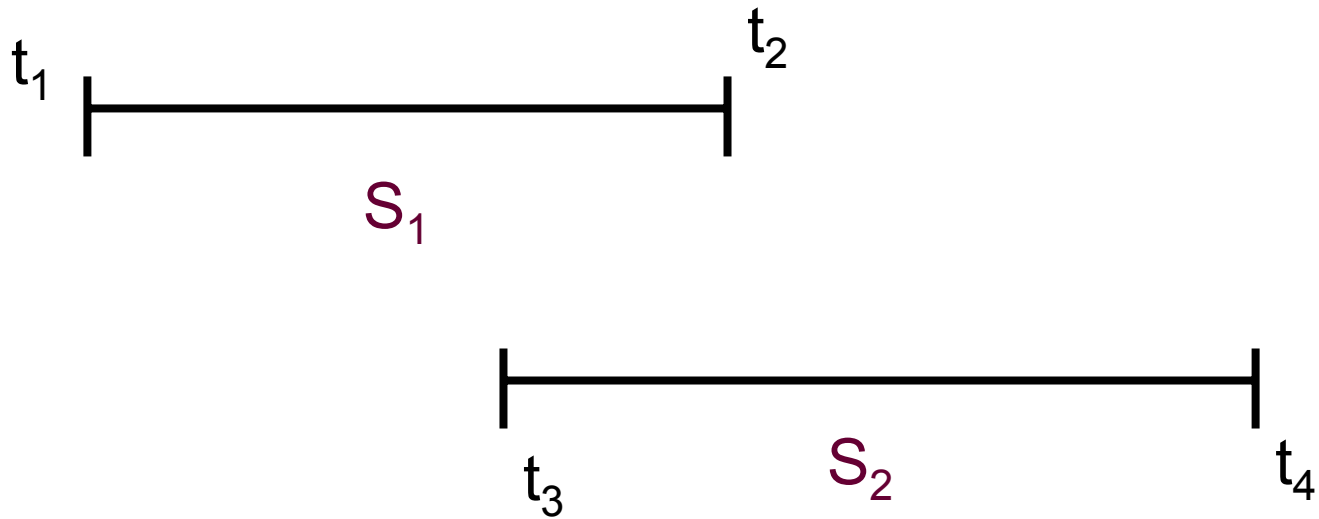


Age



number of hops

# Interval Arithmetic



# Summary

- Handles network partitioning
- Does not need special topologies
- Low message overhead
- Improvements based on probability distribution of uncertainty possible

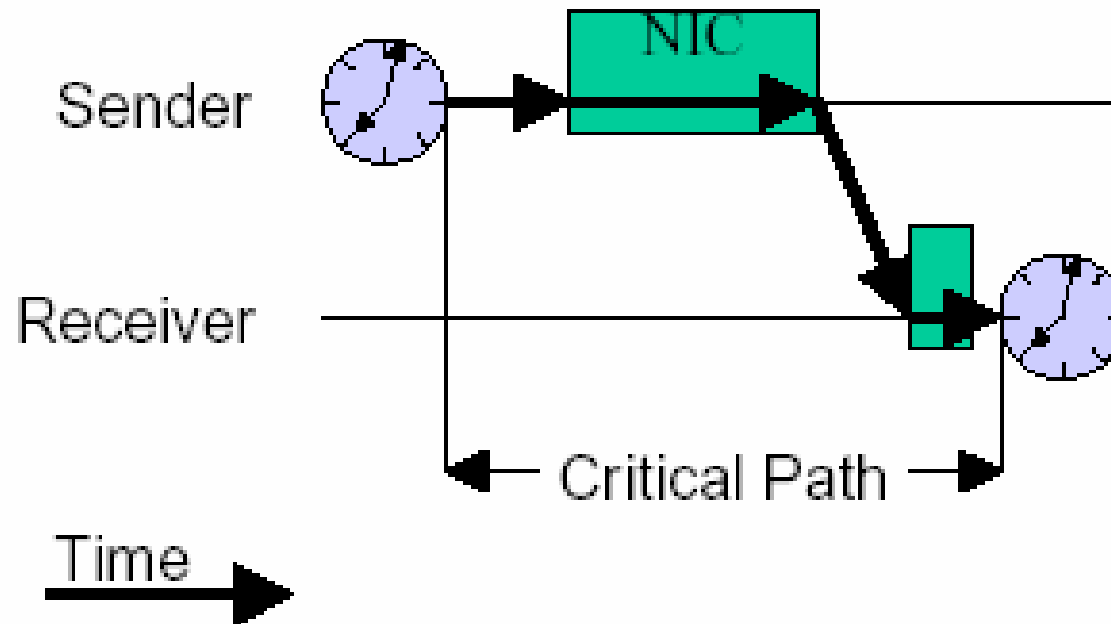
# J. Elson *et al.* – Reference Broadcasts

- Nodes send **reference beacons** to neighbors using physical-layer broadcasts
- Receivers use **arrival time** as a point of reference to compare clocks
- **Multi-hop** synchronization can be achieved

## Limitation:

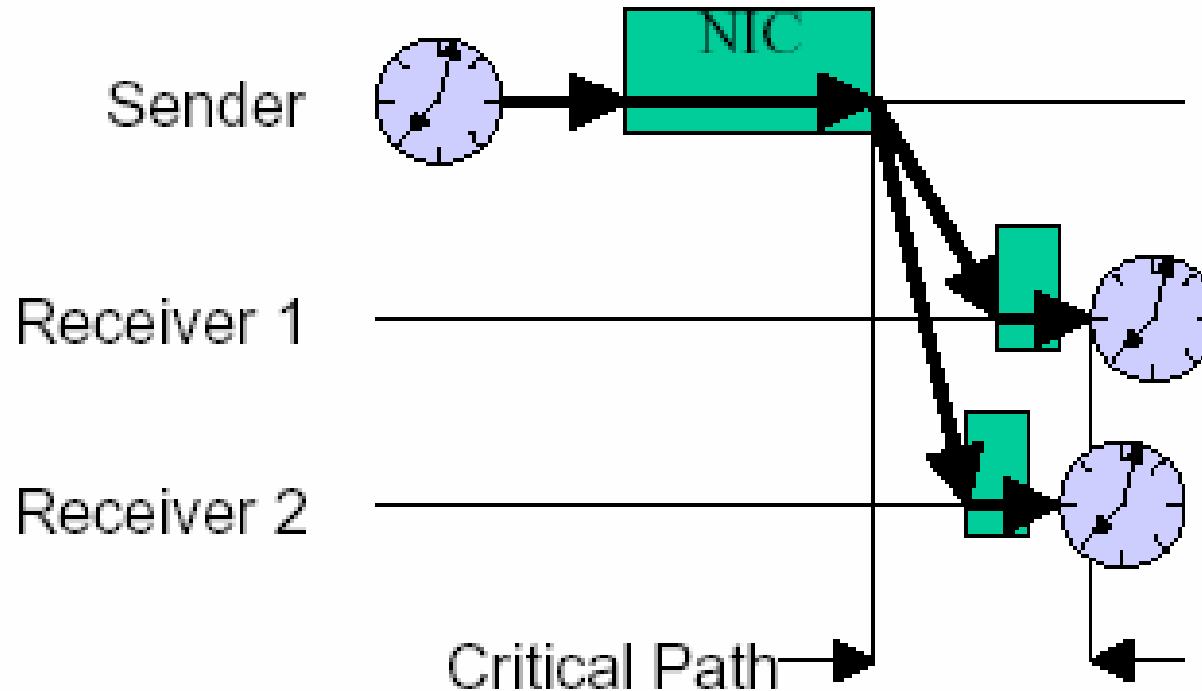
Need physical broadcast channel

# Traditional Synchronization

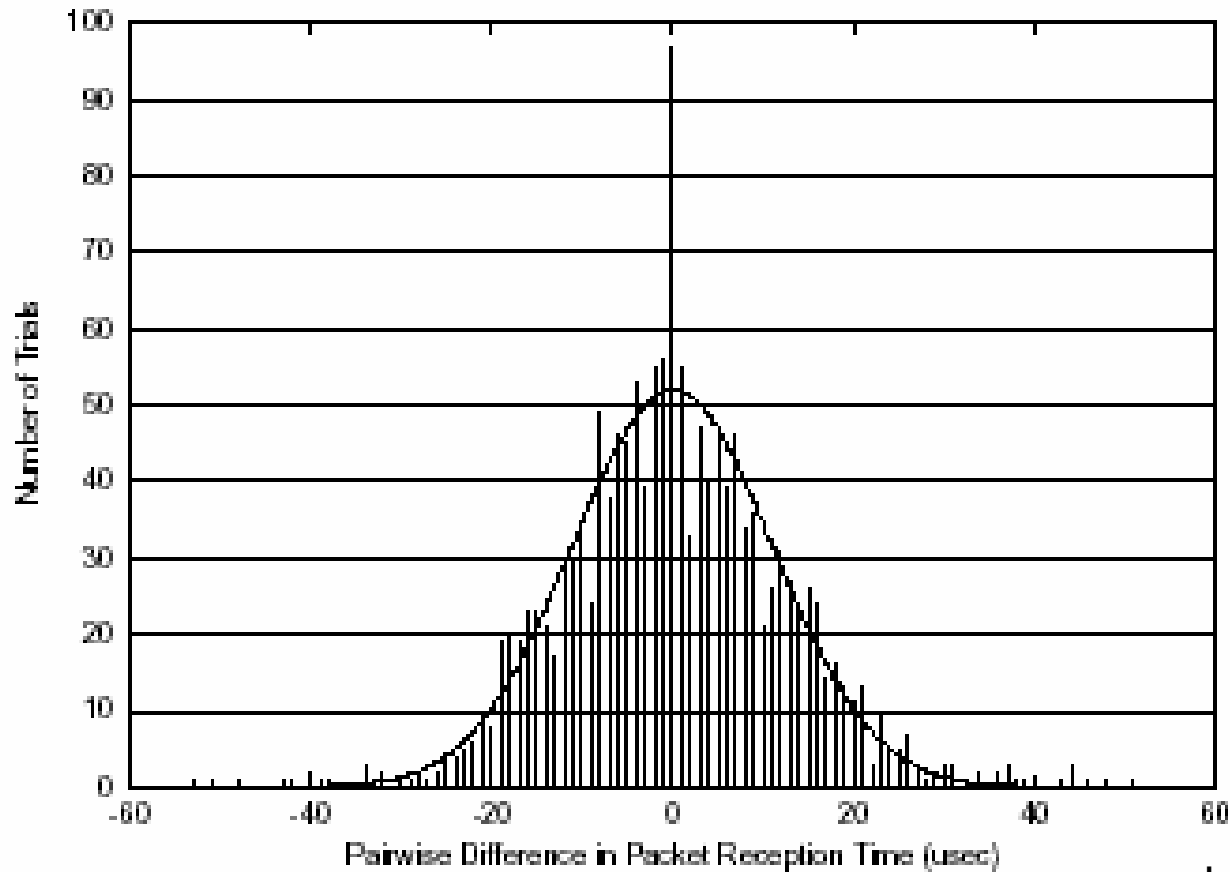




# RB Synchronization



# Inter-receiver Phase Offset



s.d. = 11.1 microsec

Bit time is 52 micro-seconds

# Estimation of Phase Offset - Algorithm

- A transmitter transmits  $m$  reference packets
- Each of the  $n$  receivers records the time that the reference was observed according to the local clock
- Receivers exchange observation
- Each pair of receivers compute the phase offset by averaging the  $m$  readings

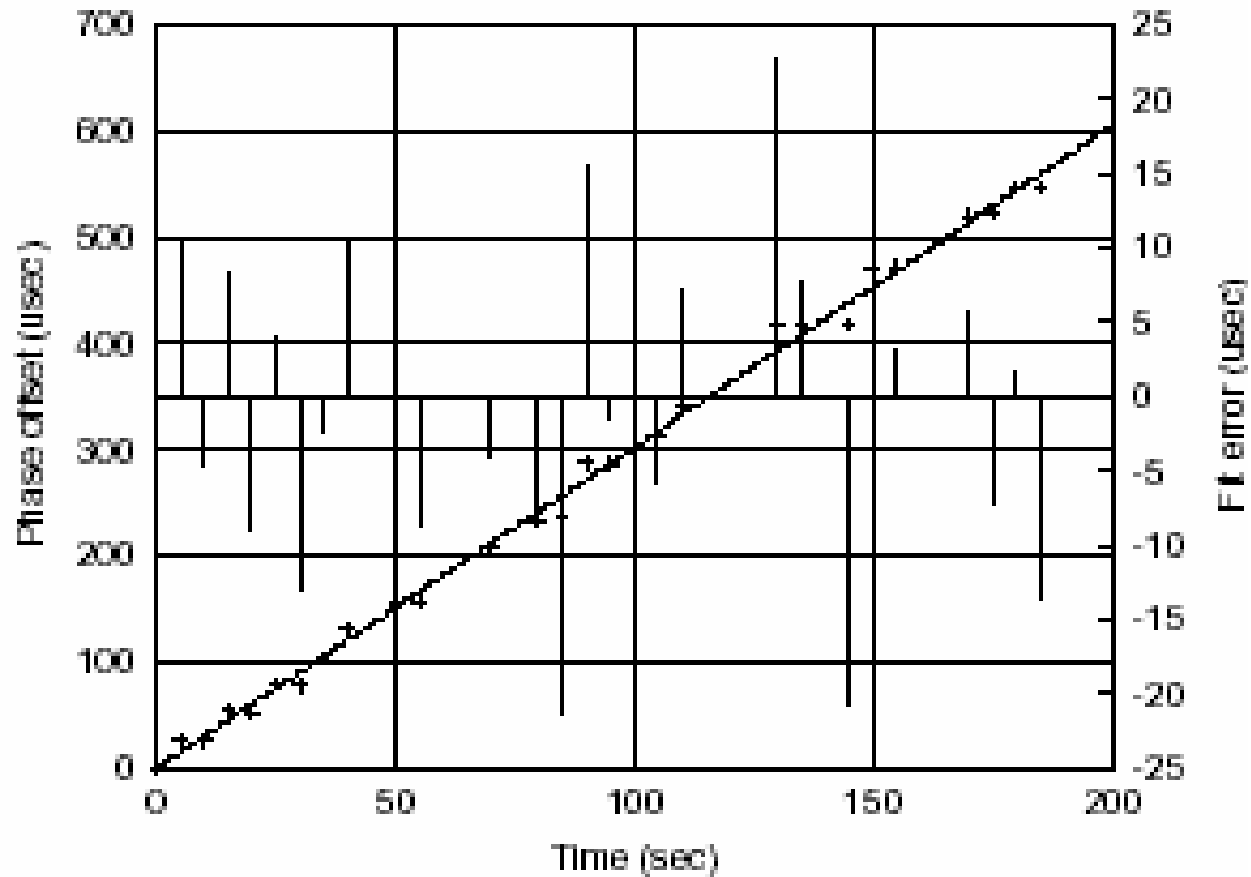
Group Dispersion: max. of phase error between all receiver pairs

# Estimation of Clock Skew

## Oscillator:

- Accuracy
  - Stability
- 
- Assume frequency is stable over small time window
  - Use **least-squares linear regression**
    - Slope gives frequency offset
    - Intercept gives phase offset

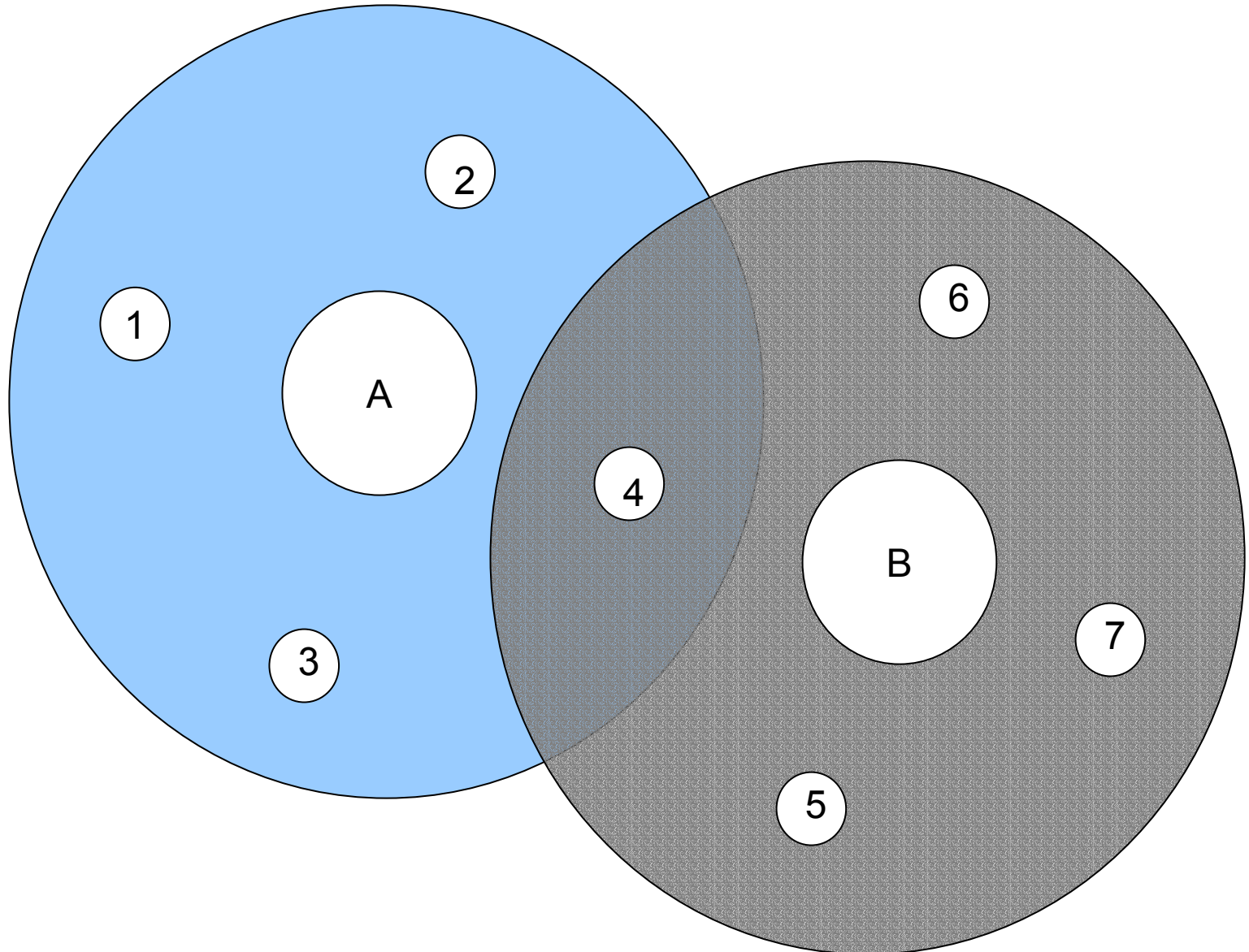
# Estimation of Clock Skew - Experiment



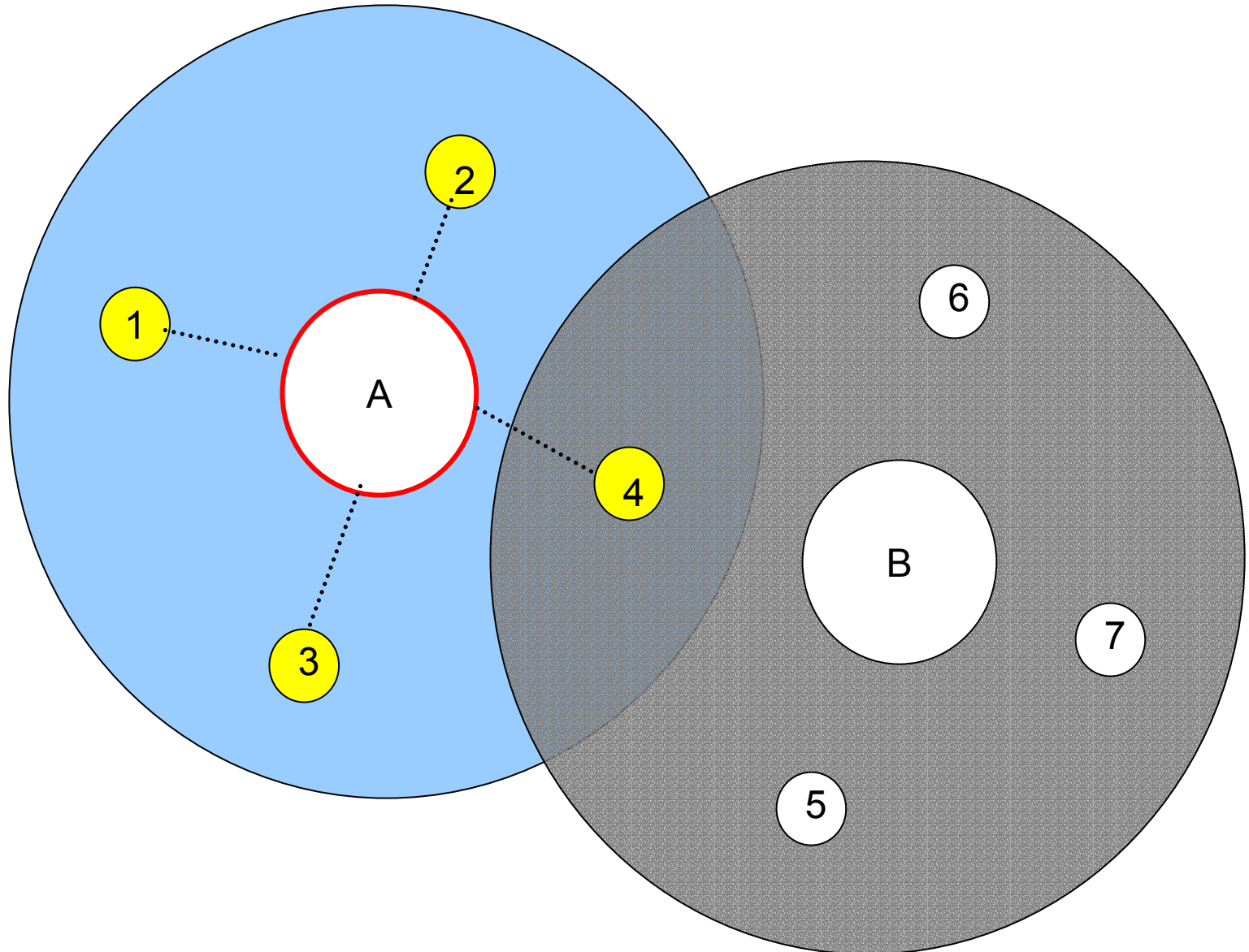
# Basic RBS Summary

- Removes largest sources of nondeterministic latency
- Multiple broadcasts decrease error, *post-facto synchronization* possible
- Outliers can be handled
- Local timescales can be constructed (absolute synchronization is a natural extension).

# Multihop Time Synchronization

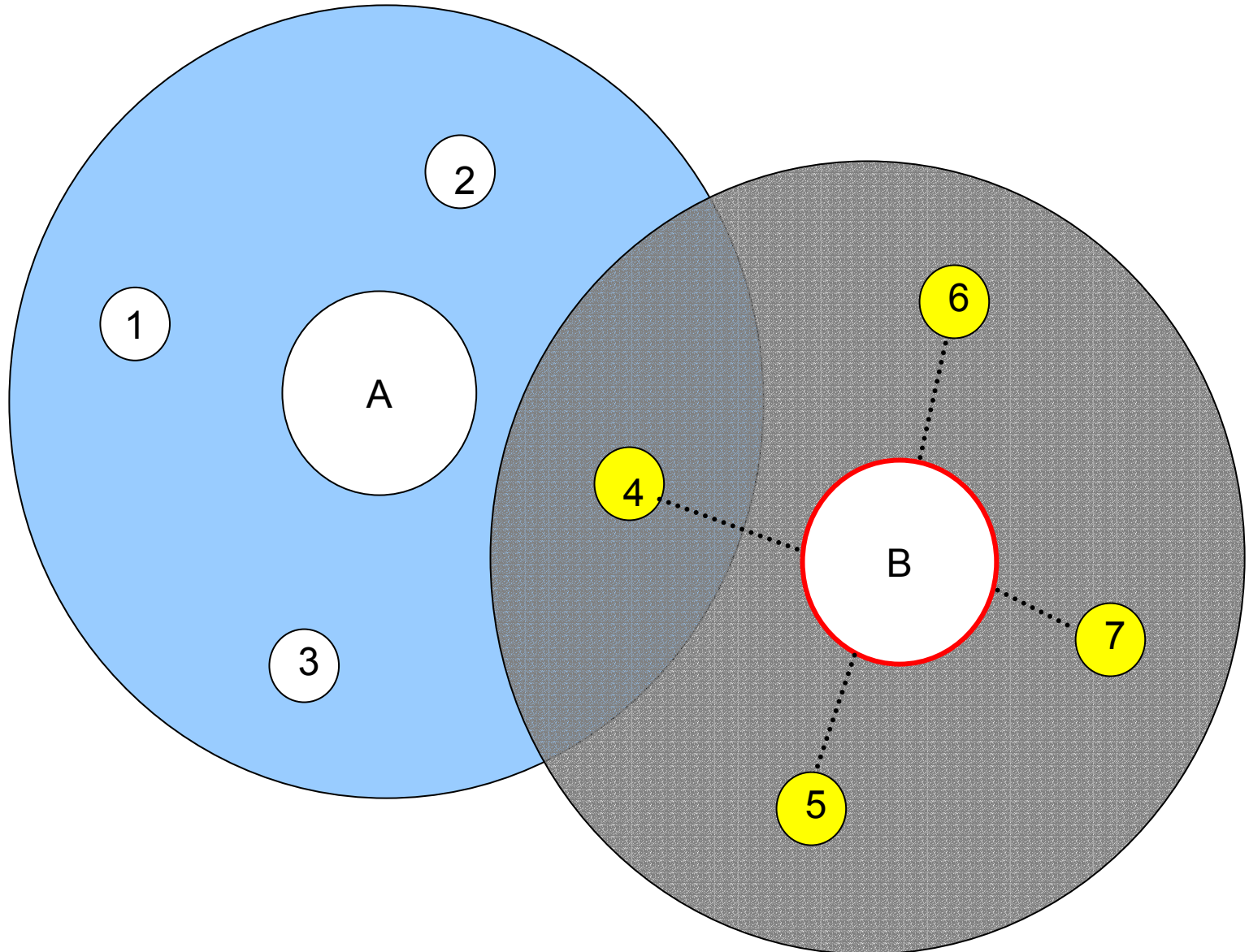


# Multihop Time Synchronization

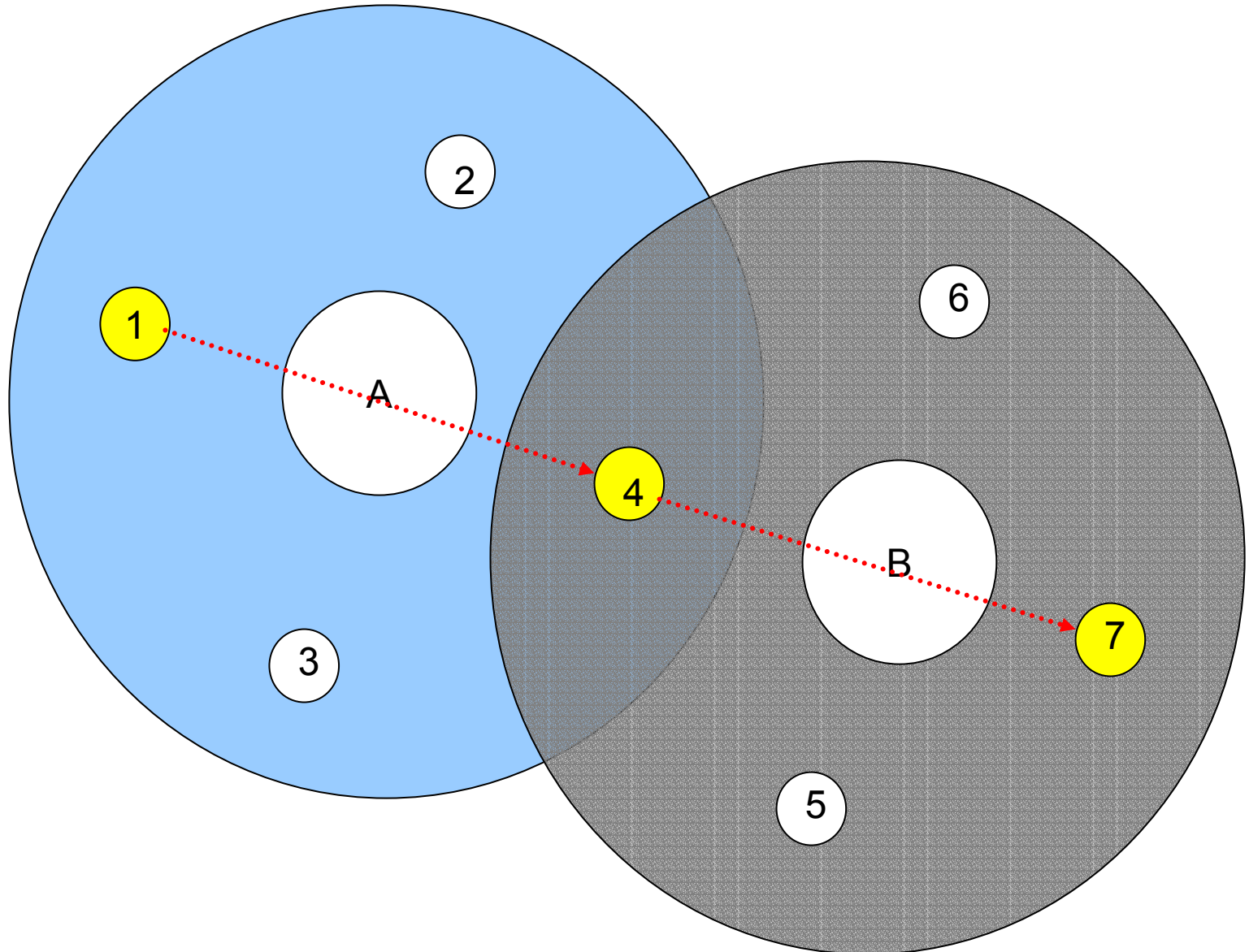




# Multihop Time Synchronization



# Multihop Time Synchronization



# Summary

- RBS synchronizes a set of receivers with one another
- No explicit timestamp needed
- Largest sources of non determinism in latency removed
- Residual error is Gaussian (experimental observation)
- *Post-facto* synchronization possible