

Ray Tracing

Ray Tracing 1

- Basic algorithm
- Overview of pbrt
- Ray-surface intersection (triangles, ...)

Ray Tracing 2

- Problem: brute force = $|\text{Image}| \times |\text{Objects}|$
- Acceleration data structures

Primitives

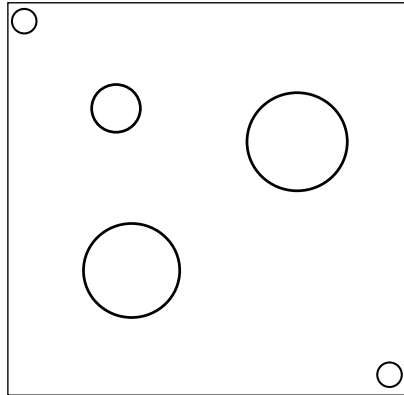
pbrt primitive base class

- Shape
- Material (reflection and emission)

Subclasses

- Primitive instance
 - Transformation and pointer to a primitive
- Aggregate (collection)
 - Treat collections just like single primitives
 - Incorporate acceleration structures into collections
 - May nest accelerators of different types
 - Types: grid.cpp and kdtree.cpp

Uniform Grids



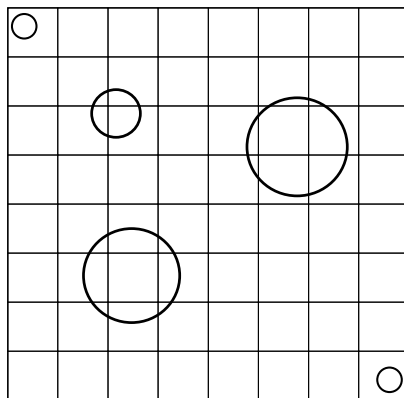
Preprocess scene

1. Find bounding box

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Uniform Grids



Preprocess scene

1. Find bounding box

2. Determine resolution

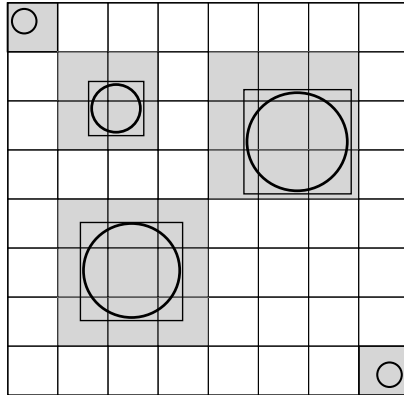
$$n_v = n_x n_y n_z \propto n_o$$

$$\max(n_x, n_y, n_z) = d \sqrt[3]{n_o}$$

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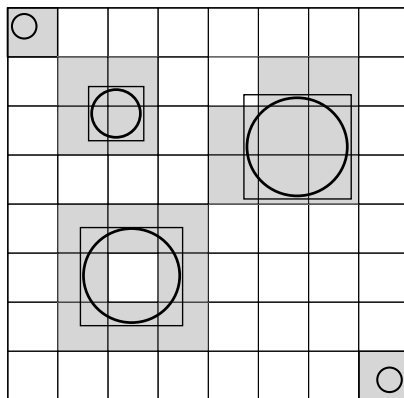
Preprocess scene

1. Find bounding box
2. Determine resolution
$$\max(n_x, n_y, n_z) = d\sqrt[3]{n_o}$$
3. Place object in cell,
if object overlaps cell

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Uniform Grids



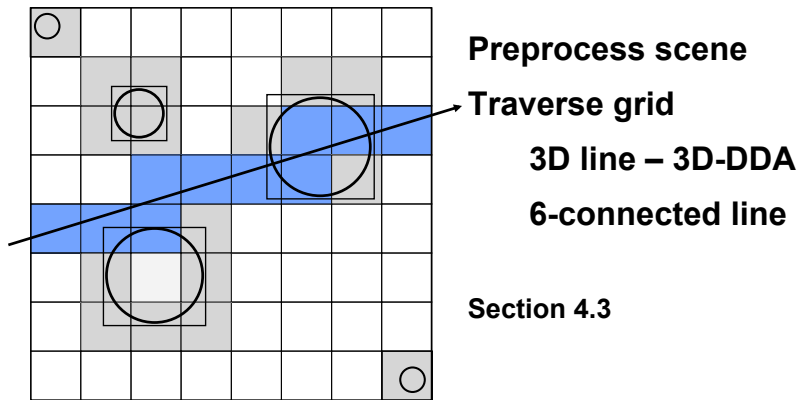
Preprocess scene

1. Find bounding box
2. Determine resolution
$$\max(n_x, n_y, n_z) = d\sqrt[3]{n_o}$$
3. Place object in cell,
if object overlaps cell
4. Check that object's
surface intersects cell

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Uniform Grids

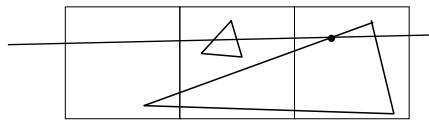


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Caveat: Overlap

Problem: Don't output first intersection found!

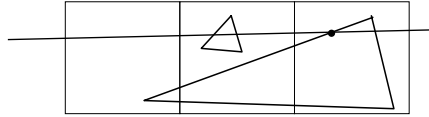


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Caveat: Overlap

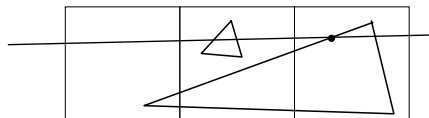
Problem: Don't output first intersection found!



Problem: Redundant intersection tests

Caveat: Overlap

Problem: Don't output first intersection found!

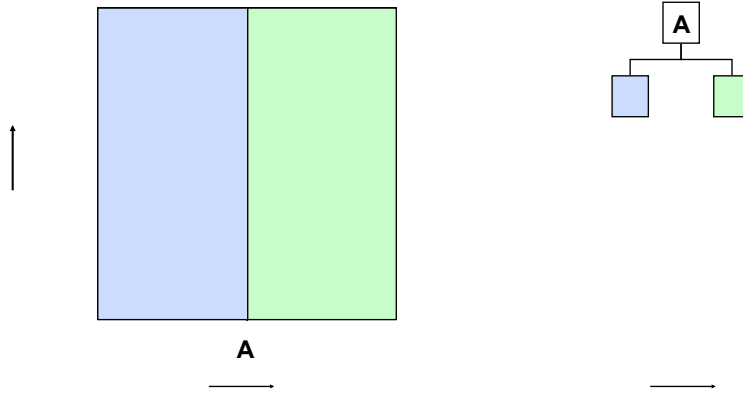


Problem: Redundant intersection tests

Solution: Mailboxes

- Assign each ray an increasing number
- Primitive intersection cache (mailbox)
 - Store last ray number tested in mailbox
 - Only intersect if ray number is greater

Spatial Hierarchies

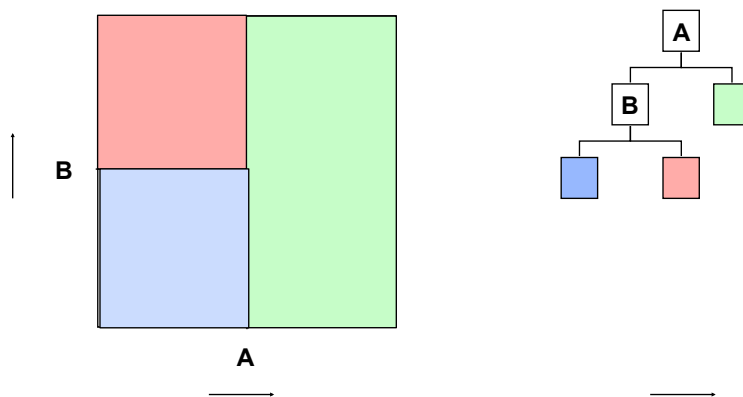


Letters correspond to planes (A)

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Spatial Hierarchies

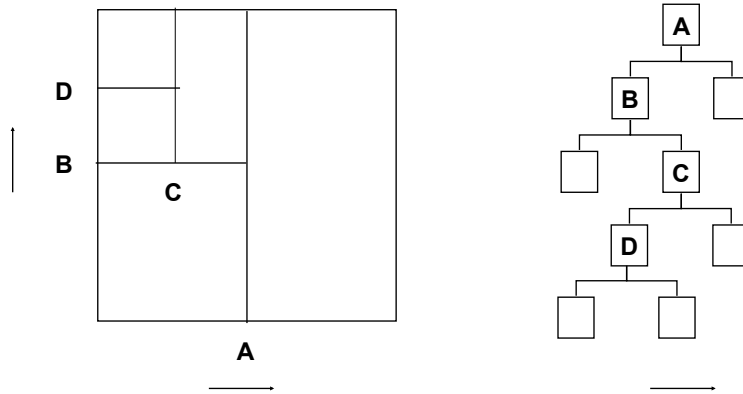


Letters correspond to planes (A, B)
Point Location by recursive search

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Spatial Hierarchies

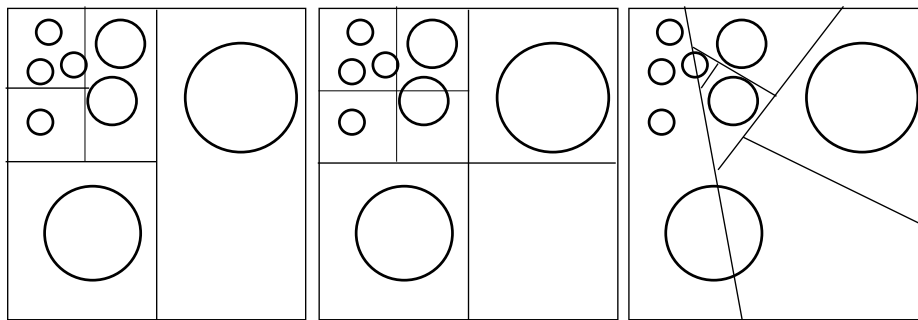


Letters correspond to planes (A, B, C, D)

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Variations



kd-tree

oct-tree

bsp-tree

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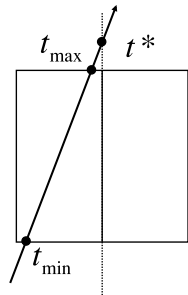
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Ray Traversal Algorithms

Recursive inorder traversal

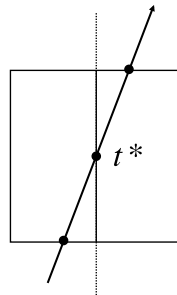
[Kaplan, Arvo, Jansen]

$$t^* = (S - O[a]) / D[a]$$



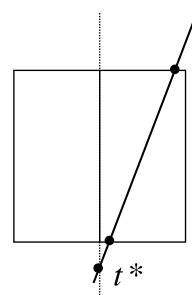
$$t_{\max} < t^*$$

Intersect(L, tmin, tmax)



$$t_{\min} < t^* < t_{\max}$$

Intersect(L, tmin, t*)
Intersect(R, t*, tmax)



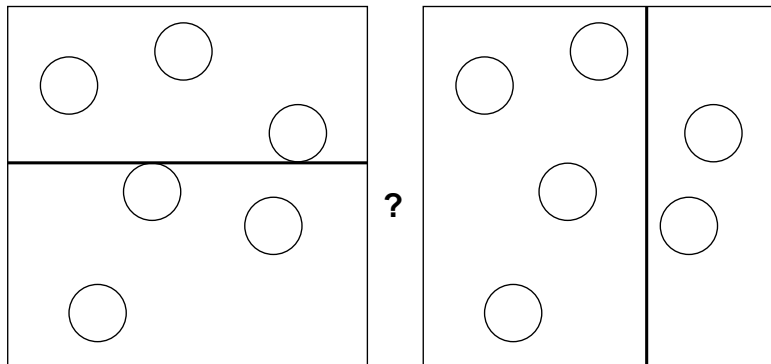
$$t^* < t_{\min}$$

Intersect(R, tmin, tmax)

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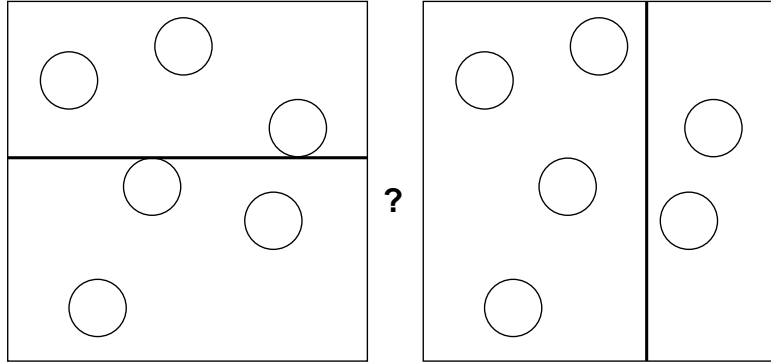
How to Build the Hierarchy?



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Build Hierarchy Top-Down



Methods to choose axis and splitting plane

- Midpoint
- Median cut (balanced)
- Surface area heuristic

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Cost

What is the cost of tracing a ray through a node?

$$\text{Cost}(\text{node}) = C_{\text{trav}} + \text{Prob}(\text{hit L}) * \text{Cost}(\text{L}) + \text{Prob}(\text{hit R}) * \text{Cost}(\text{R})$$

C_{trav} = cost of traversing a cell

$\text{Cost}(\text{L})$ = cost of traversing left child

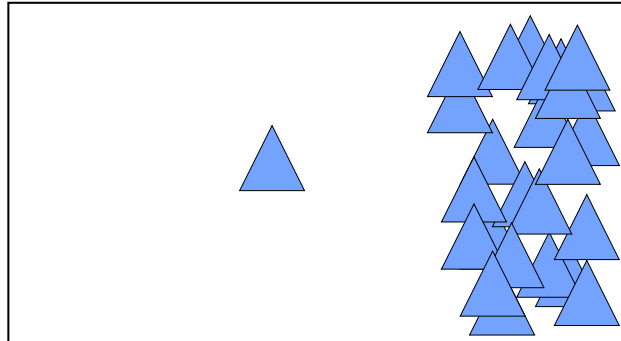
$\text{Cost}(\text{R})$ = cost of traversing right child

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Splitting with Cost in Mind

From Gordon Stoll

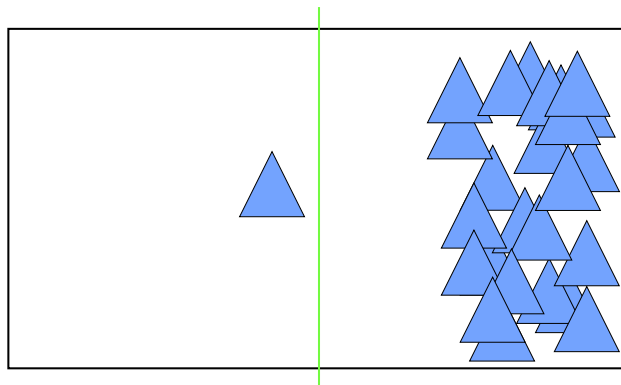


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Split in the Middle = Bad!

From Gordon Stoll



Makes the L & R probabilities equal

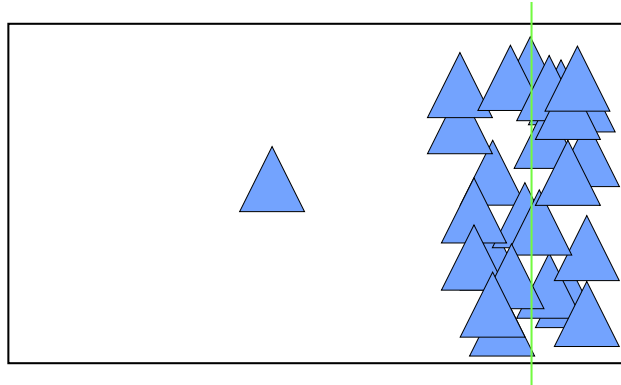
Cost of R greater than cost of L

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Split at the Median = Bad!

From Gordon Stoll



Makes the L & R costs equal

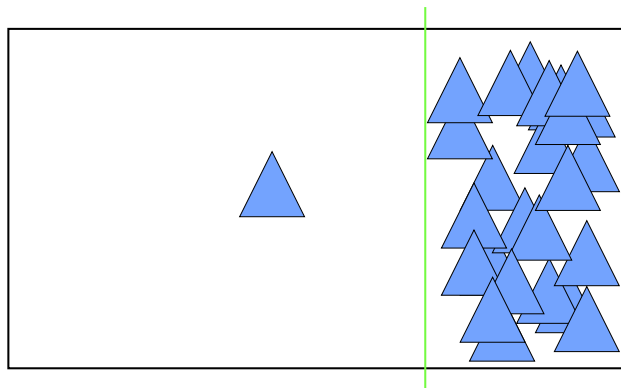
Probability of hitting L greater than R

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Cost-Optimized Split = Good!

From Gordon Stoll



$$\text{Cost}(\text{cell}) = C_{\text{trav}} + \text{Prob}(\text{hit L}) * \text{Cost}(\text{L}) + \text{Prob}(\text{hit R}) * \text{Cost}(\text{R})$$

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Cost

Need the probabilities

- Turns out to be proportional to surface area

Need the child cell costs

- Triangle count is a good approximation

$$\text{Cost}(\text{cell}) = C_{\text{trav}} + \text{SA}(\text{L}) * \text{TriCount}(\text{L}) + \text{SA}(\text{R}) * \text{TriCount}(\text{R})$$

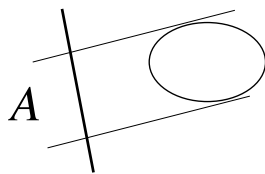
C_{trav} is the ratio of the cost to traverse to the cost to intersect

$$C_{\text{trav}} = 1:80 \text{ in pbrt}$$

$$C_{\text{trav}} = 1:1.5 \text{ in a highly optimized version}$$

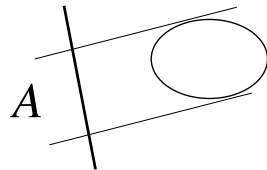
Projected Area and Ray Intersection

Number of rays in a given direction that hit an object is proportional to its projected area



Projected Area and Surface Area

Number of rays in a given direction that hit an object is proportional to its projected area



The total number of rays hitting an object is $4\pi\bar{A}$

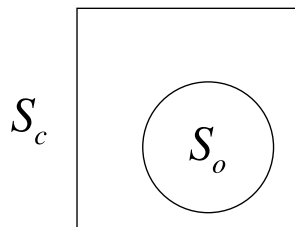
Crofton's Theorem:

For a convex body $\bar{A} = \frac{S}{4}$

For a sphere $S = 4\pi r^2$ and $\bar{A} = A = \pi r^2$

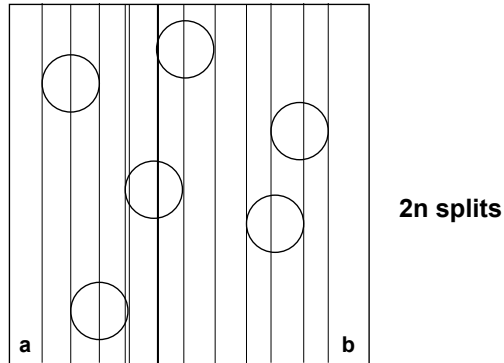
Surface Area and Ray Intersection

The probability of a ray hitting a convex shape enclosed by another convex shape is



$$\Pr[r \cap S_o | r \cap S_c] = \frac{S_o}{S_c}$$

Sweep Build Algorithm



$$p_a = \frac{S_a}{S}$$
$$N_a$$

$$p_b = \frac{S_b}{S}$$
$$N_b$$

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Basic Build Algorithm (Triangles)

1. Pick an axis, or optimize across x, y, z
2. Build a set of “candidate” split locations
 - Note: Cost extrema must be at bbox vertices
 - Vertices of triangle
 - Vertices of triangle clipped to node bbox
3. Sort the triangles into intervals
4. Sweep to incrementally track L/R counts, cost
5. Output position of minimum cost split

Running time: $T(N) = N \log N + 2T(N/2)$

$$T(N) = N \log^2 N$$

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Termination Criteria

When should we stop splitting?

- Bad: depth limit, number of triangles
- Good: When split does not lower the cost

Threshold of cost improvement




- Stretch over multiple levels
- For example, if cost doesn't go down after three splits in a row, terminate

Threshold of cell size

- Absolute probability $SA(\text{node})/SA(\text{scene})$ small

Best Reported Timings

Millions of Rays per Second

Framerate (FPS) @ 1024x1024 resolution			OpenRT @	MLRTA @	MLRTA @
			2.5 GHz P4	2.4 GHz P4	3.2 GHz P4 with HT
scene # of triangles and shader (+/-)			1 thread	1 thread	2 threads
Erw6 804		- shader	7.1	70.2	109.8
		+ shader	2.3	37.8	50.7
Confe- rence 274K		- shader	4.55	11.2	19.5
		+ shader	1.93	9.5	15.6
Soda Hall 2195K		- shader	4.12	21.1	35.5
		+ shader	1.8	15.3	24.1

Reshetov, Soupikov, Hurley, SIGGRAPH 2005

Superoptimizations

Lots of optimizations

- Carefully written inner loop (no recursion)
- Use vector instructions SSE2
- 64 bits per kd-tree node
 - 32 bit position
 - 32 bit pointer to pair of child nodes
 - 2 bits for split plane direction (x, y, or z)
- Trace packet of rays
 - 4 or more rays at a time
- Intersect beam at top of tree
- Encourage empty nodes
- Special case axis-aligned triangles
- ...

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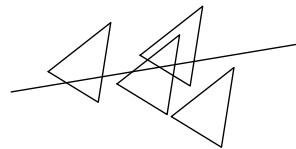
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Theoretical Nugget 1

Computational geometry of ray shooting

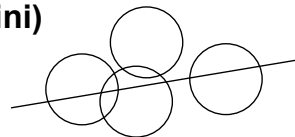
1. Triangles (Pellegrini)

- Time: $O(\log n)$
- Space: $O(n^{5+\epsilon})$



2. Sphere (Guibas and Pellegrini)

- Time: $O(\log^2 n)$
- Space: $O(n^{5+\epsilon})$



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Theoretical Nugget 2

Optical computer = Turing machine

Reif, Tygar, Yoshida

Determining if a ray
starting at y_0 arrives
at y_n is undecidable

