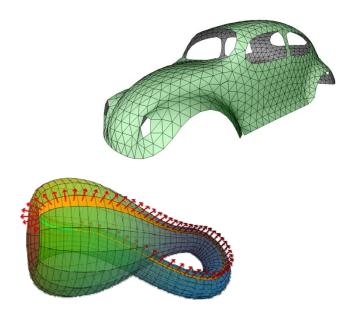




Basic Concepts



Today

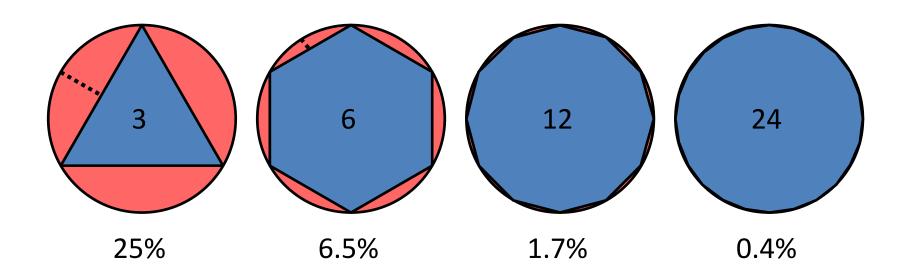
- Mesh basics
 - **–** Zoo
 - Definitions
 - Important properties
- Mesh data structures

• HW1

Polygonal Meshes

- Piecewise linear approximation
 - Error is $O(h^2)$

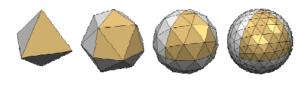




Polygonal Meshes

Polygonal meshes are a good representation

– approximation $O(h^2)$



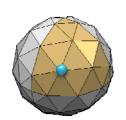
- arbitrary topology
- piecewise smooth surfaces
- adaptive refinement
- efficient rendering



Triangle Meshes

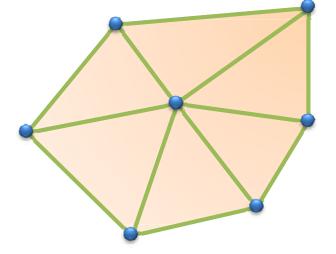
Connectivity: vertices, edges, triangles

$$\mathcal{V} = \{v_1, \dots, v_n\}$$
 $\mathcal{E} = \{e_1, \dots, e_k\} , \quad e_i \in \mathcal{V} \times \mathcal{V}$
 $\mathcal{F} = \{f_1, \dots, f_m\} , \quad f_i \in \mathcal{V} \times \mathcal{V} \times \mathcal{V}$

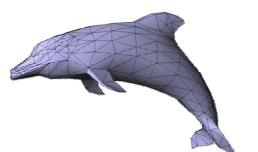


Geometry: vertex positions

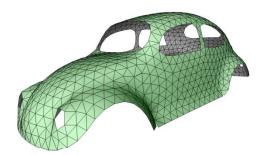
$$\mathcal{P} = \{\mathbf{p}_1, \dots, \mathbf{p}_n\}$$
, $\mathbf{p}_i \in \mathbb{R}^3$



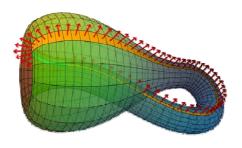
Mesh Zoo



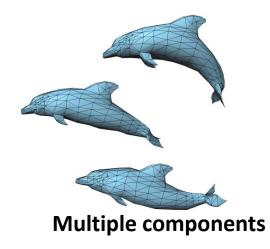
Single component, closed, triangular, orientable manifold

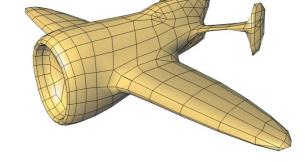


With boundaries

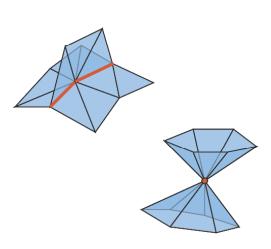


Not orientable





Not only triangles



Non manifold

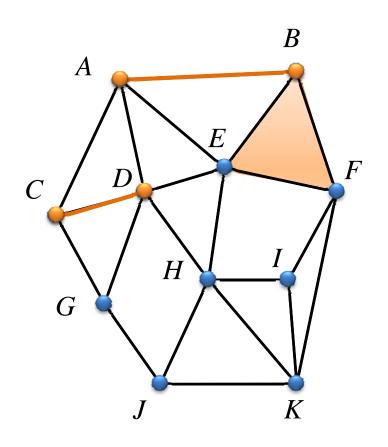
Mesh Definitions

Need vocabulary to describe zoo meshes

The connectivity of a mesh is just a graph

We'll start with some graph theory

Graph Definitions



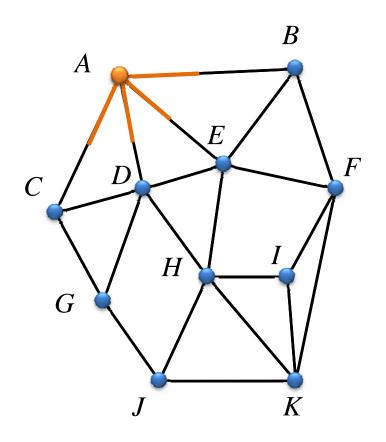
$$G = graph = \langle V, E \rangle$$

$$V = \text{vertices} = \{A, B, C, ..., K\}$$

$$E = edges = \{(AB), (AE), (CD), ...\}$$

$$F = \text{faces} = \{(ABE), (DHJG), ...\}$$

Graph Definitions



Vertex degree or **valence** = number of incident edges

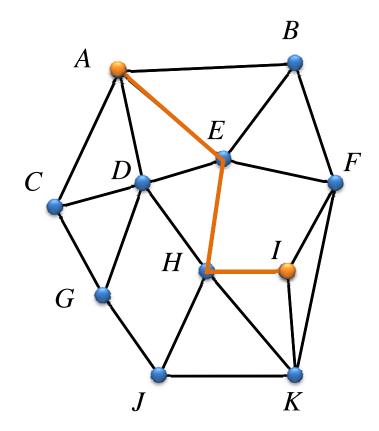
$$deg(A) = 4$$

$$deg(E) = 5$$

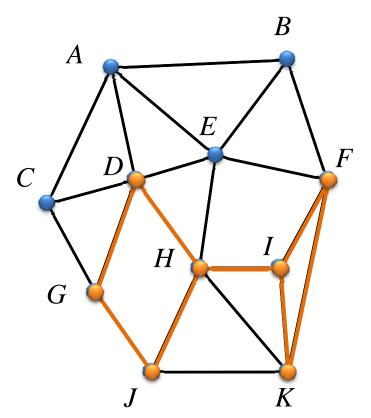
Regular mesh =

all vertex degrees are equal

Connected =



path of edges connecting every two vertices



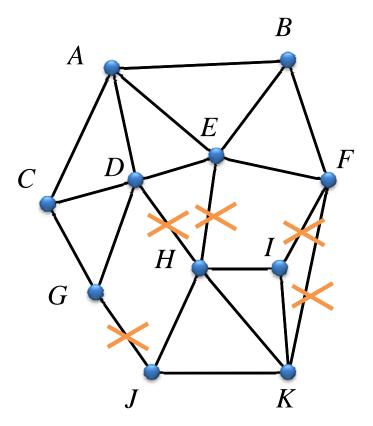
Connected =

path of edges connecting every two vertices

Subgraph =

G'=<V',E'> is a subgraph of G=<V,E> if V' is a subset of V and

E' is the subset of E incident on V

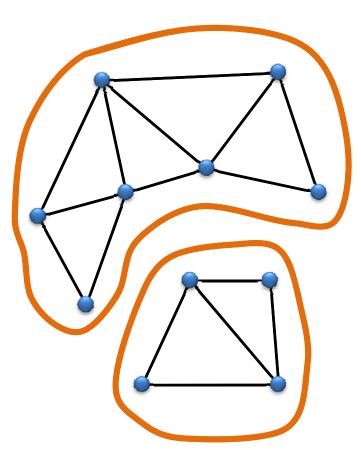


Connected =

path of edges connecting every two vertices

Subgraph =

G'=<V',E'> is a subgraph of G=<V,E> if V' is a subset of V and E' is a subset of E incident on E'



Connected =

path of edges connecting every two vertices

Subgraph =

G'=<V',E'> is a subgraph of G=<V,E> if V' is a subset of V and

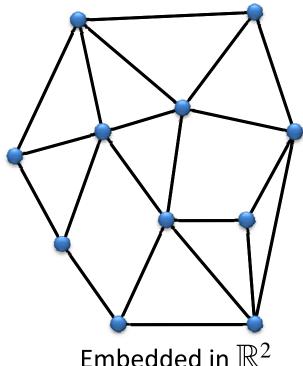
E' is the subset of E incident on V

Connected Component =

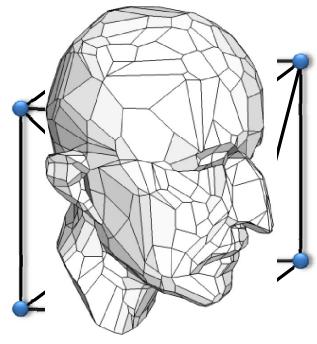
maximally connected subgraph

Graph Embedding

Embedding: G is embedded in \mathbb{R}^d , if each vertex is assigned a position in \mathbb{R}^d







Embedded in \mathbb{R}^3

Planar Graphs

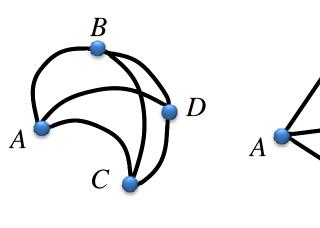
Planar Graph =

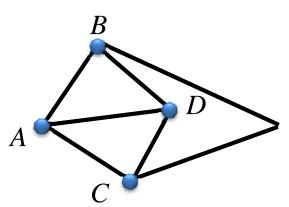
Graph whose vertices and edges can be embedded in R² such that its edges do not intersect

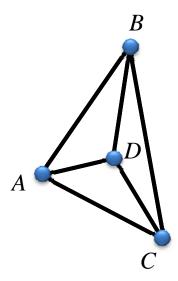
Planar Graph

Plane Graph

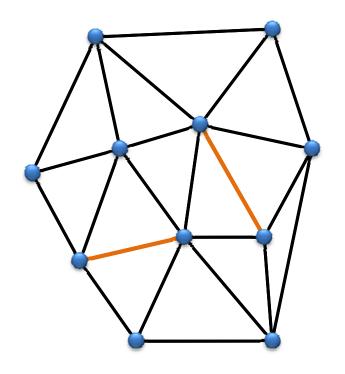
Straight Line Plane Graph







Triangulation



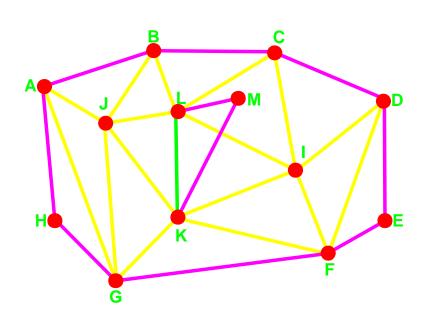
Triangulation:

Straight line plane graph where every face is a *triangle*.

Mesh

Mesh:

straight-line graph embedded in R³



Boundary edge:

adjacent to exactly one face

Regular edge:

adjacent to exactly two faces

Singular edge:

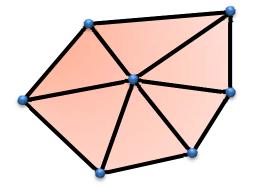
adjacent to more than two faces

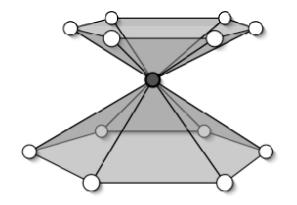
Closed mesh:

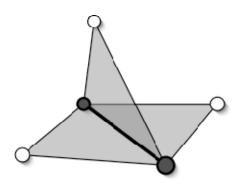
mesh with no boundary edges

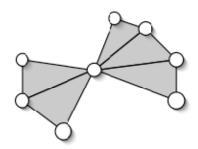
2-Manifolds Meshes

Disk-shaped neighborhoods







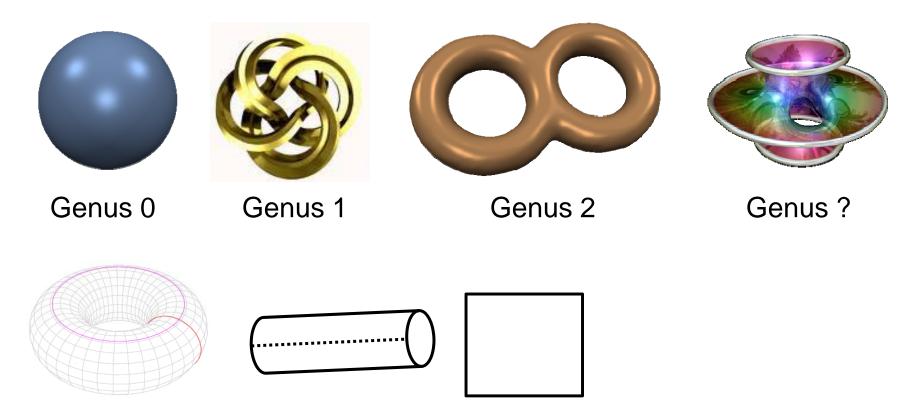


non-manifolds

Global Topology: Genus

Genus:

Half the maximal number of closed paths that do not disconnect the mesh (= the number of holes)

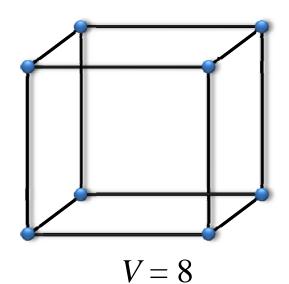


Closed 2-Manifold Polygonal Meshes

Euler-Poincaré formula

$$V+F-E=\chi$$

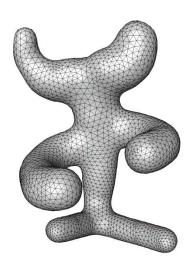
Euler characteristic



$$E=12$$

$$F = 6$$

$$\chi = 8 + 6 - 12 = 2$$



$$V = 3890$$

$$E = 11664$$

$$F = 7776$$

$$\chi = 2$$

Closed 2-Manifold Polygonal Meshes

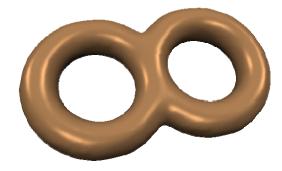
Euler-Poincaré formula

$$V + F - E = \chi = 2$$



$$V = 1500, E = 4500$$

 $F = 3000, \mathbf{g} = \mathbf{1}$
 $\chi = \mathbf{0}$



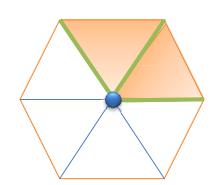
$$\mathbf{g} = \mathbf{2}$$
$$\chi = -2$$

Closed 2-Manifold Triangle Meshes

• *Triangle* mesh statistics

$$E \approx 3V$$

$$F \approx 2V$$



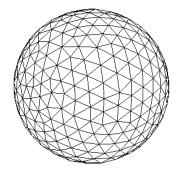
Avg. valence ≈ 6
 Show using Euler Formula



 When can a closed triangle mesh be 6regular?



Exercise



Theorem: Average vertex degree in closed manifold triangle mesh is ~6

Proof: In such a mesh, 3F = 2E by counting edges of faces.

By Euler's formula: V+F-E = V+2E/3-E = 2-2g. Thus E = 3(V-2+2g)

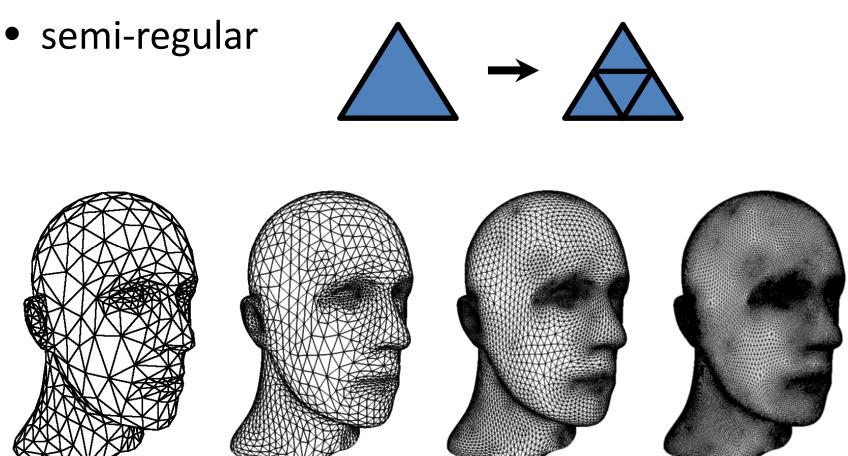
So Average(deg) = $2E/V = 6(V-2+2g)/V \sim 6$ for large V



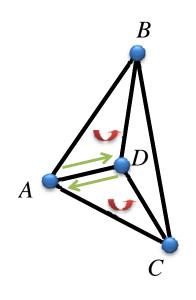
Corollary: Only toroidal (g=1) closed manifold triangle mesh can be regular (all vertex degrees are 6)

Proof: In regular mesh average degree is *exactly* 6. Can happen only if g=1

Regularity



Orientability



Face Orientation =

clockwise or anticlockwise order in which the vertices listed

defines direction of face normal

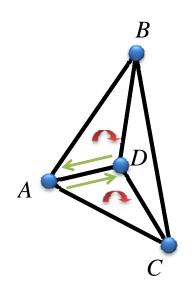
Oriented **CCW**:

{(C,D,A), (A,D,B), (C,B,D)}

Oriented CW:

 $\{(C,A,D), (D,A,B), (B,C,D)\}$

Orientability



Face Orientation =

clockwise or anticlockwise order in which the vertices listed

defines direction of face normal

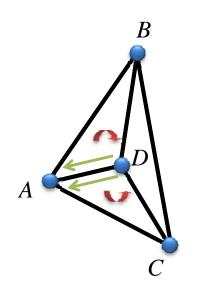
Oriented **CCW**: {(C,D,A), (A,D,B), (C,B,D)}

Oriented CW:

{(C,A,D), (D,A,B), (B,C,D)}

Not oriented: {(C,D,A), (D,A,B), (C,B,D)}

Orientability



Face Orientation =

clockwise or anticlockwise order in which the vertices listed

defines direction of face normal

Oriented **CCW**: {(C,D,A), (A,D,B), (C,B,D)}

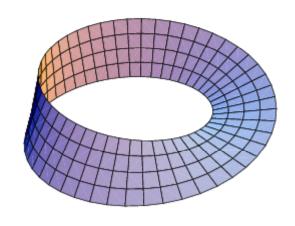
Oriented **CW**: {(C,A,D), (D,A,B), (B,C,D)}

Not oriented: {(C,D,A), (D,A,B), (C,B,D)}

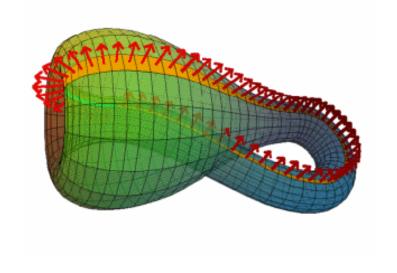
Orientable Plane Graph =

orientations of faces can be chosen so that each non-boundary edge is oriented in *both* directions

Non-Orientable Surfaces

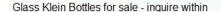


Mobius Strip



Klein Bottle

Garden Variety Klein Bottles



Need a zero-volume bottle?

Searching for a one-sided surface?

Want the ultimate in non-orientability?

Get an ACME KLEIN BOTTLE!



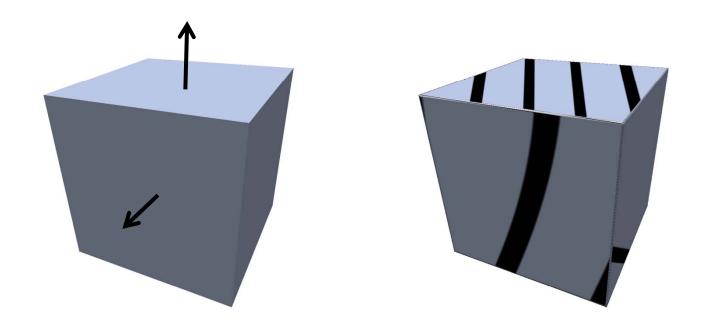




http://www.kleinbottle.com/

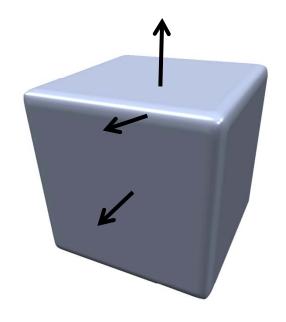
Smoothness

Position continuity = C⁰



Smoothness

- Position continuity = C⁰
- Tangent continuity $\approx C^1$





Smoothness

- Position continuity = C⁰
- Tangent continuity ≈ C¹
- Curvature continuity $\approx C^2$

